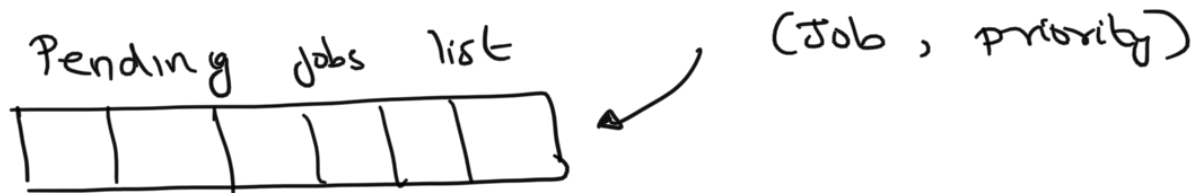


# Priority queues



→ new jobs may join list at any time

→ Extract job with highest priority in list to be executed

How to maintain list of pending jobs and priorities so that jobs can be added and extracted efficiently?

Priority queue: Abstract Data structure storing

pending jobs with priorities with 2

operations

insert()

delete-max()

Naive:

$n = \text{size of list}$

(Linear structure)

Unsorted list:

insert  $O(1)$

delete-max  $O(n)$

sorted list:

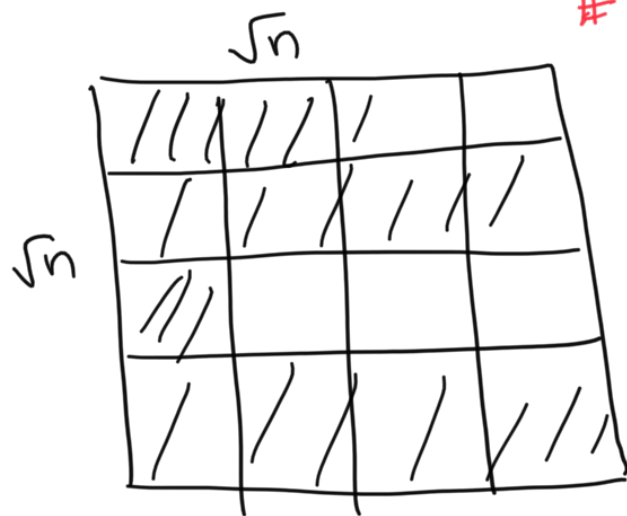
insert  $O(n)$

delete-max  $O(1)$

Processing  $n$  jobs  $\rightarrow$   $n$  insertions  
 $n$  deletions

$\rightarrow O(n^2)$  for both  
naive approaches

Better (if you know beforehand  
# of jobs  $\leq n$ ) (2-D structures)



$\sqrt{n} \times \sqrt{n}$  array

each row  
maintained in sorted  
array and

also maintain how  
many jobs in each  
row are there.

Insert  $\rightarrow$  check end element of each row  
and insert first time there is  
space and inserting element  $\geq$   
end element of current row

$O(\sqrt{n})$

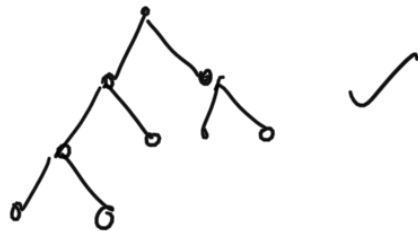
delete  $\rightarrow$  pick max (end elements of rows)  
and delete it.  $O(\sqrt{n})$

To process  $n$  jobs:  $\left. \begin{array}{l} n \text{ inserts} \\ n \text{ delete-max} \end{array} \right\} O(n\sqrt{n})$   
 $= O(n^{3/2})$

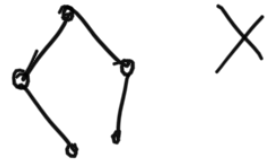
much better

## Heaps

- binary tree with  $n$  nodes
- "balanced"  $\leftarrow$  height  $\log n$



balanced



unbalanced

given  $n$ , unique balanced binary tree with  $n$  nodes  
 $\uparrow$   
 $\#$  of jobs

Can implement priority queues so that

insert :  $O(\log n)$

delete :  $O(\log n)$

processing  $n$  jobs :  $O(n \log n)$

(can grow heap as we go along.

no need to fix  $\#$  of nodes in heap in advance)

