

Matrix multiplication (DP)

$$A = (a_{ij})_{m \times n}$$

$$B = (b_{ij})_{n \times p}$$

$$C = AB = (c_{ij})_{m \times p}$$

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj} \quad - \quad O(n)$$

So computing C $- O(mnp)$

many matrix multiplications

$$(AB)C = A(BC)$$

$$\underbrace{m \times n \quad n \times p \quad p \times q}_{\text{}} \quad O(mnp)$$

$$\underbrace{m \times p \quad p \times q}_{\text{}} \quad O(mpq)$$

$$m \times n$$

$$\underbrace{n \times p \quad p \times q}_{\text{}} \quad O(npq)$$

$$\underbrace{m \times n \quad n \times q}_{\text{}} \quad O(mnq)$$

$$\downarrow \\ m \times q$$

$$O(npq + mnq)$$

$$O(mnp + mpq)$$

$$m = 1 \quad n = 100 \quad p = 1 \quad q = 100$$

$$mnp = 100$$

$$mpq = 100$$

$$\underline{200}$$

$$npq = 10000$$

$$mnq = 10000$$

$$\underline{20000}$$

more efficient to find
 $ABC = (AB)C$

than $A(BC)$

Problem Given $M_1 \quad M_2 \quad M_3 \quad \dots \quad M_n$

* dim of $M_i : r_i \times c_i$

* $c_i = r_{i+1} \quad \forall \quad 1 \leq i < n$

Find optimal order to compute product

$$(M_1 M_2 \dots M_k) (M_{k+1} \dots M_n)$$

$$\begin{array}{ccc} \downarrow & \text{last bracket} & \uparrow \\ r_1 \times c_k & & r_{k+1} \times c_n \end{array} \quad \left. \vphantom{\begin{array}{ccc} \downarrow & \text{last bracket} & \uparrow \\ r_1 \times c_k & & r_{k+1} \times c_n \end{array}} \right] \quad \begin{array}{l} \text{last multiplication} \\ \text{cost} \\ r_1 c_k c_n \end{array}$$

$$c_k = r_{k+1}$$

$$\begin{aligned} \text{Total cost} &= r_1 c_k c_n + \text{cost}(M_1 \dots M_k) \\ &\quad + \text{cost}(M_{k+1} \dots M_n) \end{aligned}$$

if last bracket is separated @ k

$$\text{Optimal cost} = \min_{1 \leq k < n} \left[r_1 c_k c_n + C(M_1 \dots M_k) + C(M_{k+1} \dots M_n) \right]$$

As you solve subproblems, you'll have to find $C(M_i \dots M_j) \dots$

So DP - find $C(i, j) := \text{cost}(M_i M_{i+1} \dots M_j)$

$$C(i, i) = 0$$

(no multiplication)

$$C(i, j) = \min_{i \leq k < j} [C(i, k) + C(k+1, j) + r_i c_k c_j]$$

only want $C(i, j)$ $i < j$

$$\text{eg } C(i, i+1) = \min_{\substack{k=i \\ \rightarrow i \leq k < i+1}} [C(i, k) + C(k+1, i+1) + r_i c_k c_{i+1}]$$

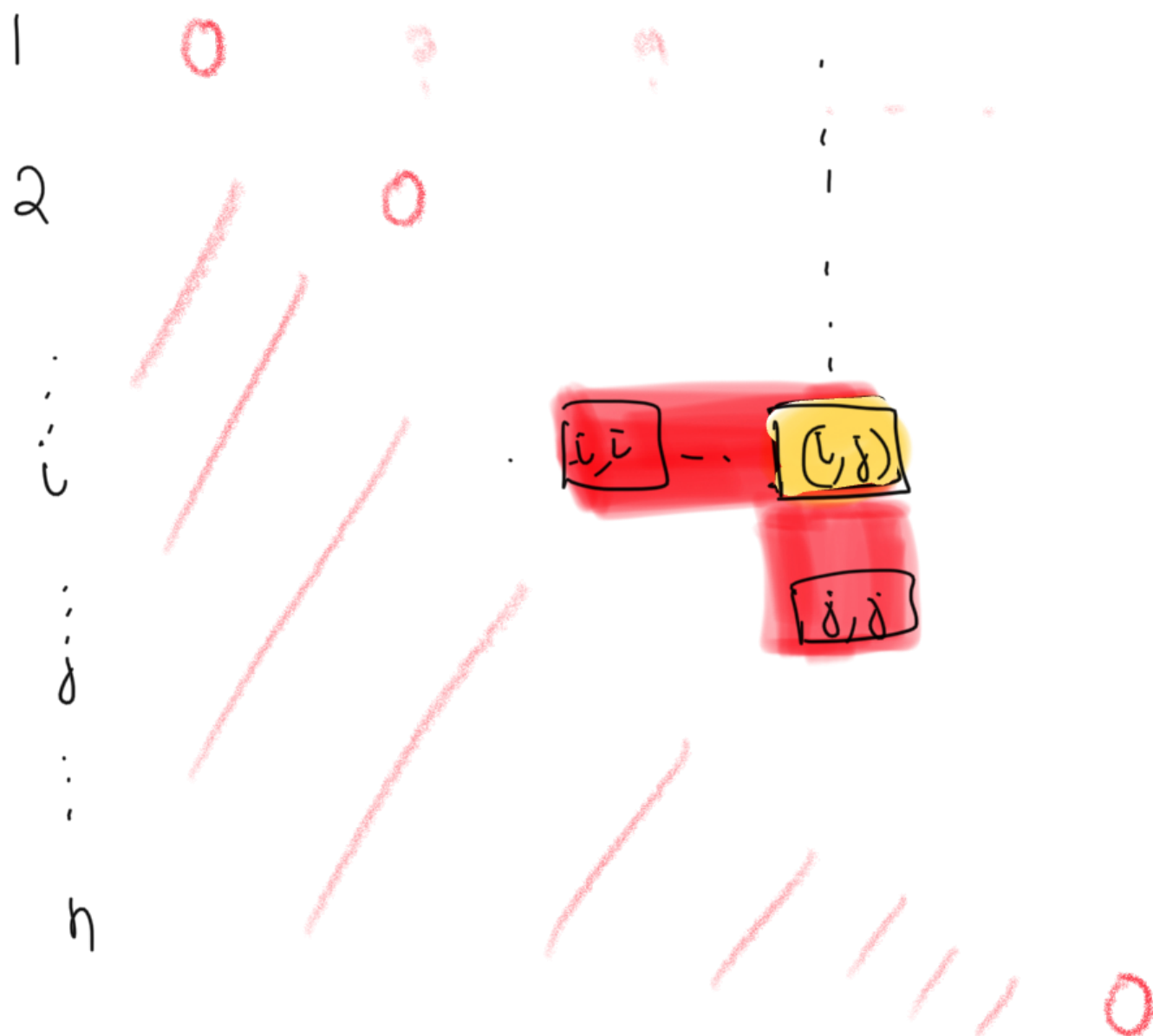
$$= C(i, i) + C(i+1, i+1) + r_i c_i c_{i+1}$$

$$= r_i c_i c_{i+1}$$

j (2nd index)

1 2 ... i - j ... n

1st index
 i

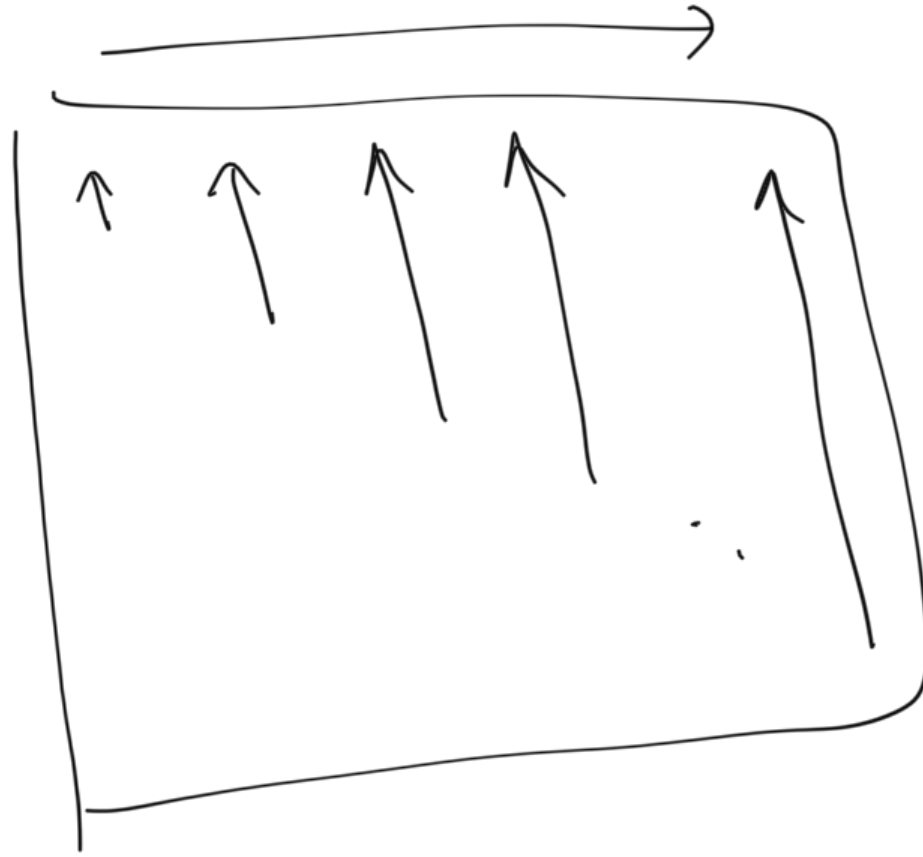


$C(i, j)$

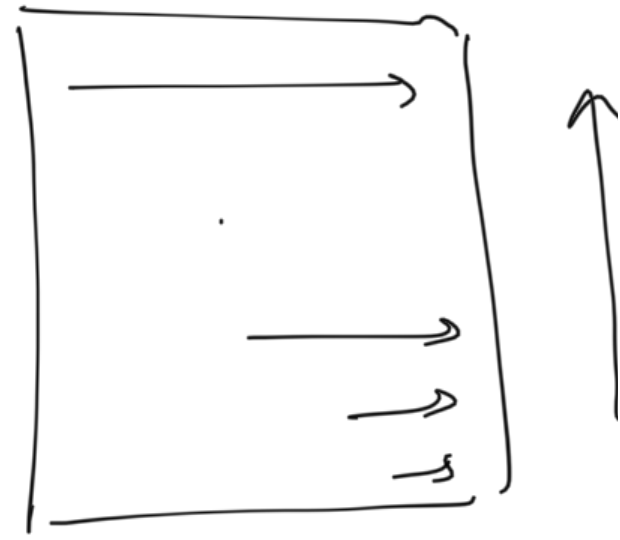
uses

$C(i, k)$
 $C(k+1, j)$ for $i \leq k < j$

Fill up

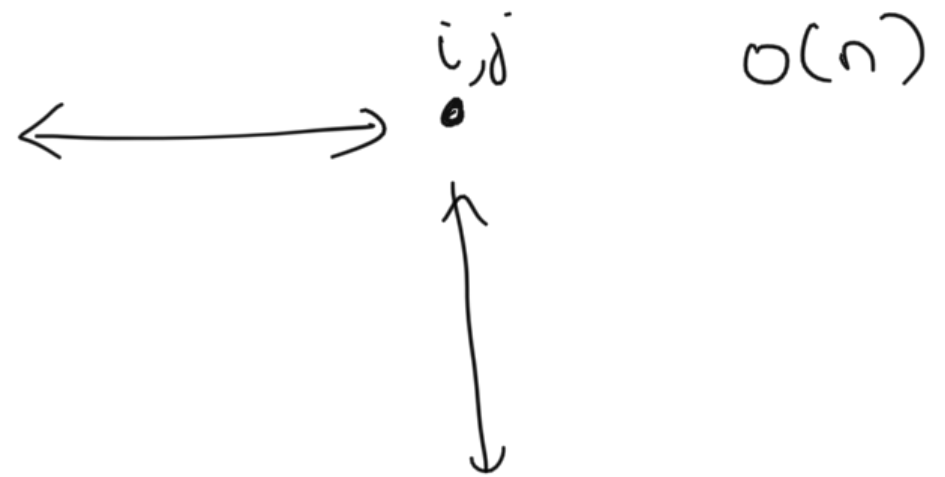


or



Complexity analysis : Filling up $O(n^2)$ table

Filling up 1 entry



so $O(n^3)$ overall complexity