## Sliding windows

( Ref: https://medium.com/outco/how-to-solvesliding-window - problems-28d 67 601 a66)

## Min. window substang

Given string s, and a set T of characters, bind min window in S which will contain all characters in T

Naive: Check S(i) S(i+1)... S(j) & ixit

If it works. Let n=1S1

O(n2) pairs (i, i) to check, each pair takes

O(n) to check. So o(n3) algorithm

I with memoization, can become o(1)]

Sliding window solution: O(n)

window

Anny

-> more fast pointer to grow your window till
you see all charaders

fast pointer

- -) once you find a valid window, start sliding slow pointer up Shrinking window till no longer valid
  - -> Take min / all valid windows

Slow pointer

(eg) s= adobe code banc! t = { a, b, c}

adobe code banc

a do be code banc

adobe code banc

a do be code banc

ado be code banc

ado be code banc

ado be code banc

a do be code banc

a do be code banc

First valid window ] Length = 6

[ shrink to set first [ Twobow]

[ second valid window length = 10]

Is hounk valid, length=9]

Is hrunk valid, length 8]

Ishrunk valid, length 7]

Is hounk valid, length of

[ shrink to set second Invalid window]

[ yalid wordow, length 7]

[ valid window, length 6]

[ valid window, lengths]

[ valid wordow, length 4]

[Terminate, cant resure fout pointer up]

min window Length = 4

## Impenentation

-rKeep a count dichonary

- -) update it so that it reflects the count of characters in convert window.
- -> This will help determine validity of window

Complexity: O(n) because both pointers
moving forward [at most n times]

cooredness:

en t

I then any window 2 B,f] will be valid but longer of so don't case about it for computing min valid window length