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## Problem Set #2

5 questions



1.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n) = 7T(n/3) + n^2$ . What's the overall asymptotic running time (i.e., the value of T(n))?

- \$\$\theta(n\log n)\$\$
- \$\$\theta(n^2)\$\$
- \$\\theta(n^2 \log n)\$\$
- \$\theta(n^{2.81})\$\$

1 point

2.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence  $T(n) = 9T(n/3) + n^2$ . What's the overall asymptotic running time (i.e., the value of T(n))?

- O \$\$\theta(n\log n)\$\$
- \$\$\theta(n^2\log n)\$\$
- \$\theta(n^{3.17})\$\$
- \$\$\theta(n^2)\$\$

1 point

3.

This question will give you further practice with the Master Method. Suppose the running time of an algorithm is governed by the recurrence T(n) = 5T(n/3) + 4n, What's the overall asymptotic running time (i.e., the value of T(n))?

- **O** \$\$\theta(n^2)\$\$
- \$\$\theta(n^{\log\_{3}(5)})\$\$
- \$\$\theta(n^{\frac{\log 3}{\log 5}})\$\$
- \$\\theta(n^{2.59})\$\$
- \$\theta(n^{5/3})\$\$
- **O** \$\$\theta(n\log(n))\$\$

1 point

## 4.

Consider the following pseudocode for calculating \$\$a^b\$\$ (where a and b are positive integers)

```
FastPower(a,b):
1
2
      if b = 1
3
        return a
4
      else
5
        c := a*a
       ans := FastPower(c,[b/2])
6
      if b is odd
        return a*ans
      else return ans
10 end
```

Here [x] denotes the floor function, that is, the largest integer less than or equal to x.

Now assuming that you use a calculator that supports multiplication and division (i.e., you can do multiplications and divisions in constant time), what would be the overall asymptotic running time of the above algorithm (as a function of b)?

- **O** \$\$\Theta(b)\$\$
- \$\Theta(\sqrt b)\\$\$
- \$\Theta(\log(b))\$\$
- **O** \$\$\Theta(b\log(b))\$\$

1 point

5.

Choose the smallest correct upper bound on the solution to the following recurrence: \$T(1) = 1\$\$ and  $\$T(n) \le T([ \sqrt{n} ]) + 1\$\$$  for  $\$n \ge 1\$\$$ . Here [x] denotes the "floor" function, which rounds down to the nearest integer. (Note that the Master Method does not apply.)

O	\$\$O(\sqrt{n})\$\$
0	\$\$O(1)\$\$
0	\$\$O(\log n)\$\$
0	\$\$O(\log \log n)\$\$

5 questions unanswered

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