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My Experience

Monday, December 24, 2012

Segment Trees and lazy propagation

In this topic i will explain a very interesting data structure that can be used to solve a specific set of problems. I will start by explaining its definition and the proceeding with an example problem to solve with it.

Table of contents:

- What is segment trees?
- · Order of growth of segment trees operations
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What is segment trees?

Segment Trees is a Tree data structure for storing intervals, or segments, It allows querying which of the stored segments contain a given point. It is, in principle, a static structure; that is, its content cannot be modified once the structure is built. It only uses O(N lg(N)) storage.

A segment trees has only three operations: build tree, update tree, query tree.

Building tree: To init the tree segments or intervals values **Update tree:** To update value of an interval or segment Query tree: To retrieve the value of an interval or segment

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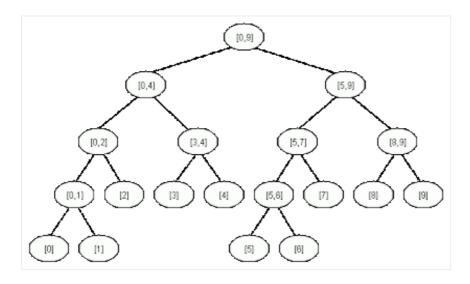
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2014 (21)

Example Segment Tree:

- The first node will hold the information for the interval [i, j]
- If i<j the left and right son will hold the information for the intervals [i, (i+j)/2] and [(i+j)/2+1, j]

Notice that the height of a segment tree for an interval with N elements is [logN] + 1. Here is how a segment tree for the interval [0, 9] would look like:



Order of growth of segment trees operations

build_tree: O(N lg(N)) • update_tree: O(lg(N + k)) query_tree: O(lg(N + k))

K = Number of retrieved intervals or segments

Show me your code

```
* In this code we have a very large array called arr, and very large set of
3 * Operation #1: Increment the elements within range [i, j] with value val
4 * Operation #2: Get max element within range [i, j]
  * Build tree: build tree(1, 0, N-1)
   * Update tree: update tree(1, 0, N-1, i, j, value)
   * Query tree: query tree(1, 0, N-1, i, j)
```

- **2013 (4)**
- **2012 (17)**
 - ▼ December (4)

Segment Trees and lazy propagation

Lucas' Theorem to solve binomial coefficients

Hacking linux keyboard

Apache JMeter along with jsf pages

- October (3)
- ► September (1)
- **▶** July (1)
- ► April (5)
- ► January (3)
- **2011 (10)**

About Me

Hussein El-Sayed

Cairo, Egypt

Software engineer at Vodafone international services, topcoder handle is husseincoder

View my complete profile

```
10 #include<iostream>
11 #include<algorithm>
12 using namespace std;
14 | #include<string.h>
15 | #include<math.h>
16
17 #define N 20
18 | #define MAX (1+(1<<6)) // Why? :D
19 #define inf 0x7fffffff
21 | int arr[N];
22 int tree[MAX];
24 /**
25 * Build and init tree
27 | void build tree(int node, int a, int b) {
        if(a > b) return; // Out of range
29
           if(a == b) { // Leaf node
31
                    tree[node] = arr[a]; // Init value
                    return;
           }
34
            build tree (node*2, a, (a+b)/2); // Init left child
            build tree(node*2+1, 1+(a+b)/2, b); // Init right child
            tree[node] = max(tree[node*2], tree[node*2+1]); // Init root value
39 }
40
41 /**
42 * Increment elements within range [i, j] with value value
43 */
44 void update tree(int node, int a, int b, int i, int j, int value) {
45
46
           if(a > b \mid \mid a > j \mid \mid b < i) // Current segment is not within range [
47
                    return;
48
49
           if(a == b) { // Leaf node}
                    tree[node] += value;
51
                    return;
            }
54
            update tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
```

```
update tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right of
            tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root with
58 }
59
60 /**
* Query tree to get max element value within range [i, j]
63 | int query tree(int node, int a, int b, int i, int j) {
64
           if (a > b \mid | a > j \mid | b < i) return -inf; // Out of range
67
            if(a >= i && b <= j) // Current segment is totally within range [i,
68
                    return tree[node];
            int q1 = query tree(node*2, a, (a+b)/2, i, j); // Query left child
            int q2 = query tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right chi
            int res = max(q1, q2); // Return final result
74
           return res;
76 }
78 | int main() {
79
            for (int i = 0; i < N; i++) arr[i] = 1;
81
            build tree(1, 0, N-1);
82
            update tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
84
            update tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12
            update tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by
            cout << query tree(1, 0, N-1, 0, N-1) << endl; // Get max element in</pre>
88 }
segment_tree.cpp hosted with by GitHub
                                                                         view raw
```

Lazy Propagation

Sometimes a segment tree operation wouldn't survive if the problem constraints is too large, here it come lazy propagation along with the segment tree.

In the current version when we update a range, we branch its childs even if the segment is covered

within range. In the lazy version we only mark its child that it needs to be updated and update it when needed.

Note: Read my solution for problem Can you answer these queries I in this article.

```
1 /**
2 * In this code we have a very large array called arr, and very large set of
3 * Operation #1: Increment the elements within range [i, j] with value val
4 * Operation #2: Get max element within range [i, j]
5 * Build tree: build tree(1, 0, N-1)
6 * Update tree: update tree(1, 0, N-1, i, j, value)
7 * Query tree: query tree(1, 0, N-1, i, j)
8 */
9
10 #include<iostream>
11 #include<algorithm>
12 using namespace std;
14 | #include<string.h>
15 | #include<math.h>
16
17 | #define N 20
18 | #define MAX (1+(1<<6)) // Why? :D
19 #define inf 0x7fffffff
21 int arr[N];
22 int tree [MAX];
23 int lazy[MAX];
24
25 /**
26 * Build and init tree
28 void build tree (int node, int a, int b) {
           if(a > b) return; // Out of range
           if(a == b) { // Leaf node}
                   tree[node] = arr[a]; // Init value
                   return;
34
           build tree (node*2, a, (a+b)/2); // Init left child
           build tree(node*2+1, 1+(a+b)/2, b); // Init right child
39
           tree[node] = max(tree[node*2], tree[node*2+1]); // Init root value
```

```
40 }
41
42 /**
    * Increment elements within range [i, j] with value value
    * /
   void update tree(int node, int a, int b, int i, int j, int value) {
46
47
            if(lazy[node] != 0) { // This node needs to be updated
48
                    tree[node] += lazy[node]; // Update it
49
                    if(a != b) {
51
                            lazy[node*2] += lazy[node]; // Mark child as lazy
                            lazy[node*2+1] += lazy[node]; // Mark child as lazy
53
                    }
54
                    lazy[node] = 0; // Reset it
57
            if(a > b || a > j || b < i) // Current segment is not within range
59
                    return;
61
            if(a >= i && b <= j) { // Segment is fully within range</pre>
62
                    tree[node] += value;
64
                    if(a != b) { // Not leaf node
65
                            lazy[node*2] += value;
                            lazy[node*2+1] += value;
67
69
                    return;
71
72
            update tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
           update tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right
74
            tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root with
76 }
78 /**
79 * Query tree to get max element value within range [i, j]
81 int query tree(int node, int a, int b, int i, int j) {
83
           if (a > b \mid | a > j \mid | b < i) return -inf; // Out of range
84
```

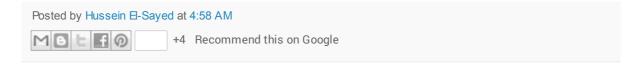
```
85
             if(lazy[node] != 0) { // This node needs to be updated
 86
                     tree[node] += lazy[node]; // Update it
                     if(a != b) {
 89
                             lazy[node*2] += lazy[node]; // Mark child as lazy
                             lazy[node*2+1] += lazy[node]; // Mark child as lazy
                     lazy[node] = 0; // Reset it
 94
             if (a \ge i \&\& b \le j) // Current segment is totally within range [i,
                     return tree[node];
             int q1 = query tree(node*2, a, (a+b)/2, i, j); // Query left child
             int q2 = query tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right ch
             int res = max(q1, q2); // Return final result
104
             return res;
105 }
107 | int main() {
             for (int i = 0; i < N; i++) arr[i] = 1;
110
             build tree (1, 0, N-1);
             memset(lazy, 0, sizeof lazy);
113
114
             update tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
             update tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12
116
             update tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1]
117
118
             cout << query tree(1, 0, N-1, 0, N-1) << endl; // Get max element in
119 }
lazy_segment_tree.cpp hosted with by GitHub
                                                                         view raw
```

Sample Problems to try

- Quadrant Queries
- D-Query
- Can You answer these queries I

References

- Wiki
- Topcoder tutorials



26 comments:



donnghi December 24, 2012 at 12:25 PM

In lazy's version, i thinks it's better if you replace update tree[2*node] and tree[2*node+1] in 49th, 50th, 80th and 81th line by lazy[2*node] and lazy[2*node+1].

Its reason is your query is not really come down to higher level, so lazy[] should be updated Reply



Hussein El-Sayed December 24, 2012 at 10:44 PM

I can't understand you:)

Reply



donnghi December 25, 2012 at 3:03 AM

So, in line 80th: tree[node*2] += lazy[node]; // Mark child as lazy tree[node*2+1] += lazy[node]; // Mark child as lazy

=> replaced by: lazy[node*2] += lazy[node]; lazy[node*2+1] += lazy[node];

It's correct?

Reply



Hussein El-Sayed December 25, 2012 at 3:29 AM

Yes you are totally right:).. thanks for correcting me;)..

Reply



Hussein El-Sayed December 25, 2012 at 3:30 AM

The same at line 49 and 50, updated check it now and tell me:)

Reply



Sandipan Manna December 31, 2012 at 12:13 PM

update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5 update_tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12 update_tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100

but your program gives output as 117 !!!

Reply



Hussein El-Sayed January 1, 2013 at 3:06 AM

No it should be 113, however the size of the array needs to be (1+(1<<6)) as it should be $2^{1+\lg N}$..

Also there was some checks needed to be added in the lazy version.. please check it and get back to me.

Reply



Sandipan Manna January 11, 2013 at 11:19 PM

Yes your segment tree size should be int x = (int)(ceil(log2(N)))+1; size = (1 << x); This one!

Reply



InfiniteComplexity February 16, 2013 at 3:27 AM

Wow! Thanks for this post, it was very helpful! However, I'm trying to implement another update_tree_val function that sets the values from a range to one value. e.g. update_tree_val(3,7,4) would set the range [3,7] to the value of 4. How can I do this using lazy propagation on your tree?

Thanks

Reply

Replies



Hussein El-Sayed February 16, 2013 at 11:33 PM

It would be the same but without incrementing..



Aditya Choudhary November 9, 2014 at 8:58 AM

This means in line 48, I have to modify it to tree[node] = lazy[node]; and in line 62, tree[node] = value;?

Reply



PRASHANTHSOUNDAR February 19, 2013 at 9:37 AM

hi, as a beginner to seg tree and lazy propagation i found this post very useful . i tried solving http://www.spoj.com/problems/HORRIBLE/

using slightly modified but same method as this but i am getting WA. can u help?.. http://ideone.com/3rBgSC

Reply



Tilak Raj Singh July 26, 2013 at 6:08 PM

in line 86 it should be tree[node] += (b-a+1)*lazy[node];

Reply

Replies



ravi November 2, 2014 at 8:22 PM

I think you forgot that query returns maximum from given range not the sum of elements of given range

Reply



rl September 15, 2013 at 10:15 AM

In order to do the following:

- 1.) Add x to A[i], A[i+1],...,A[j]
- 2.) Output the sum of A[i],A[i+1],...A[j]

I simply replaced max() with sum() however i am not getting the correct answer.

Reply

Replies



ravi November 2, 2014 at 8:25 PM

also replace line 48 and 86 by tree[node] += (b - a + 1) * lazy[node]; and line 62 by tree[node] += (b - a + 1) * value;

Reply



ইমতিয়াজ November 13, 2013 at 11:37 AM

This comment has been removed by the author.

Reply



Gautam Singh May 24, 2014 at 12:55 AM

in query_tree() method we could have updated the lazy value to the tree before we look for the out of range condition.... it would result in better performance...

am i right about it....please correct me if I am wrong!!

Reply



VIPUL JAIN August 1, 2014 at 7:16 AM

Can you please elaborate how to implement DQUERY?

Reply



Ashish Tilokani September 12, 2014 at 1:36 AM

http://discuss.codechef.com/questions/50866/segment-trees-doubt

Reply



Ashu Pachauri September 25, 2014 at 10:11 PM

Very helpful

Reply



Rahul Kumar December 18, 2014 at 1:23 AM

we have to take array size of 2ⁿ for make segment tree of n element array, then how to make segment tree of 30 or greater elements, because 2^30 = 1073741824 then how to take array of this larger size for make segment tree, how to implement?

Reply



Hussein El-Sayed December 18, 2014 at 1:27 AM

It should be 2^(1+lg N) not 2^N

Reply



ম্যাট্টিক্স.কোড January 6, 2015 at 2:36 PM

This comment has been removed by the author.

Reply



ম্যাট্টিক্স.কোড January 6, 2015 at 2:43 PM

Very good and nice blog..

I am following your code structure in my SG tree implementation:)

Reply



Jayaram Prabhu Durairaj January 7, 2015 at 7:49 PM

This comment has been removed by the author.

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