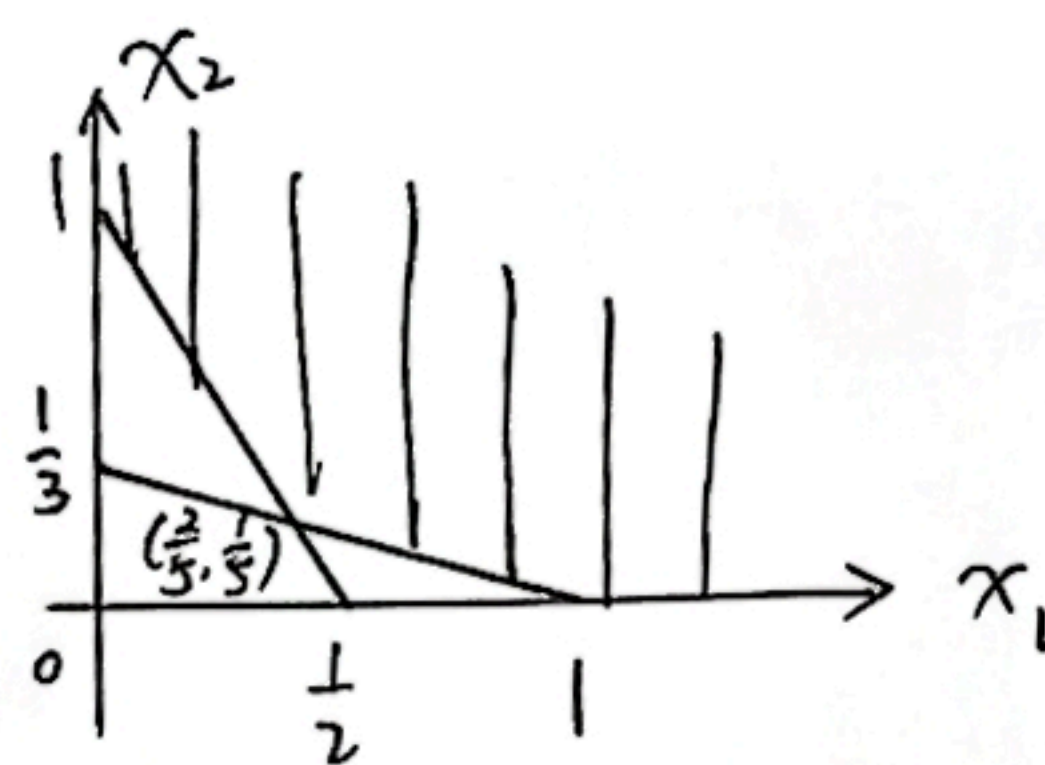


4.1 凸包的边界点为 $(0, 1)$, $(1, 0)$, $(\frac{2}{5}, \frac{1}{5})$, $(0, \infty)$, $(\infty, 0)$.

(a). $f_0(x_1, x_2) = x_1 + x_2$



①. $x_1 \in [0, \frac{2}{5}]$ 时, $x_2 \in [1-2x_1, +\infty)$.

要求 $f_0(x_1, x_2)$ 的最小值, 即 $x_2 = 1-2x_1$

$$f_0(x_1, x_2) = x_1 + 1-2x_1 = 1-x_1$$

当 $x_1 = \frac{2}{5}$ 时, $\min f_0(x_1, x_2) = \frac{3}{5}$.

②. $x_2 \in [0, \frac{1}{5}]$ 时, $x_1 \in [1-3x_2, +\infty)$.

要求 $f_0(x_1, x_2)$ 的最小值, 即 $x_1 = 1-3x_2$,

$$f_0(x_1, x_2) = 1-3x_2 + x_2 = 1-2x_2.$$

当 $x_2 = \frac{1}{5}$ 时, $\min f_0(x_1, x_2) = \frac{3}{5}$.

综上. $x^* = (\frac{2}{5}, \frac{1}{5})$. $\min f_0(x_1, x_2) = \frac{3}{5}$.

(b). 假设存在 (x_1, x_2) 使得 $f_0(x_1, x_2)$ 最小,
则必存在 (x_1+1, x_2+1) 使得 $f_0(x_1+1, x_2+1) < f_0(x_1, x_2)$
 $\therefore f_0(x_1, x_2)$ 不存在边界.

(c). $f_0(x_1, x_2) = x_1$

$$\because x_1 \geq 0.$$

$$\therefore f_0(x_1, x_2) \geq 0.$$

$$\therefore \min f_0(x_1, x_2) = 0.$$

$$\text{当 } x_1 = 0 \text{ 时, } x_2 \geq 0$$

$$\therefore x_{opt} = \{(0, x_2) | x_2 \geq 0\}.$$

(d). $f_0(x_1, x_2) = \max\{x_1, x_2\}$.

①. $x_1 \geq x_2$ 时. $x_1 = x_2$ 与 $2x_1 + x_2 = 1$ 的交点为 $(\frac{1}{3}, \frac{1}{3})$

$$\therefore f_0(x_1, x_2) = x_1$$

而此时 $x_1 \in [\frac{1}{3}, +\infty)$.

因此 $x^* = (\frac{1}{3}, \frac{1}{3})$. $\min f_0(x_1, x_2) = \frac{1}{3}$

②. $x_1 \leq x_2$ 时.

$$f_0(x_1, x_2) = x_2.$$

而此时 $x_2 \in [\frac{1}{3}, +\infty)$

$$\text{因此 } x^* = (\frac{1}{3}, \frac{1}{3}), \min f_0(x_1, x_2) = \frac{1}{3}$$

$$\text{综上. } x^* = (\frac{1}{3}, \frac{1}{3}), \min f_0(x_1, x_2) = \frac{1}{3}.$$

(e). $f_0(x_1, x_2) = x_1^2 + 9x_2^2$

①. 当 $x_1 \in [0, \frac{2}{5}]$ 时, $x_2 \in [1-2x_1, +\infty)$.

要求 $f_0(x_1, x_2)$ 的最小值, $x_2 = 1-2x_1$.

$$f_0(x_1, x_2) = x_1^2 + 9(1-2x_1)^2 \quad \text{求导得} \quad 74x_1 - 36 = 0. \quad x_1 = \frac{36}{74} > \frac{2}{5} \text{ (舍)}.$$

$$\text{当 } x_1 = \frac{2}{5} \text{ 时, } \min f_0(x_1, x_2) = \frac{13}{25} \quad \frac{1}{2} x_1 = \frac{2}{5} \text{ 时, } \min f_0(x_1, x_2) = \frac{13}{25}.$$

②. 当 $x_2 \in [0, \frac{1}{6}]$ 时, $x_1 \in [1-3x_2, +\infty)$.

要求 $f_0(x_1, x_2)$ 的最小值, $x_1 = 1-3x_2$.

$$f_0(x_1, x_2) = (1-3x_2)^2 + 9x_2^2 = 18x_2^2 - 6x_2 + 1.$$

$$\text{求导得} \quad 36x_2 - 6 = 0.$$

$$x_2 = \frac{1}{6}$$

$$x_1 = 1-3x_2 = \frac{1}{2}.$$

$$\text{此时 } \min f(x_1, x_2) = \frac{1}{2}$$

综上所述. $\frac{13}{25} > \frac{1}{2}$. 即选 ②. 为最优.

$$x^* = (\frac{1}{2}, \frac{1}{6}). \quad \min f_0(x_1, x_2) = \frac{1}{2}.$$

4.3. 欲证 $x^* = (1; \frac{1}{2}; -1)$ 为最优解, 等同于证明.

$$\nabla f_0(x)^T (y-x) \geq 0 \quad \text{对于所有 } y \in X$$

$$\nabla f_0(x) = x^T p + a^T$$

$$\text{当 } x^* \text{ 代入可得. } \nabla f(x^*) = (-1; 0; 2).$$

因此, 最优化的条件为 $-(y_1 - 1) + 2(y_3 + 1) \geq 0$.

$$\therefore y_1, y_2, y_3 \in [-1, 1].$$

$\therefore \nabla f(x^*)^T (y-x) \geq 0$ 必然成立, 证毕.

4.9. 设计变量 $y = Ax$, 则本题中问题可以改写为

$$\begin{array}{ll} \text{minimize} & c^T A^{-1} y \\ \text{subject to} & y \leq b. \end{array}$$

若 $A^{-T}c \leq 0$, 则最优解为 $y = b$, 则 $p^* = c^T A^{-1}b$,

反之, 则该问题无下界.

4.12. 题中所述可以表现为以下线性优化问题:

$$\begin{array}{ll} \text{minimize} & C = \sum_{i,j=1}^n c_{ij} x_{ij} \\ \text{subject to} & b_i + \sum_{j=1}^n x_{ij} - \sum_{j=1}^n x_{ji} = 0, \quad i=1, \dots, n. \\ & l_{ij} \leq x_{ij} \leq u_{ij}. \end{array}$$

4.13. 题中所述问题等价于.

$$\begin{aligned} & \text{minimize } c^T x \\ & \text{subject to } \bar{A}x + V|x| \leq b. \end{aligned}$$

其中 $|x| = (|x_1|, |x_2|, \dots, |x_n|)$.

上述问题反过来也等价于

$$\begin{aligned} & \text{minimize } c^T x \\ & \text{subject to } \bar{A}x + Vy \leq b \\ & \quad -y \leq x \leq y \end{aligned}$$

其中 $x \in \mathbb{R}^n$, $y \in \mathbb{R}^n$.

4.19. (a). 若题中问题为拟凸优化问题,
则可以等价为 $f_0(x) \leq \alpha$ 为凸.

也即 $\|Ax - b\|_1 - \alpha(c^T x + d) \leq 0$ 为凸.

显然, $\|Ax - b\|_1 - \alpha(c^T x + d) \leq 0$ 为一个凸约束.

(b). 由于 $\|x\|_\infty \leq 1$, $d > \|c\|_1$, 可以推导出 $c^T x + d > 0$.

定义 $y = \frac{x}{c^T x + d}$, $t = \frac{1}{c^T x + d}$.

则在凸问题中 y 和 t 是可行的. 其中

$$\|Ay - bt\|_1 = \frac{\|Ax - b\|_1}{c^T x + d}.$$

相反, 假设 y 和 t 对于凸问题 y 和 t 是可行的.

由于当 $t = 0$ 时, $y = 0$, 这与 $c^T y + dt = 1$ 矛盾, 因此, $t > 0$.

定义 $x = \frac{y}{t}$.

由于 $\|y\|_\infty \leq t$, $\therefore \|x\|_\infty \leq 1$, 且 $c^T x + d = \frac{1}{t}$. 因此,

$$\frac{\|Ax - b\|_1}{c^T x + d} = \|Ay - bt\|_1.$$

4.23. 题中所述的问题可以写为以下二次约束二次规划问题:

$$\begin{aligned} & \text{minimize} \quad \sum_{i=1}^m z_i^2 \\ & \text{subject to} \quad a_i^T x - b_i = y_i, \quad i = 1, \dots, m. \\ & \quad \quad \quad y_i^2 \leq z_i, \quad i = 1, \dots, m. \end{aligned}$$

4.33.

(a). 题中所述问题等同于几何规划问题:

$$\begin{aligned} & \text{minimize} \quad t \\ & \text{subject to} \quad \frac{p(x)}{t} \leq 1, \quad \frac{q(x)}{t} \leq 1. \end{aligned}$$

对 $x_i = e^{y_i}$ 进行对数变换.

(b). 题中所述问题等同于.

$$\begin{aligned} & \text{minimize} \quad \exp(t_1) + \exp(t_2). \\ & \text{subject to} \quad p(x) \leq t_1, \quad q(x) \leq t_2. \end{aligned}$$

对 $x_i = e^{y_i}$ 进行对数变换.

(c). 题中所述问题等同于

$$\begin{aligned} & \text{minimize} \quad t \\ & \text{subject to} \quad p(x) \leq t(r(x) - q(x)), \end{aligned}$$

和

$$\begin{aligned} & \text{minimize} \quad t \\ & \text{subject to} \quad \frac{(p(x)/t + q(x)/t)}{r(x)} \leq 1. \end{aligned}$$

这也是一个几何规划问题.