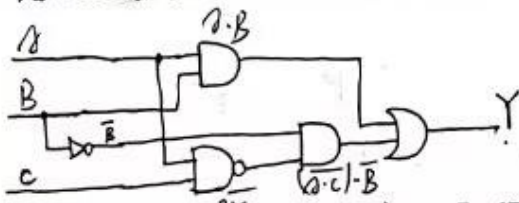


2.19.00180 题解

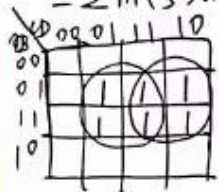
1. 用全逻辑电路实现以下函数: $Y = (A \cdot B) + (\overline{A \cdot C}) \cdot \overline{B}$.



2. 用卡诺图化简, 只用 NAND 门实现. $F = ABC + B\overline{C}D + \overline{A}BC$.

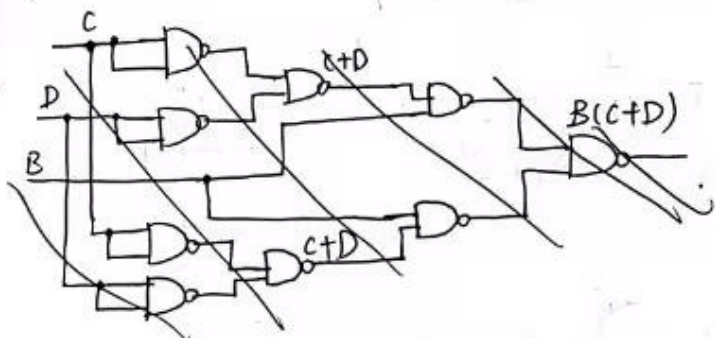
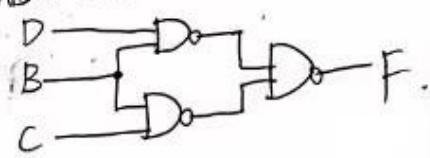
$$F = ABC + B\overline{C}D + \overline{A}BC = ABC(D + \overline{D}) + (\overline{A} + \overline{C})BCD + \overline{A}BC(D + \overline{D}) = ABCD + ABC\overline{D} + \overline{A}BCD + \overline{A}BC\overline{D} + \overline{A}BCD + \overline{A}BC\overline{D}$$

$$= \sum m(5, 6, 7, 13, 14, 15).$$



$$F = BD + BC.$$

$$F = B(C + D).$$

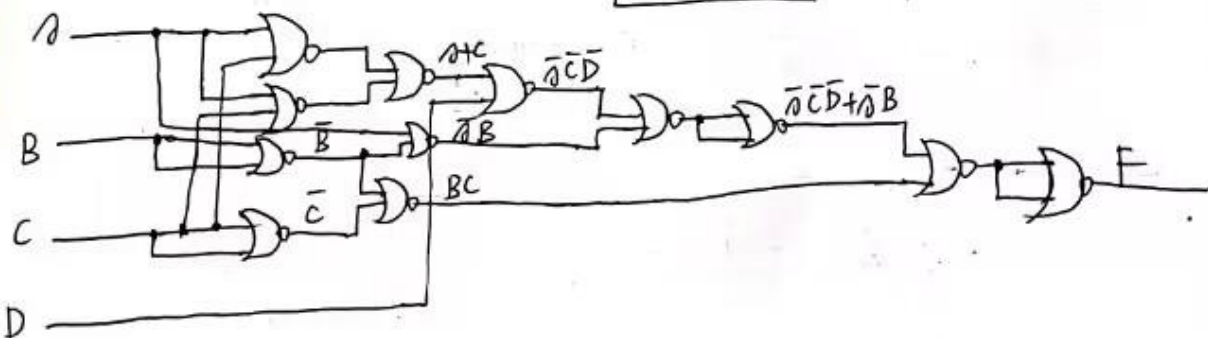
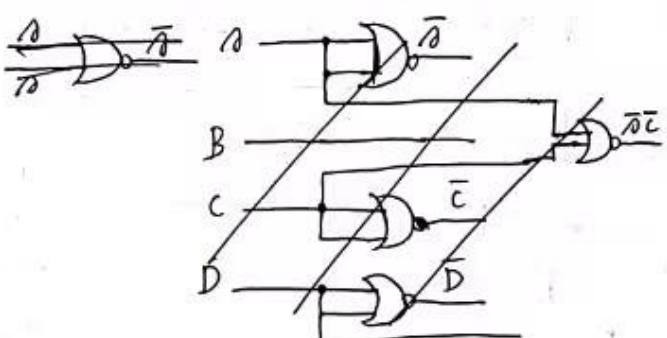


3. 用卡诺图化简以下函数, 只用 NOR 门实现.

$$F(A, B, C, D) = \prod M(1, 2, 3, 8, 9, 10, 11, 14) \cdot d(7, 15) = \sum m(0, 4, 5, 6, 12, 13) + d(7, 15).$$



$$F = \overline{A} \overline{C} \overline{D} + BC + \overline{A} B = \overline{A} \overline{C} \overline{D} + B \cdot C + \overline{A} \cdot B$$



4. 设计组合电路，功能为统计三个输入A, B, C中“1”出现的次数，写出表达式及电路图。

Handwritten solutions for problem 4, including circuit diagrams, truth tables, and logic expressions.

Circuit Diagrams:

- Initial attempt: A circuit with three AND gates and an OR gate, labeled "Count".
- Revised circuit: A circuit with three AND gates (Count1, Count2, Count3) and an OR gate, labeled "F".
- Final circuit: A circuit with three AND gates (Count1, Count2, Count3) and an OR gate, labeled "F".

Truth Tables:

Truth Table 1 (Crossed out):

A	B	C	Count1	Count2
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	0	0
1	0	1	0	0
1	1	0	0	0
1	1	1	0	0

Truth Table 2:

A	B	C	Count1	Count2	Count3
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	0	1
0	1	1	0	1	0
1	0	0	0	0	1
1	0	1	0	1	0
1	1	0	1	0	0
1	1	1	1	1	1

Logic Expressions:

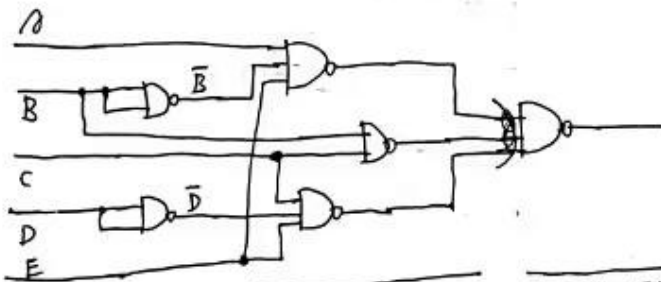
Final Count = Count1 + Count2 + Count3 (7为普通加法)

∴ 次数 = Count3 + 2 × (Count1 + Count2) (+与×为普通意义上的加与乘)

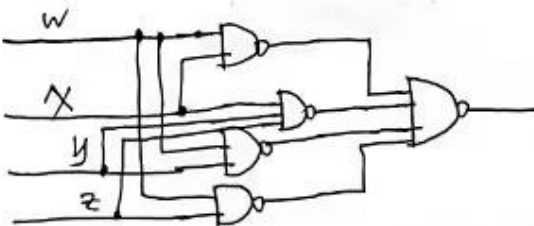
5. 以多级与非门原图:

① $(\overline{A}\overline{B} + \overline{C}\overline{D})E + B(\overline{A} + B)$

$$\begin{aligned} (\overline{A}\overline{B} + \overline{C}\overline{D})E + B(\overline{A} + B) &= \overline{A}\overline{B}E + \overline{C}\overline{D}E + B\overline{A} + B \\ &= (\overline{A} + B + \overline{E}) \cdot (\overline{C} + D + \overline{E}) \cdot (\overline{A} + \overline{B} + \overline{C}) \cdot (\overline{B} + \overline{D}) \\ &= (\overline{A} \cdot \overline{B} \cdot \overline{E}) \cdot (\overline{C} \cdot \overline{D} \cdot \overline{E}) \cdot (\overline{A} \cdot \overline{B} \cdot \overline{C}) \cdot (\overline{B} \cdot \overline{D}) \end{aligned}$$



② $w(x+y+z) + xyx = wx + wy + wz + xyx = (\overline{w} + x) \cdot (\overline{w} + y) \cdot (\overline{w} + z) \cdot (\overline{x} + y \cdot \overline{z}) = (\overline{w} \cdot x) \cdot (\overline{w} \cdot y) \cdot (\overline{w} \cdot z) \cdot (\overline{x} + y \cdot \overline{z})$

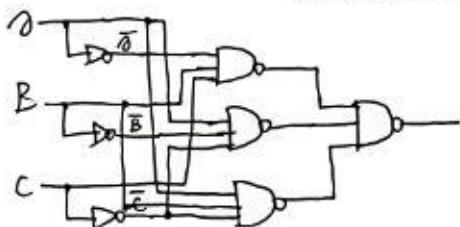


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6. 0~9, $L_1, L_2, L_3, L_4, L_5, L_6, L_7$, 写值表.

输入	L_1	L_2	L_3	L_4	L_5	L_6	L_7
0	1	0	1	1	1	1	1
1	0	0	0	0	0	1	1
2	1	1	1	0	1	1	0
3	1	1	1	0	0	1	1
4	0	1	0	1	0	1	1
5	1	1	1	1	0	0	1
6	1	1	1	1	1	0	1
7	1	0	0	0	0	1	1
8	1	1	1	1	1	1	1
9	1	1	1	1	0	1	1

7. 奇/偶数个 1, 3 输入. 奇数个 1 输出 0, 反之亦然.

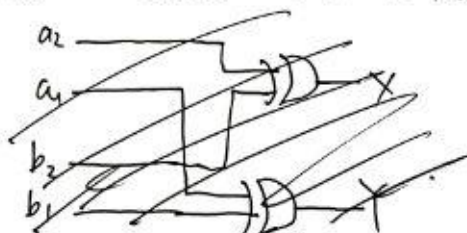
$$F = \overline{A}BC + A\overline{B}C + AB\overline{C} = \overline{(ABC + \overline{A}BC + \overline{A}B\overline{C})}$$



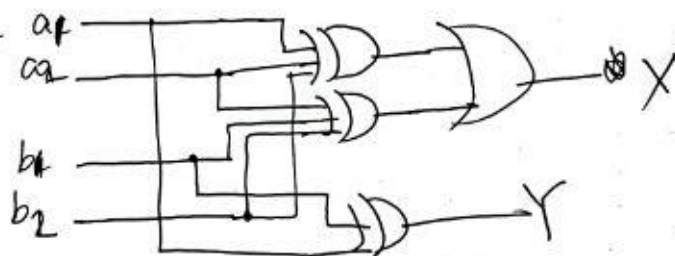
(亦即异或门)

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

8. $A(a_2 a_1) B(b_2 b_1)$, 求 $|A-B|$ 电路.



$|A-B| = XY$ 表示为 XY 这个数 (2).

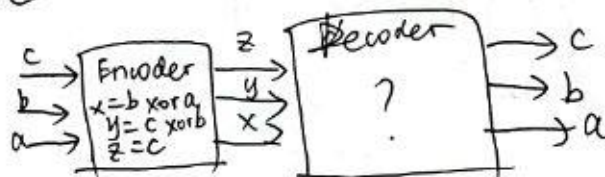


9. 证明 n 输入或门可用 $n-1$ 个 2 输入或门实现, 说明对与或非门是否仍适用.

证明: $\overline{(A+B+C+D+\dots+X)} = \overline{(A+B+C+D+\dots+X)}$, $\therefore \square$.

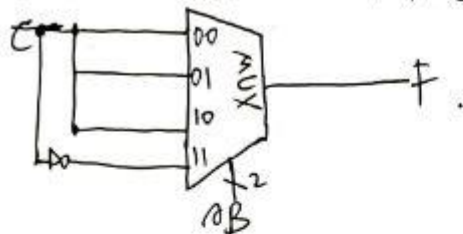
不适用. $\because \overline{A+B+C} \neq \overline{A+B+C}$.

10. 设计解码器, $(z, y, x) \rightarrow (c, b, a)$.

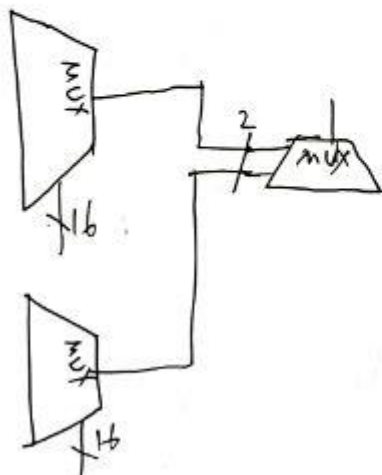


$$\begin{cases} c = z \\ b = y \text{ xor } z \\ a = x \text{ xor } (y \text{ xor } z) \end{cases}$$

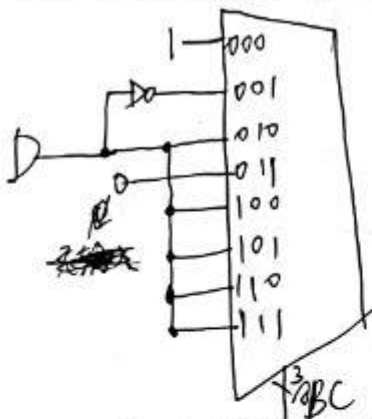
11. 基于4-路选择器实现 $F(A, B, C) = \sum m(0, 3, 5, 6)$.



12. 用两个16-路和一个2-路选择器设计22-路选择器.



13. 基于8-路选择器实现 $Y(A, B, C, D) = \sum m(0, 1, 2, 5, 9, 11, 13, 15)$.



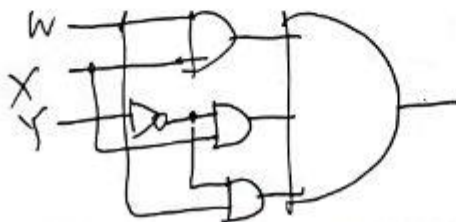
14. 针对下列表达式, 找出对应2级 AND-OR 或 OR-AND 的所有静态冒险, 设计能实现同样逻辑的无冒险电路.

(1) $F = W \cdot X + W' \cdot Y'$ (2) $F = W \cdot Y + W' \cdot Z' + X \cdot Y' \cdot Z$ (3) $F = (W + X + Y') \cdot (X' + Z')$ (4) $F = (W + Y + Z') \cdot (W + X + Y + Z) \cdot (X' + Y') \cdot (X + Z)$.

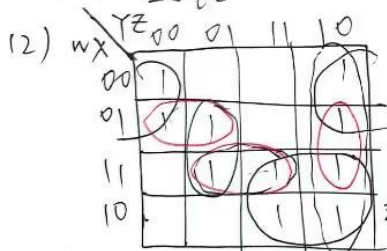
(1)

XY	00	01	11	10
W				
0	1		1	
1			1	1

$F = W \cdot X + W' \cdot Y' + X \cdot Y'$

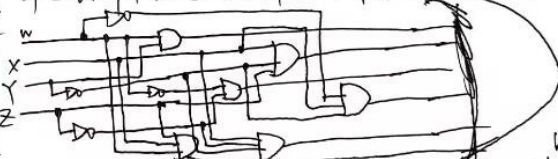


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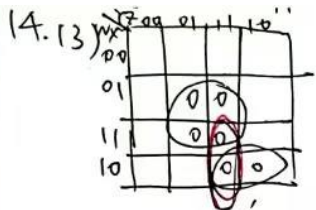


$$F = wY + w'z' + xY'z + w'xY + wXz + xYz'$$

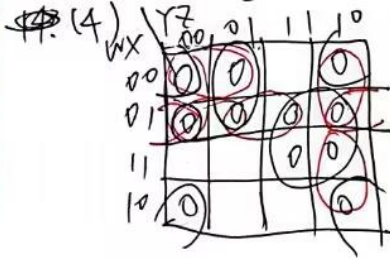
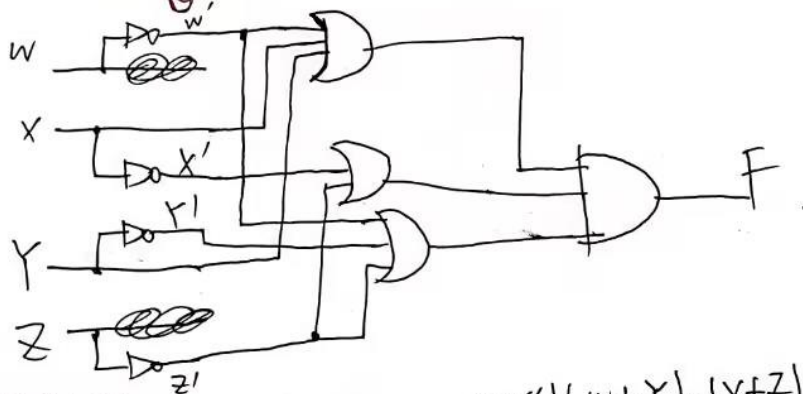
$$F = wY + w'z' + xY'z + Yz' + wXz$$



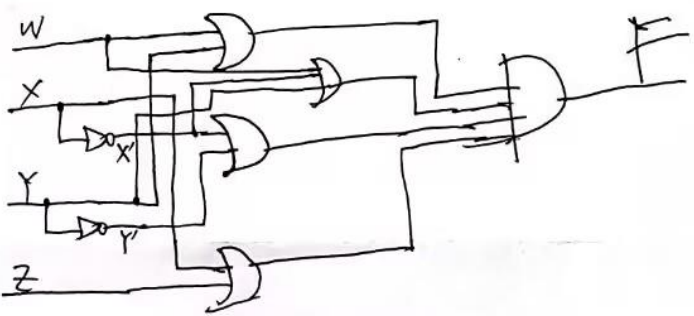
$$F = (w' + x + y')(x' + z')(w' + y' + z')$$



$$F = (w' + x + y')(x' + z')(w' + y' + z')$$

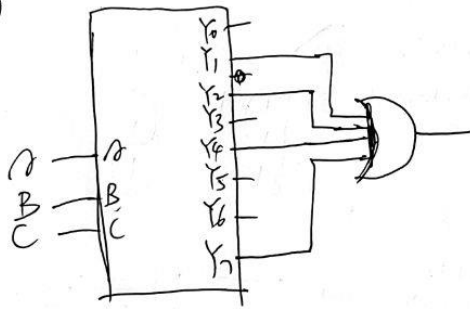


$$F = (x' + y') \cdot (w + y) \cdot (x + z) \cdot (w + x' + y)$$



15. 借助一个3-8译码器实现 (1) $S(A, B, C) = \sum m(1, 2, 4, 7)$ (2) $C(A, B, C) = \sum m(3, 5, 6, 7)$, 并说明作用

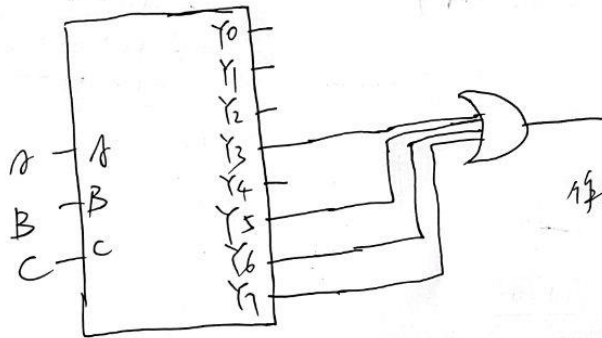
(1)



作用:

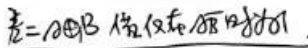
有奇数个1则输出1,
有偶数个1则输出0.

(2)



作用: 1的个数大于等于2
则输出1.

A	B	差	借位
0	0	0	0
0	1	1	1
1	0	1	0
1	1	0	0



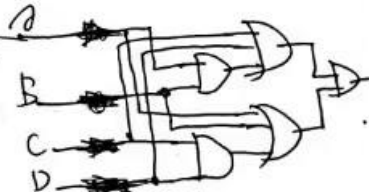
结果输出 = 数列数 ~~ABC~~

①求最大项表达式, ②使用三级门(NAND-OR-NAND)实现, 要求门最少

① $f(a, b, c, d) = TTM(0, 1, 2, 4, 8)$.

② $f = \sum m(3, 5, 6, 7, 9, 10, 11, 12, 13, 14, 15)$.

$$F = CD + AB + BC + AD + AC + BD$$



20. 构造一个5-32译码器, 使用4个带使能端的3-8译码器和一个2-4译码器

