

A more efficient and effective MFS Identification method

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Abstract—Combinatorial testing (CT) aims to detect the failures which are triggered by the interactions of various factors that can influence the behaviour of the system, such as input parameters, configuration options, and specific events. Many studies in CT focus on designing a elaborate test suite (called covering array) to reveal such failures. When a interaction failure is detected, it is desired to identify the root cause of that failure, i.e., the failure-inducing interactions. Although covering array can assist testers to systematically check each possible factor interactions (without masking effects), however, it provides weak support to locate the failure-inducing interactions. Hence, more auxiliary information are needed to filter out those failure-irrelevant interactions and obtain the failure-inducing ones. Recently some elementary researches are proposed to handle the failure-inducing interaction identification problem. However, some issues, such as unable to identify multiple failure-inducing interactions, and generating too many additional test cases, negatively influence the applicability of these approaches. In this paper, we propose a novel failure-inducing identification approach which aims to handle those issues. The key of our approach is to search for a proper factor interaction at each iteration to check whether it is failure-inducing or not until all the interactions in a failing test cases are checked. To reach this target, we maintain two data structures to guide the selection, i.e., CMINFS(represent for candidate minimal faulty schemas) and CMAXHS(represent for candidate maximal healthy schemas). With these two data set, an optimal interaction can be selected to check so that we can reduce as many additional test cases as possible, as well as not ignore any unchecked-interactions, e.g., the interactions with large number of factors. Our approach will repeat updating the interactions in these two set after a interaction is checked. Moreover, we conduct empirical studies on both widely-used real highly-configurable software systems and synthetic softwares. Results showed that our approach obtained a higher quality at the failure-inducing interaction identification, while just needed a smaller number of additional test cases.

Index Terms—failure-inducing interactions, fault location, debugging aids, combinatorial testing

1 INTRODUCTION

The behavior of modern software are affected by many factors, such as input parameters, configuration options, and specific events. To test such software system is challenging, as in theory we should test all the possible interaction of these factors to ensure the correctness of the System Under Test (SUT)[1]. When the number of factors is large, the interactions that are needed to check increase exponentially. Hence it is, if possible, not practical to conduct such exhaustive testing. Combinatorial testing (CT) is a promising solution to handle the combinatorial explosion problem. Instead of testing all the possible interactions in a system, it focus on checking those interactions with number of involved factors no more than a prior number. Commonly when applying CT,

a elaborate test suite (called covering array) will be generated to cover the valid interaction that needs to be checked. Although covering array is effective and efficient as a test suite, however, it provides weak support to distinguish the failure-inducing interactions from all the interactions.

Consider the following example [2], Table 1 presents a pair-wise covering array for testing an MS-Word application in which we want to examine various pair-wise interactions of options for ‘Highlight’, ‘Status Bar’, ‘Bookmarks’ and ‘Smart tags’. Assume the third test case failed. We can get five pair-wise suspicious interactions that may be responsible for this failure. They are respectively (Highlight: Off, Status Bar: On), (Highlight: Off, Bookmarks: Off), (Highlight: Off, Smart tags: Off), (Status Bar: On, Bookmarks: Off), (Status Bar: On, Smart tags: Off), and (Bookmarks: Off, Smart tags: Off). Without additional information, it is difficult to figure out the specific interactions in this suspicious set that caused the failure. In fact, considering that the interactions consist of other number of factors could also be failure-inducing interactions, e.g., (Highlight: Off) and (Highlight: Off, Status Bar: On, Smart tags: Off), the problem becomes more complicated. Generally, to definitely determine the failure-inducing interactions in a failing test case of n factors, we need to check all the $2^n - 1$ interactions in this test case, which is not possible when n is a large

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TABLE 1
MS word example

id	Highlight	Status bar	Bookmarks	Smart tags	Outcome
1	On	On	On	On	PASS
2	Off	Off	On	On	PASS
3	Off	On	Off	Off	Fail
4	On	Off	Off	On	PASS
5	On	Off	On	Off	PASS

number.

To address this problem, prior work [3] specifically studied the properties of the failure-inducing interactions in SUT, based on which additional test cases were generated to identify them. Other approaches to identify the failure-inducing interactions in SUT include building a tree model [4], adaptively generating additional test cases according to the outcome of the last test case [5], ranking suspicious interactions based on some rules [6], and using graphic-based deduction [7], among others. These approaches can be partitioned into two categories [8] according to how the additional test cases are generated: *adaptive*—additional test cases are chosen based on the outcomes of the executed tests [9], [3], [6], [10], [5], [11], [12], [13] or *nonadaptive*—additional test cases are chosen independently and can be executed parallel [4], [8], [7], [14], [15].

These approaches, however, are essentially approximate solutions to failure-inducing identification (Theoretically, definite solution is of exponential computational complexity). Hence, many issues may affect their effectiveness when applied in practice. Generally, *non-adaptive* approaches usually can accurately identify the failure-inducing interactions, even when there are multiple ones in the SUT. Their effectiveness are based on some mathematical objects [8], [7], [14]. The shortcomings of these approaches are they are very ad-hoc, that is, they are restricted to many limitations, such as the number of failure-inducing interactions must be given as well as the number of the maximal factors involved in a failure-inducing interaction. Moreover, these approaches usually consume many test cases [5]. *Adaptive* approaches are much more flexible, and they mainly focus on one failing test case. Commonly they are needs much less test cases than *non-adaptive* approaches. However, some problems, such as cannot handle multiple failure-inducing interactions (especially they are have overlapping factors), and cannot handle the newly introduced failure-interactions in the additional generated test cases, negatively affect their performance (both precise and recall).

In this paper, we propose a novel failure-inducing interaction identification approach (*adaptive*) which aims to alleviate these issues. Our approach is based on the notions of *faulty schemas* and *healthy schemas* (A schema is equivalent to interaction), among which

the former will result in a failure if a test case contain it while the latter will not. Furthermore, two important data structures are maintained in our approach, i.e., CMINFS(represent for candidate minimal faulty schemas) and CMAXHS(represent for candidate maximal healthy schemas). Through using the two data structures, some schemas in the failing test case can be determined to be healthy, or faulty. Those schemas that cannot be determined by these two data structures will be selected and checked. When an undetermined schema is checked, we will update these two structures. This process is repeated until all the schemas are determined to be healthy or faulty, and the failure-inducing schemas will be selected from those faulty schemas. By doing so, we will not omit any schema that can be candidate failure-inducing interaction in the failing test case. As a result, our approach can handle the cases such as a failing test case containing multiple failure-inducing interactions and interactions consisting of any different number of factors.

The order of selecting un-determined schemas to be checked is important as it affects the update of CMINFS and CMAXHS. Particularly, with the same CMINFS and CMAXHS, choosing different schema to check will result in different number of remaining undetermined schemas, which is strongly correlated to cost of our approach, i.e., the number of additional generated test cases. In this paper, we give a selection strategy which tries to optimise the performance of our approach, i.e., makes the amount of generated additional test cases as small as possible. The selection consists of two steps: (1) We select a sequence of closely related un-determined schemas from the remaining un-determined schemas and sort them according to their relationships with each other. Generally, there is more than one such sequence. As such, we will select the one with the maximum number of schemas for that we can determine the schemas as much as possible in the iteration. (2) For the sequence we select in the first step, we use binary search to find the schema in the sequence as the one to be analyzed each turn, and after a schema is determined, about half of the schemas which are correlated to this schema in the sequence will be determined. At last, the complexity of our approach can reach to $O(d \times t \times \log k + t^d)$, where d is the number of failure-inducing interactions in a failing test case, t is the maximal factors a failure-inducing consists of, and k is number of factors in a failing test case. This complexity is very efficient when compared to other approaches. Besides this, we also provide a variation of our approach which support a feedback mechanism, aiming at handling the newly introduced failure-inducing schemas.

we conducted several empirical studies to evaluate our approach. These experiment objects consists of a group of synthetic softwares with injected faults and some industrial software (small-scale

and medium-scale) with the real faults. Our results generally suggest that our approach obtains a better failure-inducing schema identification results when compared with other existed works, while just generates smaller number of additional generated test cases.

Contributions of this paper:

- We studied the properties of failure-inducing schemas and showed the issues to identify them.
- We proposed a novel adaptive approach which can identify the failure-inducing schemas effectively and efficiently.
- We conducted a series of empirically studies to evaluate the performance of our approach.

The remainder of this paper is organized as follows: Section 2 introduce some preliminary definitions and propositions. Section 4 presents some limitations of exiting failure-inducing interactions identification approaches which motivate this paper. Section 6 describe our approaches for identify failure-inducing schemas. Section 7 give the comparisons in theoretical metrics. Section 8 describe the experiments design on real subjects. Section 9 summarize the related works. Section 10 conclude this paper and discuss the future works.

2 PRELIMINARY

Before we talk about our approach, we will give some formal definitions and proposals first, which is helpful to understand the background of our description of our approach.

Assume that the behaviour of SUT is influenced by k parameters, and each parameter p_i has a_i discrete values from the finite set V_i , i.e., $a_i = |V_i|$ ($i = 1, 2, \dots, k$). In practice, these parameters can represent many factors, such as input variables, run-time options, building options, etc. Next we will give some formal definitions, some of them (Definitions 1, 2, 3) were originally defined in [16].

Definition 1 (Test Case). A test case of SUT is an array of k values, one for each parameter of the SUT, which is denoted as a k -tuple (v_1, v_2, \dots, v_k) , where $v_1 \in V_1, v_2 \in V_2 \dots v_k \in V_k$.

Take Table 1 for example, (*Highlight* = *on*, *Statusbar* = *false*, *Bookmarks* = *false*, *Smarttags* = *on*) is a test case, which is actually a group of values being assigned to each input parameter.

Definition 2 (Schema). For the SUT, the n -tuple $(-, v_{n_1}, \dots, v_{n_k}, -)$ is called a k -degree schema ($0 < k \leq n$) when some k parameters have fixed values and other irrelevant parameters are represented as “-”.

In effect a test case itself is a k -degree schema, when $k = n$. Furthermore, if a test case contains a schema, i.e., every fixed value in the schema is in this test case, we say this test case contains the schema.

Note that the schema is a formal description of the interaction between parameter values we discussed before.

Definition 3 (Subsume). Let c_l be a l -degree schema, c_m be an m -degree schema in SUT and $l < m$. If all the fixed parameter values in c_l are also in c_m , then c_m subsumes c_l . In this case we can also say that c_l is a sub-schema of c_m and c_m is a super-schema of c_l , which can be denoted as $c_l \prec c_m$.

For example, the 2-degree schema $(-, 4, 4, -)$ is a sub-schema of the 3-degree schema $(-, 4, 4, 5)$, that is, $(-, 4, 4, -) \prec (-, 4, 4, 5)$.

Definition 4 (Faulty, MFS). If all test cases that contain a schema, say c , trigger a particular fault, say F , then we call this schema c the faulty schema for F . Additionally, if none of sub-schema of c is the faulty schema for F , we then call the schema c the minimal failure-causing schema (MFS) [3] for F .

Note that MFS is identical to the failure-inducing interaction discussed previously. In this paper, the terms *failure-inducing interactions* and *MFS* are used interchangeably. Figuring the MFS out helps to identify the root cause of a failure and thus facilitate the debugging process.

Definition 5 (Healthy, MHS). A k -value schema, say, c , is called a healthy schema when we find at least one passed valid test case that contains this schema. In addition, if none of super-schema of c is the healthy schema, we then call the schema c the maximal healthy schema (MHS)

MHS is exactly the antithesis of MFS. These two type of schemas, i.e., MFS and MHS, are the keys to our approach, as they are essentially representations of the healthy schemas and faulty schemas in a test case. As will be proven later, other schemas can be determined to be healthy or faulty by these two type of schemas. As a result, we just need to record these two types of schemas (normally a small amount) instead of recording all the schemas in a test case (reach to 2^n , n is the number of factors in a test case) when doing MFS identification.

Definition 6 (Pending). A schema is called pending schema, if it is not yet determined to be healthy schema or faulty schema.

These pending schema is actually the *un-determined* schema discussed in Section 1. In effect, to identify the MFS in a failing test case is to figure out all the pending schemas, and then try to determine each of them to be healthy or faulty schema. After that, the MFS can be selected from those faulty schemas by definition.

When applying CT, the most important work is to determine whether the SUT suffers from the interaction faults or not, i.e., to detect the existence of the MFS. Rather than impractically executing exhaustive

test cases, CT commonly design a relatively small set of test cases to cover all the schemas with the degree no more than a prior fixed number, t . Such a set of test cases is called the *covering array*. If some test cases in the covering array failed in execution, then the interaction faults is regard to be detected. Let us formally define the covering array.

Definition 7. $MCA(N; t, n, (a_1, a_2, \dots, a_n))$ is a t -way covering array in the form of $N \times n$ table, where each row represents a test case and each column represents a parameter. For any t columns, each possible t -degree schema must appear at least once. When $a_1 = a_2 = \dots = a_n = v$, a t -way covering array can be denoted as $CA(N; t, n, v)$.

TABLE 2
A covering array

ID	Test case			
t_1	0	0	0	0
t_2	0	1	1	1
t_3	1	0	1	1
t_4	1	1	0	1
t_5	1	1	1	0

For example, Table 2 shows a 2-way covering array $CA(5; 2, 4, 2)$ for the SUT with 4 boolean parameters. For any two columns, any 2-degree schema is covered. Covering array has proven to be effective in detecting the failures caused by interactions of parameters of the SUT. Many existing algorithms focus on constructing covering arrays such that the number of test cases, i.e., N , can be as small as possible.

3 AN ANALYZE OF EXISTING MFS IDENTIFICATION APPROACHES

In this section, we will give a comprehensive analysis of existing failure-inducing interaction identification approaches, which formally describe the mechanism of each approach, as well as its characteristics and shortcomings.

Let $\mathcal{I}(t)$ denotes all the schemas that are contained in the test case t , e.g., $\mathcal{I}((111)) = \{(1--)(-1--)(--1)(11-)(1-1)(-11)(111)\}$. Additionally, let $\mathcal{I}(T)$ denotes all the schemas that are contained in test set T , i.e., $\mathcal{I}(T) = \bigcup_{t \in T} \mathcal{I}(t)$. By the definition of *faulty schema* (Definition 4), it can easily get that all the faulty schemas in the SUT, i.e., S_{faulty} , can be computed by the following formula

$$S_{faulty} = \mathcal{I}(T_F) - \mathcal{I}(T_P) \quad (\text{EQ1})$$

where T_F indicates all the failing test cases in the SUT, and T_P indicates all the passing test cases (Note that $T_P \cup T_F$ represents all the possible test cases, i.e., T_{all} , in the SUT, and $T_F \cap T_P = \emptyset$).

TABLE 3
Example of the minimal schemas

T_{all}	T_F	T_P	S_{faulty}	MFS
(0, 0, 0)	(0, 0, 0)		(0, 0, -)	(0, 0, -)
(0, 0, 1)	(0, 0, 1)		(0, -, 0)	(0, -, 0)
(0, 1, 0)	(0, 1, 0)		(0, 0, 0)	
(0, 1, 1)		(0, 1, 1)	(0, 0, 1)	
(1, 0, 0)		(1, 0, 0)	(0, 1, 0)	
(1, 0, 1)			(1, 0, 1)	
(1, 1, 0)			(1, 1, 0)	
(1, 1, 1)			(1, 1, 1)	

With EQ1 to compute faulty schemas, the formula to compute the MFS can be described as following:

$$MFS = \{s | s \in S_{faulty} \text{ and } \exists s' \prec s, \text{s.t.}, s' \in S_{faulty}\}. \quad (\text{EQ2})$$

An example of shows the faulty schemas and MFS are in Table 3. Suppose a SUT with three binary parameters, which can be denoted as $SUT(2^3)$. We can find that all the schemas in S_{faulty} are only contained in test set T_F (Those schemas that appear in T_P are filtered out). Additionally, MFS is a minimal schema set which filters those super schemas in S_{faulty} .

In practice, using EQ2 to compute MFS is not possible, as we should execute all the possible test cases of the SUT, which obviously has exponential complexity ($(\prod_{i=1}^k |V_i|)$). Hence, MFS identification approaches are essentially approximation to EQ2, taking into account the tradeoffs between the computing cost and the quality of the MFS identification result. Existing MFS approaches can be classified into two categories [8] according to how the additional test cases are generated: *nonadaptive*—additional test cases are chosen independently and can be executed parallel or *adaptive*—additional test cases are chosen based on the outcomes of the executed tests.

3.1 Non-adaptive approaches

The simplest idea is to sample some test cases instead of all the possible test cases for MFS identification. Covering array is a good sampling mechanism candidate, as it can systematically cover all the schemas with degrees no more than a prior fixed number t , as a result, those MFS with degrees no more than t can be guaranteed to be detected. With covering array, the formula to compute MFS will be revised. Formally, let T_f denotes those failing test cases in a covering array, and T_p denotes those passing test cases. Obviously, $T_f \subseteq T_F$ and $T_p \subseteq T_P$. Furthermore, $\mathcal{I}(T_f) \cup \mathcal{I}(T_p) \supseteq S_{\hat{f}}$, where $S_{\hat{f}}$ represents all the schemas with degrees no more than t . Then the formula to describe the MFS identification using covering array can be shown as EQ3.

$$S_{faulty} = \mathcal{I}(T_f) - \mathcal{I}(T_p)$$

$$MFS = \{s | s \in S_{faulty} \text{ and } \exists s' \prec s, \text{s.t.}, s' \in S_{faulty}\} \quad (\text{EQ3})$$

TABLE 4
MFS identification using covering array

T_{all}	T_f	T_p	S_{faulty}	MFS
(0, 0, 0)	(0, 0, 0)		(0, 0, -)	(0, 0, -)
(0, 0, 1)			(0, -, 0)	(0, -, 0)
(0, 1, 0)			(-, 0, 0)	(-, 0, 0)
(0, 1, 1)		(0, 1, 1)	(0, 0, 0)	
(1, 0, 0)				
(1, 0, 1)		(1, 0, 1)		
(1, 1, 0)		(1, 1, 0)		
(1, 1, 1)				

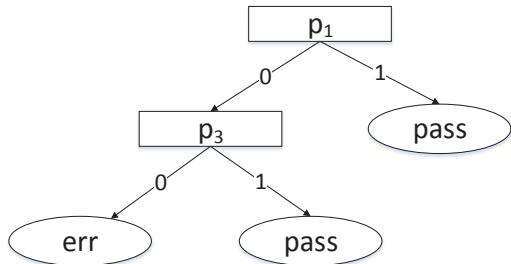


Fig. 1. CTA example

As $\mathcal{I}(T_f) - \mathcal{I}(T_p)$ is normally not equal to $\mathcal{I}(T_F) - \mathcal{I}(T_P)$, EQ3 may not get an accurate result. For example, Table 4 shows the result which directly using covering array to compute MFS. This example using the same SUT as Table 3, i.e., SUT(2^3) with MFS (0, -, 0) and (0, 0, -). In this example, a 2-way covering array consists of one failing test case (0, 0, 0) and three passing test cases (0, 1, 1), (1, 0, 1) and (1, 1, 0). Using EQ3, we can finally obtain the MFS (0, 0, -), (0, -, 0) and (-, 0, 0), which is different from the correct result in Table 3. Besides this, the complexity of EQ3 is exponential. This is because we should list all the possible schemas in each test case of T_f , which is $|T_f| * 2^{k-1}$. As directly utilizing covering array cannot give an ideal result, more efforts are desired to improve this result.

3.1.1 Original Covering array

Yilmaz [4] proposed an MFS identification approach, which can quickly characterize the MFS based on a covering array. It uses classification tree analysis (CTA) on the executed test cases of covering array to derive a tree structure. The conjunction from the root to a leaf node represents a MFS. For example, when using CTA (Implemented by Weka[17]), a tree will be built like Fig 1.

In this figure, only $p_1 = 0$ and $p_3 = 0$ will result in a failure, indicating that (0, -, 0) is MFS. Note that the way that CTA constructing decision trees usually work top-down, by choosing a parameter at each step that best splits the set of test cases [18]. Consequently, after choosing the first parameter p_1 (three parameters have equal chances at the first step), only one of p_3 and p_2 can be selected. Hence CTA will identify only one of (0, -, 0) and (0, 0, -), while ignore another.

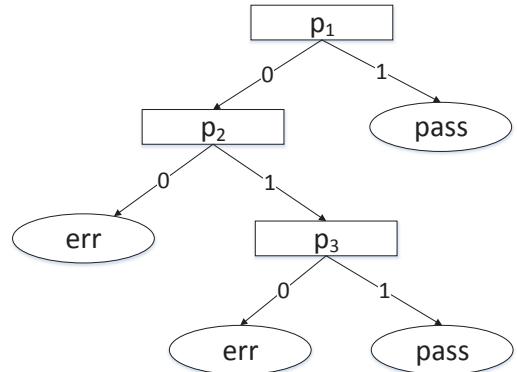


Fig. 2. Another CTA example

Another point that needs to note is that only identify the MFS with tree structure. Hence, even if CTA find two MFS in this SUT, these MFS can only be identified as Figure 2, which means that (0, 0, -) and (0, 1, 0) which is still not correct.

Another approach that are directly based on covering array is proposed by Zhang [15]. proposed approaches that are based on covering array, but with different techniques. The first one is using classification tree techniques, and the second one using sat solver.

3.1.2 Higher covering array

3.2 Adaptive approaches

4 MOTIVATION OF THE APPROACH

Take OFOT for example, then illustrate the problems, and other similar approaches problems.

4.1 Multiple MFS

4.1.1 Overlapping MFS

4.2 High degree MFS

4.3 Too many test cases

4.4 Needs too many computing resources

4.5 Introducing of newly MFS

A big table that consists of many

5 THEORY FOUNDATION

Assumption 1.

Proposition 1. All the schemas in a passed test configuration are healthy schemas. All the subschemas of a healthy schemas are healthy schemas.

Proposition 2. If schema s_a subsumes schema s_b , schema s_b subsumes schema s_c , then s_a subsumes s_c .

Proposition 3. All the parent-schemas of a faulty schema are faulty schemas.

Proposition 4. All the subschemas of a healthy schema are healthy schemas.

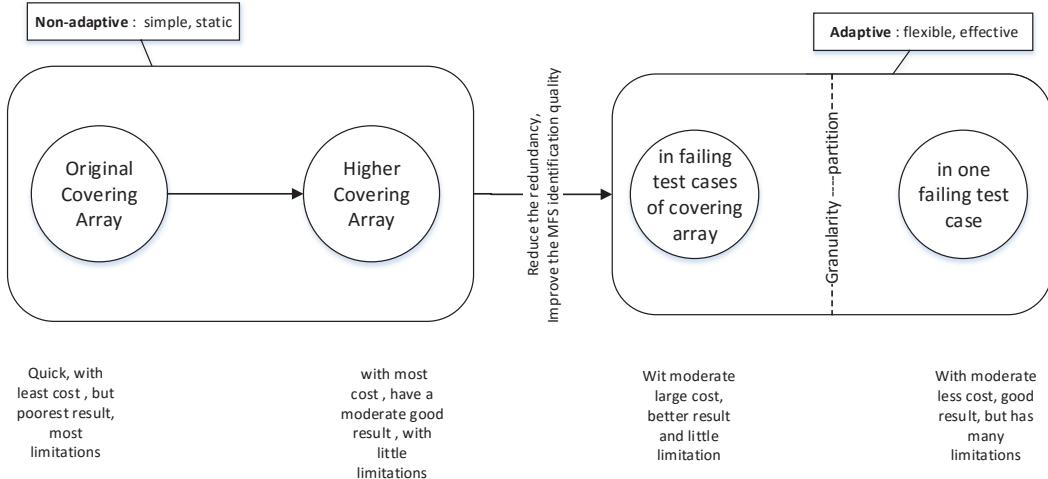


Fig. 3. A simple conclusion of the relationships between these MFS identification approaches

5.1 weaken of the assumptions

6 THE APPROACH TO IDENTIFY THE MINIMAL FAILURE-INDUCING SCHEMAS

Based on these definitions and propositions, we will describe our approach to identify the failure-inducing schemas in the SUT. To give a better description, we will start give an approach with an assumption , and later we will weak the assumption.

Definition 8 (candidate minimal faulty schema). *If a schema is a faulty schema and satisfy the followed condition: 1.none of its subschemas are faulty schemas, 2. at least one subschema is pending schema.(is need discuss!!!!!! one or none is okay?)We then call the schema a candidate minimal faulty schema(CMINFS for short).*

Definition 9 (candidate maximum healthy schema). *If a schema is a healthy schema and and satisfy the followed condition: 1.none of its parent-schemas are healthy schemas, 2. at least one subschema is pending schema. We then call the schema a candidate maximum healthy schema(CMAXHS for short).*

6.1 additional failure-inducing not be introduced

Assumption 2. *Any newly generated test configuration will not introduce additional failure-inducing schemas.*

There are similar assumptions defined in[], however, it is a strong assumption, which we will change later. And still for this assumption, we can get the followed lemma which can help us to identify whether a pending schema is a healthy schema or a faulty schema.

Lemma 1. *For a pending schema, we generate an extra test configuration that contains this schema. If the extra test configuration passes, then this schema is a healthy schema. If the extra test configuration fails, then the schema is a faulty schema.*

Proof: According to definition of healthy schema, it is obvious that this schema is a healthy schema when the extra test configuration passes.

When the extra configuration fails, this is a faulty schema (or there exists no faulty schema and this test configuration would not fail because the assumption says that this newly generated configuration will not introduce additional faulty schemas). \square

6.2 framework to identify the minimal faulty schema

As we can identify a pending schema to be healthy schema or faulty schema by generating newly test configurations, then the approach to identify the minimal faulty schema is clear. Fig.4 shows an overview of our approach. We next discuss each part of the approach in more detail.

6.2.1 Choosing a pending schema

Before we think getting a pending schema, we should assume an general scenario. That is, we have already make sure some schemas to be healthy schemas and some schemas to be faulty schemas, then assume that we still don't meet the stopping criteria, we will choose a pending to check next. But first, we should make sure what schema is pending schema?

In effect, the schemas we can't make sure whether are faulty schemas nor healthy schemas are our wanted. Step further, we sperate this condition into two parts: 1. can't make sure it is faulty schema. As we already know some faulty schemas, then we just make the schema that first not be any one of these faulty schemas and not be the parent-schema of any one of these faulty schemas(As if not, it must be a faulty schema according to the proposition 3). 2 can't make sure it is healthy schema. Similar to the first condition,we already know some healthy schemas, then we just make the schema that first not be any one

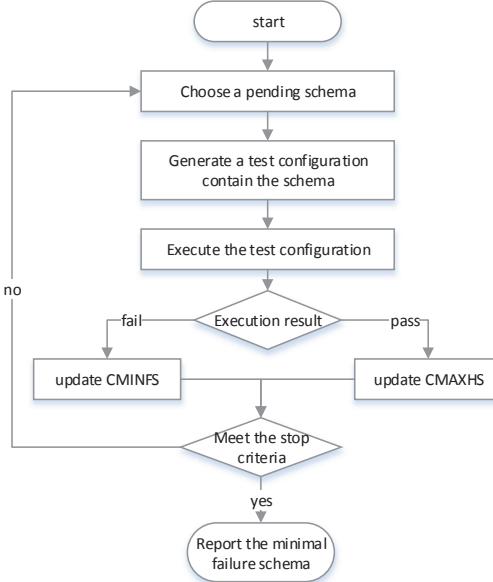


Fig. 4. Overview of approach of identifying MFS

of these healthy schemas and not be the subschema of any one of these healthy schemas. It is obvious a pending schema must meet both the two conditions.

However, this is not the end of story of checking a pending schema. With the process of identifying more and more faulty schemas and healthy schemas will be identified. It is not a small number, as it can reach to $O(2^n)$. So both for space and time restriction, we should not record all the faulty schemas and healthy schemas. Then we should find another way to check the pending schema.

To eliminate this problem, we propose an method which can check a pending schema with a small cost. It need to record the CMINFS and CMAXHS all through the process. Instead record all the faulty schemas and healthy schemas, CMINFS and CMAXHS are rare in amount, which can help to largely reduce the space need to record. Then according to the followed two propositions, we can easily check a pending schema.

Proposition 5. *If a schema is neither one of nor the parent-schema of any one of the CMINFS, then we can't make sure whether it is a faulty schema.*

Proof: Take a schema s_a , a CMINFS set S_{cminks} and a faulty schema S_{fs} set which determined now. It is note that any element $fs_i \in S_{fs}$ must meet that either $fs_i \in S_{cminks}$ or fs_i be the parent-schema of one of the S_{cminks} . Assume that s_a is neither the one of nor the parent-schema of any one of the S_{cminks} . To proof the proposition, we just need to proof that s_a is neither the one of nor the parent-schema of any one of the S_{fs} .

As s_a is neither the one of nor the parent-schema of any one of the S_{cminks} , so it is not one of S_{fs} . Then we assume s_a is the parent-schema of one of S_{fs} , say,

fs_j . As fs_j must meet either $fs_j \in S_{cminks}$ or fs_j be the parent-schema of one of the S_{cminks} . So s_a is the parent-schema of one of S_{cminks} according to the proposition 2, which is contradict. So the proposition is correct. \square

Proposition 6. *If a schema is neither one nor the the subschema of any one of the CMAXHS, then we can't make sure whether it is a healthy schema.*

We ignore this proof as it is very similar to the previous one.

Up to now, we can judge whether a schema is a pending schema, but in effect, there are many pending schemas in a test configurations, especially at the beginning of our process. To choose which one has an impact on our approach. To better illustrate this problem, we consider the followed example:

Assume the failing test configuration: (1 2 1 1 1 2 1 2), that the CMINFS set: (1 2 1 1 - 2 - -) (- 2 - - 1 2 - 2), the CMAXHS set:(- 2 - - - - -) (- - - 1 - - -). We list some pending schemas followed (not all, as the number of all the pending schemas is too much that is not suitable to list here):

(1 2 1 - - 2 1 2) (- 2 1 1 - 2 - 2) (1 2 1 1 - - - -) (- 2 1 1 - 2 - -) (1 2 1 - - - - -) (- 2 - 1 - - 2 -) (1 2 - - - - -) (- - 1 1 - - - -) (1 - - - - - -) (- - 1 - - - - -).

Choose what is really different. If we choose (1 2 1 - - - -), assume we check it as a healthy schema. Then all its subschemas, such as (1 2 - - - - -) (1 - - - - - -) (- - 1 - - - - -), are healthy schemas. It means that we did not need generate newly test configurations for them. But if we check it as a faulty schema, we can make sure all its parent-schemas are faulty schemas, in this case, they are (1 2 1 1 - - - -) and (1 2 1 - - 2 1 2) need no newly test configurations to test.

Let's look at this problem from another angle. Take the schema as a integer number, these parent-schemas of a schema is like the integer numbers bigger then this number, and these subschemas of a schemas is like the numbers smaller than this number. Take a float number as the metric, then, we can describe the faulty schema as the number bigger than this metric, and healthy schema as number smaller than this metric. So the minimal faulty schema is just the number most approximate the metric and bigger than the metric.

It seems like a search problem scenario. Then can we directly apply the efficiently algorithm binary-search? the answer is no, because there are schemas that neither parent-schema or subschema relationship, such as(1 2 1 - - 2 1 2) (- 2 1 1 - 2 - 2). So to utilize the binary search technique, we should make some change. First, should give the followed definitions:

Definition 10 (chain). *A chain is an ordered sequences of schemas in which every schemas is the direct parent-schema of the schema that follows. Moreover,if all the schemas in a chain are pending schemas, we call the chain a pending*

chain.

This definition is similar to the path in []

As all the schemas in a chain are have relationships, then we can apply binary search technique. As we all know, the longer the pending chain, the better performance binary search technique will get. So we should choose a pending chain as longer as possible each iteration.

To get a longest chain, we need to ensure that the head schema of this chain do not have any parent-schema which is a pending schema(called up pending schema), and the tail schema do not have any subschema which is a pending schema (called down pending schema). The algorithm that get the up pending schemas and down pending schemas are list in algorithm 1 and algorithm 2.

As showed in algorithm 1, we start from the failing test configuration \mathcal{T} , assign it to the $rootSchema$ (line 1)and then add it to a $lists$ (line 3). We then do some operation (line 4 -line 12)to this lists and at last get these pending schemas in lists as up pending schemas.(line 13) This operation consists of two iteration:

1. successively get one CMINFS in S_{CMINFS} . (line 4). Define a temple value $nextLists$ which initialize an empty set(line 5).We then execute the second iteration. After that, we will eliminate these same schemas list in $nextLists$ (line 10)and then assign to $lists$ (line 11).

2. Successively get one schema in $lists$ (line 6). And then mutant it according to the CMINFS to a set of schemas(line 7). Add them to the $nextLists$ (line 8).

The mutant procedure for a schema is just remove one value in it which this value is also in CMINFS. This procedure will result in k mutant schemas if the CMINFS is a k -value schema. By doing this, any mutant schema will not be the parent-schema of the corresponding CMINFS.

After this two iteration, the schemas in the $lists$ will not be the parent-schema of any CMINFS in the S_{CMINFS} .

Fig.5 gives an example to the algorithm 1.

Algorithm 2 is a bit different from algorithm 1. As our target is to get the minimal value pending schema, we get started from a 0-value schema(line 1). Then to make schemas not to be the subschema of any CMAXHS in S_{CMAXHS} , we should let the schemas must contain at least one value that is not in the corresponding CMAXHS. So the strategy is to get a reverse schema of the CMAXHS which this schema consists of all these values in the failed configuration except these are in CMAXHS(line 5). And mutant the schema by adding one value in the reversed schema to make the schma not to be the subschema of the correspindng CMAXHS(line 8). After two iteration similar to Algorithm 1, we will get all the schemas that meet that not to be subschema of any CMAXHS in S_{CMAXHS} from which we choose these pending

Algorithm 1 getting up pending schema

Input: \mathcal{T} ▷ failing test configuration
 S_{CMINFS} ▷ set of CMINFS
 S_{CMAXHS} ▷ set of CMAXHS

Output: S_{UPS} ▷ the set of up pending schemas
▷ %comment: initialize%

```

1:  $rootSchema \leftarrow \mathcal{T}$ 
2:  $lists \leftarrow \emptyset$ 
3:  $lists \cup \{rootSchema\}$ 
4: for each CMINFS in  $S_{CMINFS}$  do
5:    $nextLists \leftarrow \emptyset$ 
6:   for each schema in  $lists$  do
7:      $S_{candidate} \leftarrow mutant_r(schema, CMINFS)$ 
8:      $nextLists \leftarrow nextLists \cup S_{candidate}$ 
9:   end for
10:   $compress(nextLists)$ 
11:   $lists \leftarrow nextLists$ 
12: end for
13:  $S_{UPS} \leftarrow \{s | s \in lists \wedge s \text{ is pending}\}$ 

```

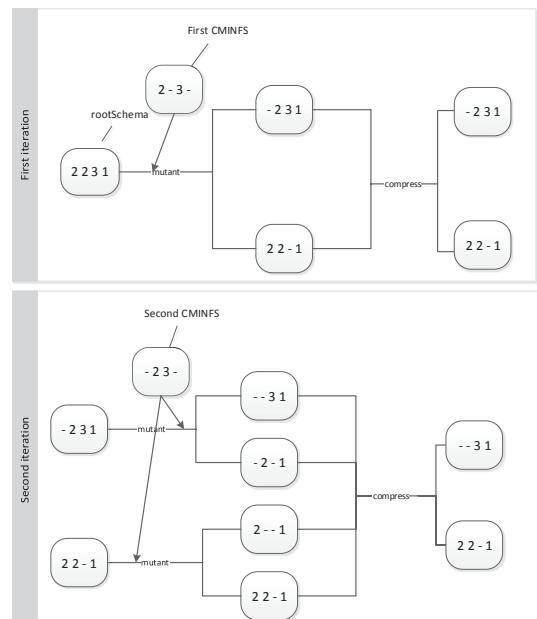


Fig. 5. example of getting up pending schemas

schemas as down pending schemas(line 14).

Fig.6 gives an example to the algorithm 2. we can see.

After we can get the up pending schemas and down pending schemas, then the algorithm of getting the longest chain can be very simple,which is list in Algorithm 3. In this algorithm we can see that we just search through the up pending schemas and down pending schemas(line 7 - 8) to find the two schemas which has the maximum distance(line 9 - 13). The distance of a k-value schema and a l-value schema ($k \neq l$) is defined as:

Algorithm 2 getting down pending schema

Input: \mathcal{T} ▷ failing test configuration
 S_{CMINFS} ▷ set of CMINFS
 S_{CMAHHS} ▷ set of CMAXHS

Output: S_{DOWNS} ▷ the set of down pending schemas
▷ %comment: initialize%

```

1: initschema  $\leftarrow ()$ 
2: lists  $\leftarrow \emptyset$ 
3: lists  $\cup \{initschema\}$ 
4: for each  $CMAHHS$  in  $S_{CMAHHS}$  do
5:    $reverse \leftarrow reverse(CMAHHS)$ 
6:    $nextLists \leftarrow \emptyset$ 
7:   for each schema in lists do
8:      $S_{candidate} \leftarrow mutant_a(schema, reverse)$ 
9:      $nextLists \leftarrow nextLists \cup S_{candidate}$ 
10:  end for
11:  compress(nextLists)
12:  lists  $\leftarrow nextLists$ 
13: end for
14:  $S_{DOWNS} \leftarrow \{s | s \in lists \wedge s \text{ is pending}\}$ 

```

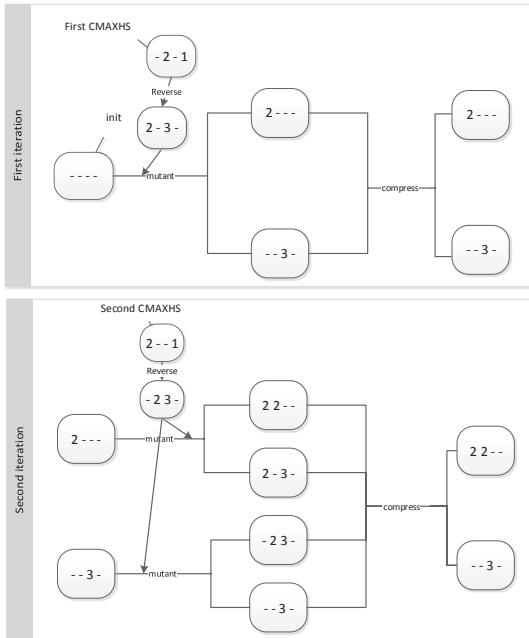


Fig. 6. example of getting down pending schemas

$$distance(S_k, S_l) = \begin{cases} -1, & S_k \text{ is not parent-schema of } S_l \\ k - l, & \text{otherwise} \end{cases}$$

Then we just use *makechain* procedure to generate the longest chain. The *makechain* procedure is very simple, it just repeat adding one schema by keeping all the value in down pending schema and removing one factor of the previous schema.

The last step of getting the schema is just choose the schema from the longest chain. To clearly describe

Algorithm 3 Finding the longest pending schema

Input: \mathcal{T} ▷ failing test configuration
 S_{CMINFS} ▷ set of CMINFS
 S_{CMAHHS} ▷ set of CMAXHS

Output: $CHAIN$ ▷ the chain

```

1:  $S_{UPS} \leftarrow getUPS(\mathcal{T}, S_{CMINFS}, S_{CMAHHS})$ 
2:  $S_{DOWNS} \leftarrow getDOWNS(\mathcal{T}, S_{CMINFS}, S_{CMAHHS})$ 
3:  $max \leftarrow 0$ 
4:  $head \leftarrow NULL$ 
5:  $tail \leftarrow NULL$ 
6:  $chain \leftarrow NULL$ 
7: for each up in  $S_{UPS}$  do
8:   for each down in  $S_{DOWNS}$  do
9:     if  $distance(up, down) \geq max$  then
10:       $max \leftarrow distance(up, down)$ 
11:       $head \leftarrow up$ 
12:       $tail \leftarrow down$ 
13:    end if
14:  end for
15: end for
16:  $CHAIN \leftarrow makechain(head, tail)$ 

```

the approach, we discuss it in the overall identifying algorithm which is list in the Algorithm 4. As we have talked previously, we first judge the end criteria, if meet we will report the minimal faulty result. Otherwise we will do the loops. In the loop, we first judge if it is the beginning or headIndex greater than tailIndex, then we will update the CMINFS and CMAXHS, followed generated the longest chain and initial the headIndex, tailIndex and middleIndex. If not we will get the middleIndex indicate the half. Then will select the schema with index of middleindex in the longest chain. And then generate a test configuration and execute the SUT under it. If the result passed, we let tail = middle - 1 else head = middle + 1. By doing this and the previous set the middleIndex we can apply the binary search to the indenting. It is noted that we intitally let middle = 0. which we want know if this chain has faulty schema as soon as possible.

6.2.2 generate a new test configuration

To generate a new test configuration to test the schema. The generated test configuration must meet the followed rules:

1. S_i must contain the selected schema.
2. must not
3. constraint(system-wide constraint and test case specific constraint, i.e., masking effect)

The first one is easy to meet. We just keep the same value which are in the schema are also in the test configuration. We must meet the second condition for that if we contain another one, combine this then this test configuration will contain an parent-schema of this selected schema in the original test configuration,

Algorithm 4 identify process

Input: \mathcal{T} ▷ failing test configuration
 \mathcal{S}_{CMINFS} ▷ set of CMINFS
 \mathcal{S}_{CMAHHS} ▷ set of CMAXHS

```

1: while hasn't meet the end criteria do
2:   if the beginning or  $headIndex > tailIndex$  then
3:      $update(\mathcal{S}_{CMINFS}, \mathcal{S}_{CMAHHS})$ 
4:      $longest \leftarrow getLongest(\mathcal{T}, \mathcal{S}_{CMINFS}, \mathcal{S}_{CMAHHS})$ 
5:      $headIndex \leftarrow 0$ 
6:      $tailIndex \leftarrow length(longest) - 1$ 
7:      $middleIndex \leftarrow 0$ 
8:   else
9:      $middleIndex \leftarrow \frac{1}{2} \times (tailIndex + headIndex)$ 
10:  end if
11:   $SCHEMA \leftarrow longest[middleIndex]$ 
12:  generate a extra test configuration  $\mathcal{T}'$  contain  $SCHEMA$ 
13:  execute SUT under  $\mathcal{T}'$ 
14:  if the test configuration passed then
15:     $tailIndex \leftarrow middleIndex - 1$ 
16:  else
17:     $headIndex \leftarrow middleIndex + 1$ 
18:  end if
19: end while
20: report the minimal faulty schemas

```

and it will confuse us whether it is indicate this schema or its parent-schemas dedicate this result. To fulfil this condition, we need to choose other available values in the SUT which are different from the original test configuration. the third one is that we should consider the constraints

6.2.3 execute SUT under the test configuration

In real software testing scenario, when we test a SUT under a test configuration, there may be many possible testing state: such as pass the testing assertion, don't pass the testing the assertion but with different failure type, can't complete the testing. To get a clear discussion, in this paper we just use *pass* represent the state that pass the testing assertion and *fail* represent all the remained state.

6.2.4 update information

The update information is followed when the current chain is checked over. Then before we generate another longest chain, we should update the CMINFS and CMAXHS set. In fact, we just need the CMINFS and CMAXHS in the longest chain.

6.2.5 stop criteria

The stop criteria is clear, our algorithm stops when there are no pending schemas left, for that when we once can checked all the schemas of a test configuration, we can get the minimal faulty schemas, which is the target of our algorithm. And whether there are

pending schemas can be easily checked by that if we can't generate a longest chain (the length must greater than 1), there must be no pending schemas.

6.2.6 report the result

In fact, the last in the CMINFS lists is must be the minimal faulty schemas. For that if these schema in the CMINFS is not minimal faulty schema, then there must be some pending schemas. However, the algorithm stop when there are no pending schemas remained. So at last these in the CMINFS must be the minimal faulty schemas.

6.2.7 new ideas

you should first identify a failure-inducing schema, and then find others. initial the find tuples, and then using the same algorithms as others. when identify is confirm , then choosing new longest path.

6.3 example

We will give an complete example listed in Table 5. Assume that a SUT is . constraints.

6.4 Without Safe Values Assumption

Up to now our algorithm is based on the assumption that additional generated test configuration does not introduce newly faulty schema. We will give an augment algorithm this section to eliminate this assumption. Our augment algorithm is inspired by the feedback machinery in the controlling system, which in high level we will validate the identify result at the end of the aforementioned algorithm, and if the schema is validated as a minimal faulty schema, we will end the algorithm, otherwise we will repeat the algorithm again to adjust the result. The detail of our augment algorithm is list in Algorithm 5. We can find the first part of augment algorithm is an unlimited loop, in the loop we first use previous identify algorithm to identify the MFS, and then we will check the result, the check process is just to addtional generate another test configurations and execute. When we find the MFS is not right, which means our process introduce newly MFS, then we first label this MFS as an healthy schema, and then update the CMAXHS. After check, if we find at least one MFS is not right, we then empty the CMINFS, and then readd these right identified MFS to this set, and then reprocess this process untill all the MFS is validated as right. The second part of our algorithm is just we look through the generated test cases one by one, if it failed, and did not contain the MFS we identified in the first part, which means it introduce newly faulty schema, and we will rerun this process to find these introduced MFS.

Table 6 illustrate this augment algorithm with an example.

TABLE 5
An example of identifying

S_{CMINFS}	S_{CMAHHS}	Longest Chain	Choosing Schema	Generating Test Configuration	Execution result
1	1	1	1	1	false
1	1	1	1	1	pass
1	1	1	1	1	false

TABLE 6
An example of augment identifying

S_{CMINFS}	S_{CMAHHS}	Longest Chain	Choosing Schema	Generating Test Configuration	Execution result
1	1	1	1	1	false
1	1	1	1	1	pass
1	1	1	1	1	false

Algorithm 5 augment identify process

```

Input:  $\mathcal{T}$                                 ▷ failing test configuration
         $S_{CMINFS}$                           ▷ set of CMINFS
         $S_{CMAHHS}$                           ▷ set of CMAXHS
1: while true do
2:    $S_{MFS} = Identify\_process(\mathcal{T})$ 
3:   for each  $MFS$  in  $S_{MFS}$  do
4:     if validate( $MFS$ ) = fail then
5:       updateHealthySchemas( $S_{CMAHHS}, MFS$ )
6:     end if
7:   end for
8:   if at least one  $MFS$  is not correct then
9:     empty( $S_{CMINFS}$ )
10:    add all the validated  $MFS$  in the  $S_{CMINFS}$ 
11:   else
12:     break
13:   end if
14: end while
15: for each testCase in EXTRATESTCASES do
16:   if testCase introduce newly  $MFS$  then
17:     augment_identify_process(testCase)
18:   end if
19: end for
20: report the minimal faulty schemas

```

7 EVALUATION WITH SIMULATED SUT

In this section, we designed a simulated SUT of which the number of parameters and the value of each parameter can both be customized. In addition, we can also manually inject faulty schemas in the simulated SUT. We will conduct a series of experiments with this simulated SUT in this section. The goal of our experiments is to evaluate the efficiency and effectiveness of our approach compared with other existed techniques. The main reason why we use simulated program is that we can easily run a bench of experiments with various states of a SUT, i.e., different parameters and different MFSs. By doing so we can thoroughly learn the performance of each algorithm

without biases.

As we discussed in the background section, the fault privilege properties just effect the way we determine a configuration is fail or pass. It doesn't effect the performance of each algorithm. So to be simple and clear, we omit the fault privilege in the simulated SUT, i.e., all the fault in the SUT has the same privilege. This is a ideal scenario which may not exist in real softwares, in the empirical studies of section 7 we will deal with the real faults in some open-source softwares with different privileges.

7.1 comparison algorithms

There are several algorithms aim to identify the MFS, they can be classified as non-adaptive methods and adaptive methods. The first set of methods do not need additional test configurations and can identify the MFSs when given an executed test configurations while the second one will generate some more additional test configurations to facilitate the process of identifying MFSs in a failing test configuration. Our algorithm is belong to the second part. To make the comparison clear and fair, we just choose the algorithms which all belong to the second part as the comparison subject, i.e., adaptive methods. These algorithms is listed as follows:

(1)OFOT, proposed by Nie[], it change one factor of a test configuration one time to generate a newly configuration and then analysis the MFSs according to the executed result of these generated configurations. (2)FIC-BS, proposed by Zhang, using a delta debugging strategy to isolate the MFSs (3) LG, proposed by Martine[], use a locatable graph to represent the faulty schema and adopt divide and conquer strategy to find the MFSs. (4) SPI[], proposed by Ghandehari, they generate configurations for these mostly suspicious schemas and give a rank of these schemas at last based on the suspicious "degree" of them (5) CTA, proposed by Shakya[], it combine the OFOT strategy and the classified tree technique which first applied in characterize the MFS in [] to analysis the MFSs. (6)RI,

proposed by Li[], it is also using delta debugging, but the factors needed to change in a iteration are different from FIC-BS at some conditions.(7)IterAIFL, proposed by Wang[], a mutant based on OFOT, which changes different number of factors instead of one each iteration for a test configuration. (8)TRT and TRT-NA, two approaches proposed in our previous work[], using a structured tree model to analysis the MFSs in a failing configurations. More details of these algorithms will be discussed in the related works.

Some parameters of these algorithm are assigned: for SP, the degree of MFSs we set is 2(except for the last group of experiment, which is 4 instead), for CTA,the confidence factor of classified tree algorithm is 0.25 as same as [].

7.2 evaluation metrics

There are three facts we care about each algorithm: 1)How many additional test configurations do an approach need to generate? 2)How many MFSs can an approach identify for a SUT? 3)Does all the schema identified by an approach are correctly? We will define two formula to represent the second and third facts respectively, i.e., recall and precise:

$$\text{recall} = \frac{\text{correctly Identified MFS}}{\text{all the MFS in SUT}}$$

$$\text{precise} = \frac{\text{correctly Identified MFS}}{\text{all the identified MFS}}$$

The recall is percentage of the correctly identified MFSs of an algorithm of all the MFSs in the SUT. The bigger the metric the better an approach can perform on finding MFSs in a SUT. There are some factors may influence this metric of an algorithm, mainly are whether an approach consider a failing configurations contain multiple MFSs, whether consider if there exist overlapped MFSs in a failing configurations and whether consider the newly introduced MFSs of the additional generate configurations. The precise is the percentage of the number of correctly identified MFSs of algorithm of all the number of MFSs this algorithm identified. This metric measure the accurately of an approach and the bigger it is the better the approach performs. This metric is related to whether a algorithm consider the influence if an generated test configuration contain newly MFSs. If the approach did not consider so, it may make a wrong judgement to determine some schemas in the original configuration to be faulty schemas which in fact are not.

Of course these factors we mentioned are just affect these the no-machine learning algorithm, as for the machine learning algorithms,i.e., CTA, the factor that matters is the classified algorithms itself.

7.3 experiment setups

We will carry out the followed five experiment in this section:

1)For the first one, we will give a set of SUTs with 8 parameters, each parameter has 3 values, all of them only have one MFS, the difference between each SUT in this set is the MFS we inject in it. All the MFSs has the degree of 2. There are $\binom{8}{2} = 28$ possible MFSs we can inject, so the number of set of SUTs is 28. For each SUT in this set, we will feed each algorithm the same failing configurations contain the MFS we inject, and then let these algorithm identify the MFS in it. We will record the additional configurations each algorithm needed. As the first experiment is simple, so all the approaches can identify the MFS with precise 1 and recall 1. This experiment will give us a initial view of the cost of each algorithm on identifying the MFS in a failing configuration.

2)The second experiment is similar to the first one, except that we will inject double different MFSs with degree 2 in each SUT. It is easily compute their are $\binom{28}{2} = 278$ possible SUTs in this experiment. For each algorithm, we also feed them with the same failing configuration contain the two MFSs and then let them identify them. Another difference from the experiment 1 is that in this experiment we will record the "recall" and "precise" of each algorithm, as these metric is not equal to 1 for all the algorithms like experiment 1. This experiment mainly focus on observing the performance of each algorithms in facing multiple MFSs in a configuration.

3)The third experiment aims at learning the capability that each algorithm handling the introducing newly MFSs. To accomplish this goal, We firstly inject two MFSs in a SUT, and then feed each algorithm a failing configuration which contain only one MFS, remaining another MFS not in it. To increase the possibility of introducing newly MFSs for each algorithm, we set the second MFS with degree 1, which may be easily introduced by just changing one factor of the original failing configuration. The first MFS in the original configuration has degree 2. So there are $\binom{8}{2} \times \binom{8}{1} = 224$ possible SUTs in this experiment. The same as experiment 2, we will record the "recall" and "precise" of each algorithm.

4)For the fourth experiment, we will vary the number of parameters of SUT. In specific we will take 8, 9, 10, 12, 15, 20, 30, 40, 60, 80, 100, 120 respectively as the number of parameters of SUT. For each number of parameters, say n, we will generate $\binom{n}{2}$ SUTs with different MFSs injected. And for these SUT, we will do the same process as experiment 1. At last we will record the average number of additional configurations each algorithm needed for each number of parameters. The goral of this experiment is to see the influence of the number of parameters on each approach.

5)In the fifth experiment we will fix the number of parameters of SUT to be 8, but vary the number of degrees of the MFS we inject in the SUT. That is, we will inject MFS in SUT with degrees 1,2,3,4,5,6,7,8 respectively, for a degree of m ($m = 1,2,3,4,5,6,7,8$), we will generate $\binom{8}{m}$ SUTs with different m-degree MFS injected. And then we will do the same thing as experiment 4 , i.e., record the average number of additional configurations each algorithm needed for each number of degree. This experiment is order to inspect whether these algorithm can handle different degree of MFS in a failing configuration.

7.4 results and discuss

We will discuss the results into five groups according to the experiment set up.

1)The result of the experiment 1 is depicted in figure 7. The y-axis report the number of configurations and the x-axis represent a SUT. Algorithms differs from each other in the signal of point and the line linking these points. For a particular algorithm, a point indicate the number of configurations corresponding to the y-axis needed for identifying the MFS in the SUT which corresponding to the x-axis of this point.

There are some intuitive information we can learn from this figure. First, the configurations needed of some algorithms didn't change along with the change of SUT. These algorithms are AIFL, CTA, OFOT,LG, while others need different configurations according to the SUT they are facing. Second, some algorithms(sp, aifl, trtNA, OFOT, CTA) shows a obvious disparity in this figure, we can easily rank them according to the number of configurations, i.e., ofot = cta < trtNA < sp < aifl. Besides, These algorithms needs more configurations than the remaining algorithms.

For the remaining algorithms, their results are tangled with each other. So in order to rank them, we give an average configurations each algorithms needed : Chain 9.25, FeedBack 11.25, fic 7.96, ri 8 ofot 16 lg 11 , sp .29, aifl 36, trt 10.75 trt-NA 23.75 CTA 16. Overall, we can conclude an initial rank in the number of configurations each algorithm needed as: *FIC-BS < RI < LG < TRT < CHAIN < OFOT = CTA < TRT - NA < AIFL < SP*.

2)Figure 8 list the result of experiment 2. This figure doesn't show the number of configurations of each algorithms,instead it just show the average precise and recall for each algorithms as not all these algorithms can both get recall 1 and precise 1 as experiment 1. It is obvious meaningless if we compare two algorithms that one algorithm can reach recall 1 and precise 1 while the other can't. So we will only list the average number configurations of these algorithms that can reach recall 1 and precise 1 later.

From this figure, we can find the following algorithms: chain, lg, sp, aifl, trt and trt-na can both get the precise and recall 1, which means that they can

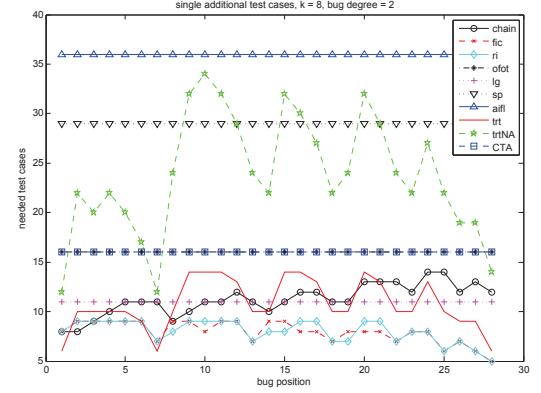


Fig. 7. $k = 8, t = 2$,single, additional test suites

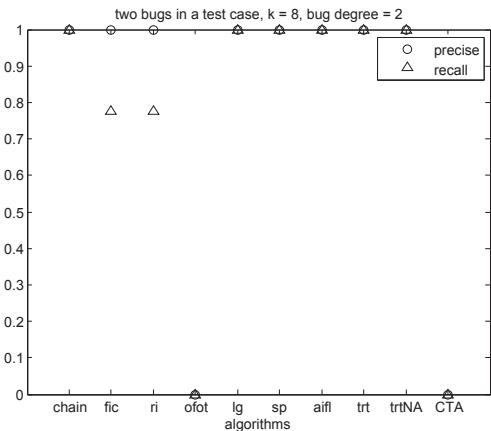


Fig. 8. $k = 8, t = 2$,double, additional test suites

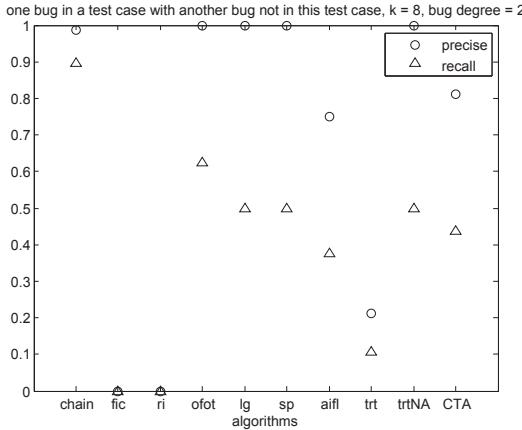
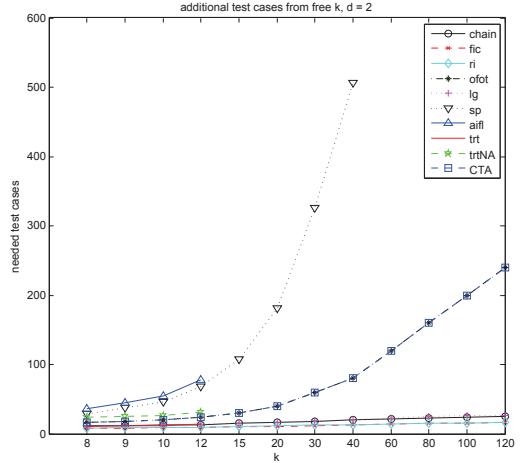
perfectly deal with the multiple MFSs in a failing configuration. As the opposite side, algorithms OFOT and CTA both get the recall 0 and precise 0. This is a signal that the two algorithms can't handle the condition that a failing configuration contain multiple MFSs. So for these two algorithms, they may need other techniques to assist them in facing such conditions. As for the remaining algorithms RI and FIC, they got precise 1 , but didn't reach recall 1, this is because they can't handle the condition if two MFSs have overlapped part. When they facing such condition, they can only identify one MFS of them.

The average number of configurations of these algorithms can both get recall and precise 1 is as follows:
we can find that our approach can also get a no-bad result among them

3)The third experiment's result is shown in fig 9. Similar to experiment 2, this figure just record the recall and precise of each algorithm. We can learn several facts of this result.

First, no algorithm can both get recall 1 and precise 1 in this condition. Thus shows that no algorithm can perfectly handle the importing problem.

Second, even though, there are some algorith-

Fig. 9. $k = 8, t = 2$,import, additional test suitesFig. 10. k is free, $t = 2$,single, additional test suites

m (OFOT, LG, SP, TRT-NA, our approach is almost get precise 1) can get precise 1, this is a signal that the importing problem has a little impact on them

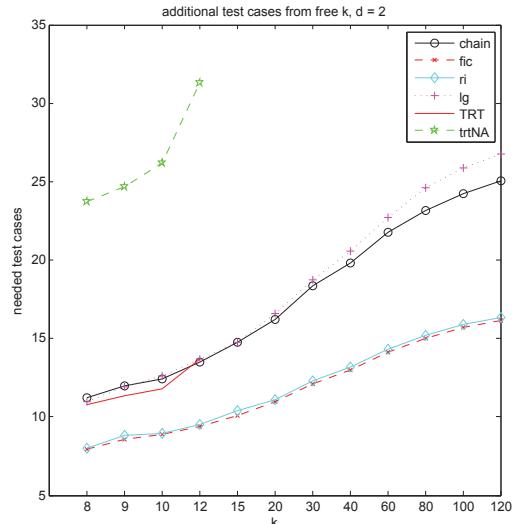
Third, our approach get the highest score in recall than others, the specific rank are *Chain* > *OFOT* > *LG* = *SP* = *TRT - NA* > *CTA* > *AIFL* > *TRT* > *FIC* = *RI*. It shows that our approach can find more MFSs than others when just feeding a failing configuration.

Fourth, algorithms *fic* and *ri* really suffers from this condition, as they get both recall and precise 0.

4)The data depicted in figure 10 shows that the average configurations needed of each algorithm for experiment 4. In this figure, the y-axis still reports the number of configurations while the x-axis represent a set of SUTs with the same number of configurations. These number of parameters are 8, 9, 10, 12, 15, 20, 30, 40, 60, 80, 100 and 120 respectively. And a point in this figure indicate the average number of configurations needed for the corresponding algorithm to identifying the MFS in these SUT with the same number of the parameters. For example, for the point whose x-axis is 8 corresponding to the algorithm "AIFL", we record the number of configurations needed to identify each SUT from the SUT with 8 parameters (totally $\binom{8}{2} = 28$ SUTs) and compute the average value of them, i.e., "".

We can easily learn that AIFL needs the largest number of configurations to identify MFS. Worse more, it need such big computing resource along with increase of k so that we even can't make it to identify the MFS in the SUT with number of parameters bigger than 12. The second largest is the algorithm "SP". Although it needs much smaller computing resource than "AIFL", we yet can't tolerate the time it need to identify the MFS with number of parameter bigger than 40.

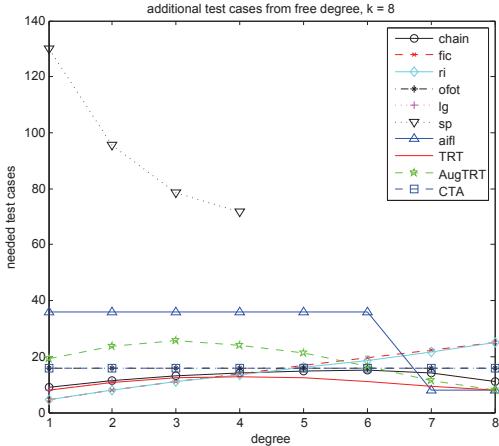
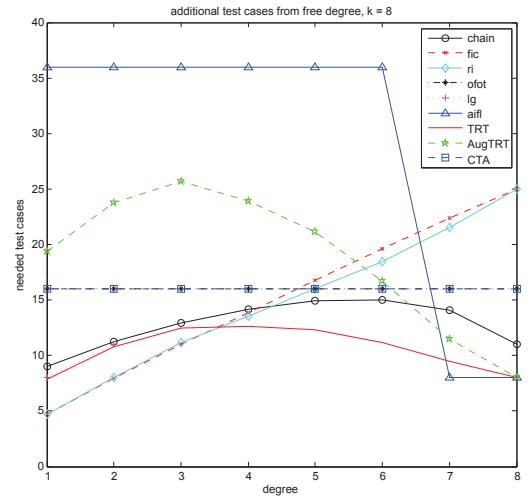
As the average number of configurations of "SP" increase quickly with the increase of k . It make the results of the remaining algorithms show unclearly in this figure. To get a better view of the remaining

Fig. 11. k is free, $t = 2$,single, additional test suites for some algorithms

algorithms, we enlarge the perspective for these algorithms in figure 11.

From this perspective, we first find that our previous approach TRT and TRT-NA also suffer from the large k . We can't let them identify the MFS with k bigger than 12. Apart form this two algorithms, all the remaining algorithms: chain, fic, lg, ri can complete this experiment. second, it also give us a intuitive rank of each algorithms according to the cost of configurations generated, i.e., *TRT - NA* > *LG* > *CHAIN* > *TRT* > *RI* > *FIC*. Combine With the two algorithms "AIFL" and "SP" we mentioned before, the final rank is *AIFL* > *SP* > *TRT - NA* > *LG* > *CHAIN* > *TRT* > *RI* > *FIC*.

5)Figure 12 gives us the result of the last experiment. This figure is organised similarly as figure 10, except the x-axis represents a set of SUTs have the same degree of MFS we inject in it. For a particular

Fig. 12. $k = 8$, t is free,single, additional test suitesFig. 13. $k = 8$, t is free,single, additional test suites for some algorithms

algorithm, a point in this figure shows the average number of configurations needed to identify the MFS in these SUTs with same degree MFS injected (there are $\binom{8}{m}$ SUTs for the degree m). Note that in this experiment we set the degree that SP can be identified 4 to distinguish with algorithm LG which can just identify the MFS with degree not bigger than 2.

From this figure, we can find that as we set parameter of SP to be 4, the number of configurations SP needed to identify the MFS in a failing configuration increase a lot than before. And as we expected, SP can just identify the MFS with degree not than 4. Apart this algorithm we enlarge the view of the remaining algorithm in figure 13 as experiment 4. From this clearer view, we can find several useful information: First, the LG method can only indentify the MFS with degree not than 2 as expected. Second, we can find the number of configurations of ri and fic increase quickly along with the degree increase. Third, CTA and CTA need the same configurations 16 regardless of what the degree is. Fourth, though AIFL perform worst in most cases, when d is 7 and 8 it can do better than other techniques. Fifth, the number of configurations needed of TRT-NA, TRT and our chain techniques first increase then decrease along with increase of d .

we will enlarge some of them to figure 13

8 EMPIRICAL CASE STUDY

While we got know that our approach have a better performance than others in several scenarios such as a failing test case contains multiple MFSs, the generated extra test case introduces newly MFS and so on from previous simulated experiments, it did not give us strong confidence in that our approach will also perform well in real software systems. One important reason of this is that we don't know whether these competitive scenarios for our approach existed in real software systems. So to eliminate the doubt we conduct a series empirical studies in this section.

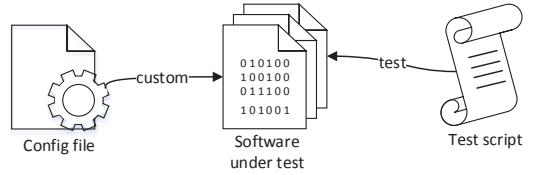


Fig. 14. proceeding

These empirical studies aimed to answer the following questions:

Q1: Is there any test case that contain multiple MFSs, and if so, do they overlapped each other? Do any of these MFSs have high degree and how likely does it introduce newly MFS when generate extra test cases.

Q2: How well of these approaches mentioned in previous sections performed when identifying the MFS in the real softwares?

Q3: If we combine the result of one approach with another one, does it give us a better result than both of them?

The subject systems for these studies are HSQLDB-B(2.0rc8). HSQLDB is a database management system written in pure java. All of them have millions of uncommented lines of code. And they all share the same proceedings depicted in Fig.14 when tested in the following studies.

As see from the figure, there are three modules in this proceeding: configuration file, software under test and the test script. The software under test is a set of software components which supply some interfaces, with which one can invoke the serving functions of the software. The configuration file is the file that can custom the properties of the software, and the test script is a executable file which test some functions of the software. In specific, the proceeding is trying different configuration of the software by changing the content of the configuration file, and

TABLE 7
input model of HSQLDB

SQL properties(TRUE/FALSE)		
table properties		
hsqldb.default_table_type	CACHED, MEMORY	
hsqldb.tx	LOCKS, MVLOCKS, MVCC	
hsqldb.tx_level	read_committed, SERIALIZABLE	
hsqldb.tx_level	read_committed, SERIALIZABLE	
Server properties		
Server Type	SERVER, INPROCESS	WEBSERVER
existed form	MEM, FILE	
Result Set properties		
resultSetTypes	TYPE_FORWARD_ONLY, TYPE_SCROLL_INSENSITIVE, TYPE_SCROLL_SENSITIVE	
resultSetConcurrencys	CONCUR_READ_ONLY, CONCUR_UPDATING	
resultSetHoldabilitys	HOLD_CURSORS_OVER_COMMIT, CLOSE_CURSORS_AT_COMMIT	
option in test script		
StatementType	STATEMENT, PREPAREDSTATEMENT	

then executing the same test script to observe the result.

8.1 experimental setup

Before we use these programs for study, we will take the following steps for each program to obtain the basic information of them, 1)Build the input configuration model of the subject program 2)execute the test case under each possible configuration and record their executed result. 3)list all the types of fault during testing and judge their priority.

We will depict these processes and results for each subject program in detail.

8.1.1 HSQLDB

Through searching the developers' community of HSQLDB in sourceforge, we found a buggy version with its faulty description in [], and a test script is attached which can reproduce the bug. The original test script only considered one option which have two values, remaining the other options set default values. To see what will happen when executing the test script under more different configurations, we need add some other options which may influence the behaviour of the database management software. From numerous options of HSQLDB, We chose some of the options that control the properties of table, some control the sql, some control the server and some control the result set. Along with the one in the original test script, we derive a input configuration model of HSQLDB listed in table 7.

TABLE 8
faulty and Priority

exception ID	lower priority	num of configs
1	2	36855
2	1	110558
3	1 2	58
4	1 2	1

We denoted this input model as $(2^{13}, 3^3)$. There are $2^{13} * 3^3 = 221184$ possible configurations in total. We executed the test script under each of the 221184 configurations and recorded their result. In specific, there are 147472 configurations triggered exceptions, while others passed this test. These exceptions can be classified to 4 types according to the exception traces. Through viewing the exception trace info and the ~~ADBTE~~ code of test script, we sorted out the priorities of these exceptions. Table 8 shows the detail of the 4 types of exceptions. The "exception ID" column shows the id of the exception. The "lower priority" column gives the IDs of exception which have lower priority than this exception. The "num of configs" means the number of test configurations under which the test script triggered this type of exception when executing.

8.2 Study 1: what the schemas like in practice

In the first study, we aimed to answer the Q1. We need to investigate the MFSs of the real softwares to see 1) Do there existed multiple MFSs in the same configuration? 2) if so, how many of them overlapped each other? 3) Do any of the MFSs have the degree larger than 2. 4) What the possibility of introducing newly MFSs when generating extra configurations. In particular we will inspect the MFSs in HSQLDB.

The first obstacle of case study 1 is that we don't know the MFSs in the real softwares. The original bug report page of these softwares[][] can give us hints of the MFSs of the software, but it is not enough for us to get accurate MFSs of the software. The reason is that we had added more options to test the script which resulted in more exceptions than reported and the original reported exception can also be related to more options that didn't be mentioned in the bug report.

So the only way to recognize the MFSs in the software is searching all the schemas in a configuration one by one, and then judging the state of the schema according to the definition and propositions. We will repeat the process in all the possible configurations of the software, this process is time-assuming which can be optimized by introducing hash table and adjusting the order when searching the schema of a configuration. We will omit these details as it is not point of this paper. Lastly after this time-assuming process we

got the accurate MFSs for each software, we recorded them for later use.

8.2.1 Measure the MFSs in the softwares

As MFSs in the softwares are figured out, we can easily count the number of configurations contain multiples MFSs and MFSs that overlapped each other, as well as count the number of MFSs that have a degree larger than 2. But for the possibility of introducing newly MFSs when generating extra configuration (brief as possibility of introducing), we yet can't measure it for we didn't define how to compute the "possibility". We will give a formula to compute the possibility of introducing next, before then we will explain how is this formula derived.

First we should understand under which condition will the event of introducing newly MFSs harms our identifying algorithm. Consider the following scenario.

For a test case of which we want to identify the MFS, say test case A, if we change one factor of it to generate a new test case, say test Case B. Then assume in this step we introduced a new failure-inducing schema, but meanwhile we didn't break the original failure-inducing schema in the test case A when we generate B, in other words, the MFS in A is still in B.

At this situation it will not influence our identifying result for that if we did not introduce the schema, our result is the same—trigger the same failure.

Consider another scenario, still for test case A, and we also change one factor of it to generate a new test case, say test Case C. Similarly, we introduced a new MFS, but different from the first scenario, this time we break the original MFS in the test case A when we generate C, which means the MFS in A is not in C.

At this situation it will influence our identifying result for that our expected result is that C should pass the test as we have broken the MFS in A. But the result failed at last. And we could owe this failure to some schema in A that is not the real MFS if we did nothing to deal with the introducing problem.

So we just need consider the possibility of the situation of simultaneously broking one MFS and introducing another MFS. This metric is related to the cost of changing one schema to another schema. The followed formula define the changing cost of two schemas:

$$\text{ChangeCost}(A, B) = |T(A, B)| + \sum_{i \in T(B, A)} (P_i / 2) + \sum_{i \in S(A, B)} (|A_i - B_i| / 2)$$

In this formula, A and B represent two different schemas. the denotation of $T(A, B)$ gives the parameters in A but not in B. And $S(A, B)$ means the parameters in both A and B, but their value is different. P_i refers to the number of values in the i th parameter of SUT, and A_i is the value of one factor in Schema A , the factor is the i th parameter of SUT.

Then the introduce rate of a SUT is defined as:

$$\frac{\sum_{a,b \in MFSs, a \neq b} (\text{ChangeCost}(a, b))}{|MFSs| \times |MFSs - 1|}$$

8.2.2 result and analysis

The statistic info of MFSs of the real software is listed in Table 9.

In this table, "exp ID" is the id of exception which the MFSs will trigger, "MFSs" gives the number of MFSs that can trigger this type of Exception show in "exp ID" column. Column "degree than 2" list the number of MFSs that have a degree larger than 2. "multiple MFSs" means the number of configurations that contain multiple MFSs. The column "overlapped MFSs" gives the number of configurations that contain MFSs that overlapped each other. And the last column "intro rate" shows the introduce possibility of newly MFSs when generating extra test configuration.

From this result, we will answer the sub-question of Case study 1 one by one.

1)Do there existed multiple MFSs in the same configuration?

Answer: yes. Although it is rare among the failing configurations(for exception 1, there are 36855 configurations trigger the same exception, and only 8 configurations contain multiple MFSs), but it do exist in the configurations of real software, we can see except the 4th exception which just has one MFS, all the other exception have configurations have multiple MFSs, which are 8,25,1 respectively.

2)if so, how many of them overlapped each other?

Answer: most of them are overlapped each other. We learned configurations which have overlapped MFSs is 8, 1, 1 respectively. As we all know, the configurations have overlapped MFSs is just one part of the configurations have multiple configurations. Considering the configurations contain multiple MFSs is rare, which are 8 , 25, 1 respectively, then the We think it is a high possibility when we encounter the situation that configurations contain multiple MFSs and these MFSs overlapped with each other.

One possible explanation for this phenomenon may be real softwares may have many branches, and these branches may have iteration, so for some MFSs , they may share the same entrance of the branch.

3)Do any of the MFSs have the degree larger than 2.

Answer: yes. It is clearly that almost all the MFSs in the software have a degree than 2, except one MFS for exception 1 (8 among 9 MFSs have degrees larger than 2) and one MFS for exception 2(23 among 24 MFSs have degrees larger than 2).

4)What the possibility of introducing newly MFSs when generating extra configurations?

Answer: the possibility varies from one to another.

We can learn from the table that the intro rate of exception 1 is 0.1028, which are the biggest than others, and the smallest is the exception 3, which is 0.0026. They differ markedly, so the possibility

TABLE 9
MFSs information of each exception

exp ID	MFSs	degree than 2	multiple MFSs	overlapped MFSs	intro rate
1	9	8	8	8	0.1028
2	24	23	25	1	0.0145
3	56	56	1	1	0.0026
4	1	1	0	0	-

of introducing newly MFSs depends on the specific exception and may have big difference among each other.

8.3 Study 2: how these algorithms behave when applied in real softwares

The second study aims to answer the Q2. We will evaluate the performance of each algorithm in identifying the MFSs of the real softwares.

8.3.1 Study setup

To conduct this case study, We will feed one failing configuration to each algorithm and use them to identify the MFSs. Then we will compare the results get by each algorithm to the real MFSs in that configuration given in the study 1 respectively. The comparison metrics is similar to the simulated experiment, which consists of the number of extra test configurations needed, the precise and the recall. To be fair, no other information is given to each algorithm except the feeded failing configuration. We will repeat this comparison for each failing configuration of a exception. At last we will report the average number of test configurations , precise and recall for each algorithm.

As we will take a large number of configurations to identify(147472 for HSQLDB), this is too big compact to some algorithms, such as TRT, IterAIFL, they can't even accomplish the task. Some other algorithm can complete this task but with a huge time cost, such as SP, it will take weeks or even months to complete. We will omit these algorithms in our comparison to make this study available in a relative short time. As a result, we just choose ChainFeedBack , FIC, RI, OFOT, LG, CTA as our comparison algorithms.

8.3.2 result and analysis

The result is shown in table 10. In this table, Column "exp ID" still means the ID for the specific exception. Column "algorithm" gives the specific algorithm measured in this row. Column "num of extra configs" shows the average number of extra configurations needed to generate to identify the MFSs. Column "recall" and "precise" respectively shows the average recall and precise which have been defined in the simulated experiment for each algorithm.

We will discuss the result by column:

1)the number of extra configurations: We may get two points in this column: a. no algorithms always need the smallest configurations in all conditions. For example,FIC needs the smallest configurations for exception 1(9.023) and exception 2(6.016) but needs the largest number of configurations for exception 4(65.0) as FIC may need a high cost dealing high way degree schema. RI performs good at exception 2(6.016), but performs bad at exception 3(62.62) and 4 (65.0). LG needs the smallest number of configurations for exception 3(5.344) and exception 4(5.0), but for exception 1 and 2 it doesn't perform as good as FIC and RI. b. OFOT and CTA performs almost the same for all the exception exceptions(almost as 19.0, exception the OFOT for the exception 1(19.0007)), this is because they are not the adaptive method, they just change one factor one time for all the conditions. And the unique different is that OFOT add some more configurations to deal the introduce problem. c. Our algorithms performs moderately better than OFOT for exception 1 and 2 but worse than others, and better than FIC and RI for exception 3 and 4 and worse than others. This is because our algorithm consider more conditions such as introduce, overlapped and high degree problem.

2)recall : our algorithms perform the best in all the circumstances for this metric, which means our algorithms can find more MFSs than others algorithms when identifying the MFSs in a failing configuration. The reason why our approach have advantages over others is that we consider more scenarios when fed a failing configurations, especially we took account of the case multiple , overlapped MFSs in a test configuration and the case when the MFSs have a high degree. Other approaches may consider some of them, but none of them consider all these scenarios. Another point that needed note is that for exception 1, 3 and 4, the recall of LG algorithm is 0, it is a signal that in this three circumstances, the MFSs in the SUT have degrees larger than 2.

3)precise : algorithm OFOT and our approach performs better than others (our approach is litter weaker than OFOT for the exception 2, which ours is 0.9996, OFOT is 0.9997). As we have discussed before, the condition that introducing newly configurations when generating extra test configurations can make the identifying result incorrectly, so it is a nature idea that if the algorithm deal with the introducing newly

TABLE 10
comparison in real softwares

exp ID	algorithm	num of extra configs	recall	precise
1	ChainFeedBack	15.061	1.0	0.9995
	FIC	9.023	0.9991	0.9988
	RI	9.024	0.9991	0.9988
	OFOT	19.007	0.9997	0.9995
	LG	11.000	0.0	0.0
	CTA	19.0	0.9997	0.9995
2	ChainFeedBack	13.058	0.9999	0.9996
	FIC	6.0160	0.9999	0.9997
	RI	6.0161	0.9999	0.9997
	OFOT	19.0	0.9997	0.9997
	LG	11.0005	0.9998	1.0
	CTA	19.0	0.9997	-
3	ChainFeedBack	19.793	0.9827	0.9827
	FIC	61.6551	0.8879	0.89655
	RI	62.6206	0.9396	0.9482
	OFOT	19.0	0.9827	0.9827
	LG	5.3448	0	-
	CTA	19.0	0.9137	0.9137
4	ChainFeedBack	19.0	1.0	1.0
	FIC	65.0	1.0	1.0
	RI	65.0	1.0	1.0
	OFOT	19.0	1.0	1.0
	LG	5.0	0.0	-
	CTA	19.0	1.0	1.0

MFSs, the algorithm can get a more accurate result than those approaches don't consider this scenario. The result data show in this table is coincide with this idea. We can also learn that the machine learning approach- CTA can also get a well result (0.9995, - , 0.9137, 1.0, slightly weaker than OFOT and our approach) in this column. Although this approach didn't consider the introducing scenario, the nature of dealing noisy data in machine learning can give this approach a some relief from those bad affects.

In addition, we can find some regular among the exceptions, i.e., for all the algorithms , the recall and precise is best in 4, second best in 2 , little worse in 1, the worst in 3. This is coincided with the intro rate, which means that the intro rate really have a influence for the performance in algorithms.

Overall, the answer to Q2 is: Our algorithms performs good at recall and precise, and moderately in num of extra configs. This is coincided with the result in simulated experiment.

8.4 Study 3: Is combination useful?

We can learn from the result in the second case study that none of these algorithms can be the best among all the scenarios. So a nature question is that can we combine the result of multiple algorithms to find a better result. In this study we will carry out some experiments to answer this question, i.e., the Q3.

8.4.1 Study setup

To utilize the result of each algorithm, we build a voting system. Thus, choose some algorithms and take a voting schedule: if the a schema is identified more than one algorithms, we take it as MFS, otherwise, we will discard it. We will vary the number of algorithms and adopt different combination of algorithms to observe the result. In specific we will take 2, 3, 4, 5 and 6 algorithms as the voters respectively. In addition, we will take all the combination of algorithms for a specific number. Take 2 for example we will choose ChainFeedBack and OFOT as a combination, ChainFeedBack and RI as another combination and so on. The total number combination of algorithms for a specific number n is $\binom{n}{2}$.

For each combination of algorithms, we will record the metric "recall" and "precise" of voting result. To measure if the voting system is useful, we compare this voting "recall" to the highest "recall" of the algorithm among the combination. The same comparison is also made to the "precise" metric. Note that the highest "recall" and the highest "precise" is not necessarily belong to the same algorithm in a combination. As a example, for the combination "OFOT" and "RI" for the exception 2. The highest value of "recall" is , which belongs to "OFOT". and the highest value of "precise" is, which is belong to the "RI". And our voting result will compare to the

recall "" and precise "" respectively.

8.4.2 result and analysis

The detail comparison result is listed in table 11. In this table, "Column" means the "+" signal. means a promotion over the highest . and the "-" signal means a decreasing of the highest.

The answer to Q3 is that: yes, it has promotion, but very little.

8.5 threats to validate

how seeded bugs are representative auto oracle, if we don't have a correct version, or we can't determine a test case is faulty and wrong, then what can I do? timing result One worrying aspect of this research is that it seems to consider only the number of tests and number of faults uncovered. In the practice of testing, it is as important or more important to know testing times.

9 RELATED WORKS AND DISCUSS

Detail list the history. In the experiment section, we just give some brief introduction.

Nie's approach in [3] and [6] first separates the faulty-possible tuples and healthy-possible tuples into two sets. Subsequently, by changing a parameter value at a time of the original test configuration, this approach generates extra test configurations. After executing the configurations, the approach converges by reducing the number of tuples in the faulty-possible sets. Delta debugging [5] proposed by Zeller is an adaptive divide-and-conquer approach to locating interaction fault. It is very efficient and has been applied to real software environment. Zhang et al. [4] also proposed a similar approach that can identify the failure-inducing combinations that has no overlapped part efficiently, Colbourn and McClary [7] proposed a non-adaptive method. Their approach extends the covering array to the locating array to detect and locate interaction faults. C. Martiez [8-9] proposed two adaptive algorithms. The first one needs safe value as their assumption and the second one remove the assumption when the number of values of each parameter is equal to 2. Their algorithms focus on identifying the faulty tuples that have no more than 2 parameters. Ghandehari.etc [10] defines the suspiciousness of tuple and suspiciousness of the environment of a tuple. Based on this, they rank the possible tuples and generate the test cases. Although their approach imposes minimal assumption, it does not ensure that the tuples ranked in the top are the faulty tuples. Yilmaz [11] proposed a machine learning method to identify inducing combinations from a combinatorial testing set. They construct a classified tree to analyze the covering arrays and detect potential faulty combinations. Beside this, Fouch? [12] and

Shakya [13] made some improvements in identifying failure-inducing combinations based on Yilmaz' work.

We list a comprehensive detail and comparison table 12as followed.

But even we get the failure-inducing schemas, it is still having a gap to fetch the failure causing root from the code. Such as int the TCAS, we get the failure-inducing schemas, and this is caused by a code mutation in the code such as followed: So in the future, we will analysis the relationship between the failure-inducing schemas and the real code causing.

10 CONCLUSION

The conclusion goes here.

APPENDIX A

PROOF OF THE FIRST ZONKLAR EQUATION

Appendix one text goes here.

APPENDIX B

Appendix two text goes here.

The procedure :

mutant_a
mutant_r
ispendingschema
makechain

ACKNOWLEDGMENTS

This work was supported by the National Natural Science Foundation of China (No. 61272079), the Research Fund for the Doctoral Program of Higher Education of China (No.20130091110032), the Science Fund for Creative Research Groups of the National Natural Science Foundation of China(No. 61321491), and the Major Program of National Natural Science Foundation of China (No. 91318301)

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TABLE 11
the combination of algorithms

	ChainFeedBack	FIC	RI	OFOT	LG	CTA
ChainFeedBack	-	\	\	\	\	\
FIC	\	-	\	\	\	\
RI	\	\	-	\	\	\
OFOT	\	\	\	-	\	\
LG	\	\	\	\	-	\
CTA	\	\	\	\	\	-

TABLE 12
The comparison of each algorithms

algorithms	time complexity	space complexity	multiple schemas	strength of schema	safe value	result style
CMINFS	$O(\log(n))$	$O(n)$	yes, no limit, can overlapped	1 n	can handle	precise
FIC	$O(\log(n))$	$O(n)$	yes, no limit, can't overlapped	1 n	can't handle	precise
TRT	$O(\log(n))$	$O(2^n)$	yes, no limit, can overlapped	1 n	can handle	precise
Spetrum	$O(n*t)$	$O(n * t)$	yes, no limit, can overlapped	1 n	can handle	a rank of possible

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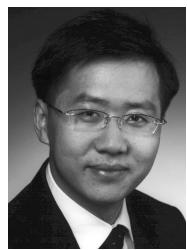
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