

VII. APPENDIX

A. Proof on Lemma 4.1 and Lemma 4.2

Lemma 4.1 $\forall u \in V, p \in H_0, 0 \leq p(u) \leq 1$, have the following relation:

$$\prod_{p \in H_0} (1 - p(u)) \geq 1 - \sum_{p \in H_0} p(u). \quad (6)$$

Proof: Let p_i be the distribution of the hash function h_i , then Equation (6) can be expressed as:

$$\prod_{i=0}^k (1 - p_i(u)) \geq 1 - \sum_{i=0}^k p_i(u). \quad (21)$$

We denote Equation (21) as Ψ . Next we use mathematical induction to prove Ψ , obviously it holds when $k = 0$, we assume that Ψ holds when $k = \alpha - 1$, then we have $\prod_{i=0}^{\alpha-1} (1 - p_i(u)) \geq 1 - \sum_{i=0}^{\alpha-1} p_i(u)$ and we can get:

$$\begin{aligned} \prod_{i=0}^{\alpha} (1 - p_i(u)) &= (1 - p_{\alpha}(u)) \prod_{i=0}^{\alpha-1} (1 - p_i(u)) \\ &= \prod_{i=0}^{\alpha-1} (1 - p_i(u)) - p_{\alpha}(u) \prod_{i=0}^{\alpha-1} (1 - p_i(u)) \\ &\geq 1 - \sum_{i=0}^{\alpha-1} p_i(u) - p_{\alpha}(u) \prod_{i=0}^{\alpha-1} (1 - p_i(u)) \\ &\geq 1 - \sum_{i=0}^{\alpha} p_i(u). \end{aligned} \quad (22)$$

Therefore, Ψ holds when $k = \alpha$, this completes the proof. ■

Lemma 4.2 $\forall 0 \leq x \leq 1$, Function $f(x) = \frac{|S| \cdot x}{(1-x)^{|S|}-1}$ is a convex function.

Proof: We rewrite the $f(x)$ as follows:

$$f(x) = \frac{|S| \cdot x(1-x)^{|S|}}{1 - (1-x)^{|S|}} = \frac{|S| \cdot (1-x)^{|S|}}{\sum_{i=0}^{|S|-1} (1-x)^i} \quad (23)$$

Let $\mu = 1 - x$ and $\theta = |S|$, so $f(\mu) = \frac{\theta \mu^{\theta}}{\sum_{i=0}^{\theta-1} \mu^i}$, and we can derive $f'(\mu)$ as follows:

$$f'(\mu) = \theta \frac{\sum_{i=\theta-1}^{2\theta-2} (2\theta-1-i) \mu^i}{(\sum_{i=0}^{\theta-1} \mu^i)^2} > 0 \quad (24)$$

Since $f'(x) = \frac{\delta f(\mu)}{\delta \mu} \frac{\delta \mu}{\delta x} = -f'(\mu) < 0$, then we can derive $f''(\mu)$ as follows:

$$\begin{aligned} f''(\mu) &= \frac{\theta}{(\sum_{i=0}^{\theta-1} \mu^i)^4} \left((\sum_{i=0}^{\theta-1} \mu^i)^2 \sum_{i=\theta-1}^{2\theta-2} i(2\theta-1-i) \mu^{i-1} \right. \\ &\quad \left. - 2 \sum_{i=0}^{\theta-1} \mu^i \sum_{i=1}^{\theta-1} i \mu^{i-1} \sum_{i=\theta-1}^{2\theta-2} (2\theta-1-i) \mu^i \right) \\ &= \frac{\theta}{(\sum_{i=0}^{\theta-1} \mu^i)^3} \left(\sum_{i=0}^{\theta-1} \mu^i \sum_{i=\theta-1}^{2\theta-2} i(2\theta-1-i) \mu^{i-1} \right. \\ &\quad \left. - 2 \sum_{i=0}^{\theta-1} i \mu^i \sum_{i=\theta-1}^{2\theta-2} (2\theta-1-i) \mu^{i-1} \right) \end{aligned} \quad (25)$$

Next, we compare $\sum_{i=0}^{\theta-1} i \mu^i$ with $\frac{\theta-1}{2} \sum_{i=0}^{\theta-1} \mu^i$,

$$\begin{aligned} &\frac{\theta-1}{2} \sum_{i=0}^{\theta-1} \mu^i - \sum_{i=0}^{\theta-1} i \mu^i \\ &= \sum_{i=0}^{\theta-1} \left(\frac{\theta-1}{2} - i \right) \mu^i \\ &= \sum_{i=0}^{\frac{\theta-1}{2}} \left(\frac{\theta-1}{2} - i \right) (\mu^i - \mu^{\theta-1-i}) \end{aligned} \quad (26)$$

Since $0 \leq \mu \leq 1$, we have $\sum_{i=0}^{\theta-1} i \mu^i < \frac{\theta-1}{2} \sum_{i=0}^{\theta-1} \mu^i$. According to Equation (25), we have:

$$\begin{aligned} f''(\mu) &> \frac{\theta}{(\sum_{i=0}^{\theta-1} \mu^i)^3} \left(\sum_{i=0}^{\theta-1} \mu^i \sum_{i=\theta-1}^{2\theta-2} i(2\theta-1-i) \mu^{i-1} \right. \\ &\quad \left. - (\theta-1) \sum_{i=0}^{\theta-1} \mu^i \sum_{i=\theta-1}^{2\theta-2} (2\theta-1-i) \mu^{i-1} \right) \\ &= \frac{\theta}{(\sum_{i=0}^{\theta-1} \mu^i)^2} \sum_{i=\theta-1}^{2\theta-2} (i - (\theta-1))(2\theta-1-i) \mu^{i-1} \end{aligned} \quad (27)$$

Therefore, $f''(\mu) > 0$ and $f''(x) = \frac{\delta^2 f(\mu)}{\delta^2 \mu} (\frac{\delta \mu}{\delta x})^2 + \frac{\delta f(\mu)}{\delta \mu} (\frac{\delta^2 \mu}{\delta^2 x}) = f''(\mu) > 0$. Since $f'(x) < 0$, $f''(x) > 0$, $f(x)$ is a convex function. ■

B. Analysis of P'_c

To simplify the analysis, we assume that each bit in Bloom filter is set to 0 with probability p_0 and 1 with probability $1-p_0$. Note that we do not consider the case of *cost exchange*.

When all buckets mapped by e_{sk} through all hash functions in H_c are *conflict after adjustment*, we cannot adjust the hash functions of e_{sk} , so we get

$$P'_c = 1 - \prod_{h \in H_c(e_{sk})} (1 - (1 - p_0^{k-1})^{\chi(h(e_{sk})))}, \quad (28)$$

where $\chi(i)$ represents the number of keys in the i^{th} bucket of Γ . Moreover, according to average value inequality, we have

$$\begin{aligned} 1 - P'_c &\leq \left(\frac{|H| - k - \sum_{h \in H_c} (1 - p_0^{k-1})^{\chi(h(e_{sk})))}{|H| - k} \right)^{|H| - k} \\ &\leq \left(1 - \frac{1}{|H| - k} \sum_{h \in H_c} (1 - p_0^{k-1})^{\chi(h(e_{sk})))} \right)^{|H| - k} \\ &\leq \left(1 - \prod_{h \in H_c} (1 - p_0^{k-1})^{\frac{\chi(h(e_{sk})))}{|H| - k}} \right)^{|H| - k}. \end{aligned} \quad (29)$$

It is easy to prove that: $\forall 0 < \alpha < 1, \beta \in \mathbb{N}, (1 - \alpha)^\beta < 1 - \alpha^\beta$, which is similar to Lemma 4.1, then we have

$$\begin{aligned} 1 - P'_c &< 1 - (1 - p_0^{k-1})^{\sum_{h \in H_c} \chi(h(e_{sk})))} \\ P'_c &> (1 - p_0^{k-1})^{\sum_{h \in H_c} \chi(h(e_{sk})))}. \end{aligned} \quad (30)$$

Since function $g''(x) = (1 - p_0^{k-1})^x$ is a convex function, by the Jensen inequality, we get

$$E(P'_c) > (1 - p_0^{k-1})^{E(\sum_{h \in H_c} \chi(h(e_{sk})))}. \quad (31)$$

Let $\psi = \sum_{h \in H_c} \chi(h(e_{sk}))$, and we assume that $\forall h \in H, e_{sk} \in S$, for a certain unit u in V , the probability that u is mapped by e_{sk} through h is only determined by $p(u)$, so we have

$$E(\psi) = E\left(\sum_{u=1}^m \sum_{p \in H_c} \chi(u)p(u)\right) = E\left(\sum_{u=1}^m \chi(u) \sum_{p \in H_c} p(u)\right), \quad (32)$$

where $\chi(u) = |O| \sum_{p' \in H_0} p'(u)$, for $\forall p_\alpha \in H_0, p_\gamma \in H_c, p_\alpha$ and p_γ are independent of each other, we have

$$\begin{aligned} E(\psi) &= \sum_{u=1}^m |O| E\left(\sum_{p \in H_0} p(u)\right) \cdot E\left(\sum_{p \in H_c} p(u)\right) \\ &< \sum_{u=1}^m \frac{|O|}{4} \left(\sum_{p \in H} E(p(u))\right)^2 = \frac{|O| \cdot |H|^2}{4m}. \end{aligned} \quad (33)$$

Since $0 < (1 - p_0^{k-1}) < 1$, then

$$E(P'_c) > (1 - p_0^{k-1})^{\frac{|O| \cdot |H|^2}{4m}}. \quad (34)$$