

期中作业

hw-6 (2023/10/31)

姓名:

学号:

p.44: (8) Let \mathcal{A} be a wf. $((\neg p_1 \rightarrow p_2) \rightarrow (p_1 \rightarrow \neg p_2))$. Show that L^+ , obtained by including this \mathcal{A} as a new axiom, has a larger set of theorems than L . Is L^+ a consistent extension of L ? (10 points)

(注意: 此题有两问)

Your proof:

p_1	p_2	$(\neg p_1 \rightarrow p_2) \rightarrow (p_1 \rightarrow \neg p_2)$									
T	T	F	T	T	T	F	T	F	F	T	
T	F	F	T	T	F	T	T	T	T	F	
F	T	T	F	T	T	T	F	T	F	T	
F	F	T	F	F	F	T	F	T	T	F	

For the first question: Obviously $\mathcal{A} = ((\neg p_1 \rightarrow p_2) \rightarrow (p_1 \rightarrow \neg p_2))$ is not a tautology by above truth table, then by the **Soundness Theorem**, \mathcal{A} is **not** a theorem of L , while it is a theorem of L^+ , therefore L^+ has a larger set of theorems than L .

For the second one: L^+ is consistent. For suppose otherwise, then there is a formula \mathcal{B} such that

$$\vdash_{L^+} \mathcal{B} \quad \text{and} \quad \vdash_{L^+} \neg \mathcal{B}.$$

Since L^+ is obtained by including $\mathcal{A} = ((\neg p_1 \rightarrow p_2) \rightarrow (p_1 \rightarrow \neg p_2))$ as an extra axiom then L , hence we have that (note that the **subscript** of \vdash)

$$\mathcal{A} \vdash_L \mathcal{B} \quad \text{and} \quad \mathcal{A} \vdash_L \neg \mathcal{B}.$$

By the **Deduction Theorem**,

$$\vdash_L \mathcal{A} \rightarrow \mathcal{B} \quad \text{and} \quad \vdash_L \mathcal{A} \rightarrow \neg \mathcal{B},$$

which means that $(\mathcal{A} \rightarrow \mathcal{B})$ and $(\mathcal{A} \rightarrow \neg \mathcal{B})$ are tautologies according to the **Soundness Theorem**. Then by the definition, for any valuation v we have that $v(\mathcal{A} \rightarrow \mathcal{B}) = T$ and $v(\mathcal{A} \rightarrow \neg \mathcal{B}) = T$, which implies that $v(\mathcal{A}) = F$, that is, \mathcal{A} is a *contradiction*. But this is impossible by the truth table of \mathcal{A} . Contradiction! \square

p.44: (10) Let L^{++} be the extension of L obtained by including as a fourth axiom *scheme*:

$$((\neg \mathcal{A} \rightarrow \mathcal{B}) \rightarrow (\mathcal{A} \rightarrow \neg \mathcal{B})).$$

Show that L^{++} is inconsistent. (Hint: see Chapter 1 exercise 7 (p.10))

(10 points)

Your proof:

method-(1)

Let

$$\top = (p \rightarrow p) \quad \text{and} \quad \varphi = (\neg \top \rightarrow \top) \rightarrow (\top \rightarrow \neg \top),$$

clearly $\vdash_{L^{++}} \varphi$ (i.e., let $\mathcal{A} = \mathcal{B} = \top$). It is easy to check, say using truth table, that φ is a contradiction, hence $\neg \varphi$ is a tautology. By the **Completeness Theorem**, $\vdash_L \neg \varphi$, and then $\vdash_{L^{++}} \neg \varphi$ since L^{++} is a extension of L .

But we have that $\vdash_{L^{++}} \varphi$ and $\vdash_{L^{++}} \neg \varphi$, by the definition, L^{++} is inconsistent as required. \square

(continue on next page)

method-(2) (下面这个证明来自 吴家儒 同学, 这种证明很直接且颇具暴力美学, 再次感谢吴家儒同学为大家带来如此精彩的证明 ♡♡♡)¹

Since $\vdash_L (p \rightarrow p)$ (cf. *Example 2.7-(a)* in page 31), we have that $\vdash_{L^{++}} (p \rightarrow p)$ obviously. And let $(L4)$ denotes the *fourth axiom scheme* of L^{++} , that is,

$$(L4) \quad ((\neg \mathcal{A} \rightarrow \mathcal{B}) \rightarrow (\mathcal{A} \rightarrow \neg \mathcal{B})).$$

Considering the following proof sequence in L^{++} :

1. $[\neg(p \rightarrow p) \rightarrow (p \rightarrow p)] \rightarrow [(p \rightarrow p) \rightarrow \neg(p \rightarrow p)]$ (instance of $L4$)
2. $[(\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow ((p \rightarrow p) \rightarrow \neg(p \rightarrow p))] \rightarrow$
 $[((\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow (p \rightarrow p)) \rightarrow ((\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow \neg(p \rightarrow p))]$ (instance of $L2$)
3. $((\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow (p \rightarrow p)) \rightarrow ((\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow \neg(p \rightarrow p))$ ($1 + 2, MP$)
4. $(p \rightarrow p) \rightarrow [(\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow (p \rightarrow p)]$ (instance of $L1$)
5. $(p \rightarrow p)$ ($(p \rightarrow p)$ is a theorem of L , so is for L^{++})
6. $(\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow (p \rightarrow p)$ ($4 + 5, MP$)
7. $(\neg(p \rightarrow p) \rightarrow (p \rightarrow p)) \rightarrow \neg(p \rightarrow p)$ ($6 + 3, MP$)
8. $(p \rightarrow p) \rightarrow (\neg(p \rightarrow p) \rightarrow (p \rightarrow p))$ (instance of $L1$)
9. $\neg(p \rightarrow p) \rightarrow (p \rightarrow p)$ ($5 + 8, MP$)
10. $\neg(p \rightarrow p)$ ($9 + 7, MP$)

Hence $\vdash_{L^{++}} \neg(p \rightarrow p)$, together with previous $\vdash_{L^{++}} (p \rightarrow p)$, L^{++} is inconsistent as desired. \square

.....作业反馈

- 大部分人还是没有区分「元语言」和「对象语言」, 所以严格来说很多人的回答都是不合法的甚至是错误的。不过改作业的时候已经采取十分宽容的态度了, 还希望大家一定要重视这点, 这对后续的逻辑学习是十分重要的。
- 依旧强烈建议不要使用「简化真值表」, 这并不是说「简化真值表」是什么洪水猛兽大家碰不得, 只不过照现在的作业来看, 同学一画「简化真值表」就容易画错。
- 虽然很多同学借鉴了教材 *p.205* 的提示, 但这种提示往往省略了超多细节, 这些细节应该要补充完整的, 直接抄书行不得! 一个证明首先要说服自己才能说服别人!
- 建议以后用黑笔作答, 期末考试时也是一样的。
- 很多同学都误用了 $(L3) : (\neg \mathcal{A} \rightarrow \neg \mathcal{B}) \rightarrow (\mathcal{B} \rightarrow \mathcal{A})$ 公理, 如下的公式并不是 $(L3)$ 公理的一个实例:

$$(p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p) \quad \text{或} \quad (p \rightarrow \neg q) \rightarrow (q \rightarrow \neg p)$$

单单只使用公理模式 $(L3)$ 得不到上述公式是 L 的定理的, 注意否定符号的位置。

¹虽然下面陈述的和家儒同学的原始版本有亿点点不同... 再赞叹一次, 家儒同学的证明真的很有暴力美学, 爱了爱了👍

- 同样容易误用的是 **Proposition 2.19**:

Let L^* be a consistent extension of L and let φ be a formula which is not a theorem of L^* . Then L^{**} is also consistent, where L^{**} is the extension of L obtained from L^* by including $(\neg\varphi)$ as an additional axiom. (p. 40)

显然 L 是其本身的一个一致扩张，并且很多人做第 8 题第二问的时候，确实证明了 $\vdash_L \neg\mathcal{A}$ ，然后直接运用 **Prop. 2.19** 就说 L^+ 是 L 的一致扩张，这中间其实还有一个 gap 要补充的。

根据 **Prop. 2.19** 和 $\vdash_L \neg\mathcal{A}$ 我们只能得到 $L \cup \{\neg\neg\mathcal{A}\}$ 是一致的（注意否定的个数），而题目中的是 $L^+ = L \cup \{\mathcal{A}\}$ 。虽然语义直观上 \mathcal{A} 和 $\neg\neg\mathcal{A}$ 是一个意思，但是仅仅作为字符串来说二者是完全不同的东西。因此，如果硬是要用 **Prop. 2.19** 的话，我们就必须还得论证： $L \cup \{\neg\neg\mathcal{A}\}$ 和 $L \cup \{\mathcal{A}\}$ 是同一个系统。然而这在教材中是没有明确说明的。²

希望大家可以从这点看出逻辑证明的严格性。

- 抄作业的情况有点严重呀！虽然鼓励同学们相互讨论，但写作业的时候也别直接抄呀，都做对就还好啦，错都错一样的话就很难说过去了 :(

²在其他逻辑著作中，这往往涉及到一种称之为「置换」的操作，要注意「置换」和「代入」不同。