# Semantic Theory Week 6 – Event semantics

Noortje Venhuizen Harm Brouwer

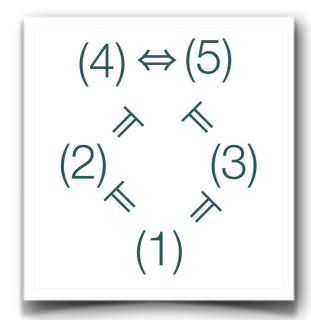
Universität des Saarlandes

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## A problem with verbs and adjuncts

(1) The gardener killed the baron

- $\rightarrow kill_1(g',b')$   $kill_1::\langle e,\langle e,t\rangle \rangle$
- (2) The gardener killed the baron in the park  $\mapsto \text{kill}_2(g',b',p')$   $\text{kill}_2::\langle e,\langle e,\langle e,t\rangle\rangle$
- (3) The gardener killed the baron at midnight  $\mapsto \text{kill}_3(g',b',m')$  kill<sub>3</sub> ::  $\langle e,\langle e,\langle e,t\rangle \rangle$
- (4) The gardener killed the baron at midnight in the park → kill<sub>4</sub>(g',b',m',p') kill<sub>4</sub> :: ....
- (5) The gardener killed the baron in the park at midnight → kill<sub>5</sub>(g',b',p',m') kill<sub>5</sub> :: ....

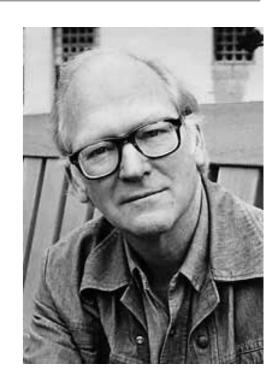


Q: How to explain the systematic logical entailment relations between the different uses of "kill"?

#### Davidson's solution: verbs introduce events.

Verbs expressing events have an additional event argument, which is not realised at linguistic surface:

• kill  $\mapsto \lambda y \lambda x \lambda e(kill'(e,x,y)) :: \langle e,\langle e,\langle e,t \rangle \rangle \rangle$  arity = n+1



Sentences denote sets of events:

•  $\lambda y \lambda x \lambda e(kill'(e,x,y))(b')(g') \Rightarrow^{\beta} \lambda e(kill'(e,g',b')) :: \langle e,t \rangle$ 

Existential closure turns sets of events into truth conditions

- $\lambda P \exists e(P(e)) :: \langle \langle e, t \rangle, t \rangle$
- $\lambda P \exists e(P(e))(\lambda e(kill'(e,g',b'))) \Rightarrow^{\beta} \exists e(kill'(e,g',b')) :: t$

## Davisonian events and adjuncts

Adjuncts express two-place relations between events and the respective "circumstantial information": time, location, ...

- at midnight  $\mapsto \lambda P \lambda e(P(e) \land time(e,m')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$
- in the park  $\mapsto \lambda P \lambda e(P(e) \land Iocation(e,p')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$

The gardener killed the baron at midnight in the park

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\Rightarrow \exists e \; (kill(e, g', b') \land time(e, m) \land location(e, p')) \} \vDash \exists e \; (kill(e, g', b') \land time(e, m')) \\ \Leftrightarrow \exists e \; (kill(e, g', b') \land location(e, p) \land time(e, m')) \} \vDash \exists e \; (kill(e, g', b') \land location(e, p')) \\ \vDash \exists e \; (kill(e, g', b') \land location(e, p')) \}
```

# Compositional derivation of event-semantic representations

#### the gardener killed the baron

```
\lambda x_e \lambda y_e \lambda e_e [\text{ kill(e, y, x) ](b')(g')} \Rightarrow^{\beta} \lambda e [\text{ kill(e, g', b') ]}
... at midnight
```

 $\lambda F_{\langle e,t\rangle} \lambda e_e$  [ F(e)  $\wedge$  time(e, m') ]( $\lambda e_1$  [ kill(e<sub>1</sub>, g', b') ])  $\Rightarrow^{\beta} \lambda e$  [ kill(e, g, b)  $\wedge$  time(e, m') ]

#### ... in the park

```
\lambda F_{\langle e,t\rangle} \lambda e_e [F(e) \wedge location(e, p')] (\lambda e_2 [kill(e<sub>2</sub>, g', b')\wedgetime(e<sub>2</sub>, m')]) \Rightarrow^{\beta} \lambda e [kill(e, g', b') \wedge time(e, m') \wedge location(e, p')]
```

#### Existential closure

 $\lambda P_{(e,t)} \exists e(P(e))(\lambda e'(K \land T \land L) \Rightarrow \beta \exists e [kill(e, g', b') \land time(e, m') \land location(e, p')]$ 

#### Model structures with events

To interpret events, we need enriched ontological information

Ontology: The area of philosophy identifying and describing the basic "categories of being" and their relations.

A model structure with events is a triple  $M = \langle U, E, V \rangle$ , where

- U is a set of "standard individuals" or "objects"
- E is a set of events
- U  $\cap$  E =  $\emptyset$ ,
- V is an interpretation function like in first order logic

# Sorted (first-order) logic

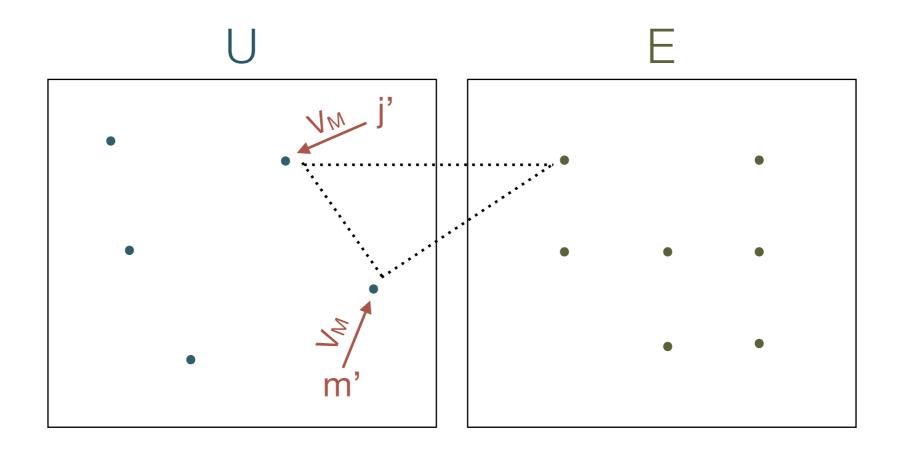
A variable assignment g assigns individuals (of the correct sortspecific domain) to variables:

- $g(x) \in U$  for  $x \in VAR_U$   $VAR_U = \{x, y, z, ..., x_1, x_2, ...\}$  (Object variables)
- $g(e) \in E$  for  $e \in VAR_E$   $VAR_E = \{ e, e', e'', ..., e_1, e_2, ... \}$  (Event variables)

Quantification ranges over sort-specific domains:

- $[\![ \exists x \Phi ]\!]^{M,g} = 1$  iff there is an  $a \in U$  such that  $[\![ \Phi ]\!]^{M,g[x/a]} = 1$
- $[\exists e \ \Phi]^{M,g} = 1$  iff there is an  $a \in E$  such that  $[\![\Phi]^{M,g[e/a]} = 1$
- (universal quantification analogous)

#### Interpreting events



## Advantages of Davidsonian events

- ☑ Intuitive representation and semantic construction for adjuncts
- Uniform treatment of verb complements
- Uniform treatment of adjuncts and post-nominal modifiers
- Coherent treatment of tense information
- Highly compatible with analysis of semantic roles

# Uniform treatment of verb complements

#### (1) Bill saw an elephant

$$\rightarrow$$
 3e 3x (see(e, b', x)  $\land$  elephant(x))

see :: 
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

(2) Bill saw an accident

see :: 
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

(3) Bill saw the children play

$$\rightarrow \exists e \exists e' (see(e, b, e') \land play(e', the-children))$$

see :: 
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

# Uniform treatment of adjuncts and post-nominal modifiers

Treatment of adjuncts as predicate modifiers, analogous to attributive adjectives:

- red  $\mapsto \lambda F \lambda x [F(x) \land red^*(x)]$   $\langle \langle e, t \rangle, \langle e, t \rangle \rangle$
- in the park  $\mapsto \lambda F \lambda e [F(e) \land location(e, park)] \langle \langle e, t \rangle, \langle e, t \rangle \rangle$
- (1) The murder in the park...
- $\rightarrow \lambda F\lambda e[F(e) \land location(e, park)] (\lambda e_1 [murder(e_1)])$
- (2) The fountain in the park ....
- $\rightarrow \lambda F \lambda x [F(x) \land location(x, park)] (\lambda y [fountain(y)])$

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## Classical Tense Logic

John walks walk(john)

John walked P(walk(john))

John will walk F(walk(john))

Syntax like in first-order logic, plus

Φ has always been the case

Φ is always going to be the case

 if Φ is a well-formed formula, then PΦ, FΦ, HΦ, GΦ are also well-formed formulae.

Φ happened in the past

Φ will happen in the future

# Classical Tense Logic (cont.)

Tense model structures are quadruples  $M = \langle U, T, \langle V \rangle$  where

- U is a non-empty set of individuals (the "universe")
- T is a non-empty sets of points in time
- $U \cap T = \emptyset$
- < is a linear order on T</li>
- V is a value assignment function, which assigns to every non-logical constant α a function from T to appropriate denotations of α

 $[P\Phi]^{M, t, g} = 1$  iff there is a t' < t such that  $[\Phi]^{M, t', g} = 1$ 

 $\llbracket F\Phi \rrbracket^{M, t, g} = 1$  iff there is a t' > t such that  $\llbracket \Phi \rrbracket^{M, t', g} = 1$ 

# Temporal Relations and Events

- (1) The door opened, and Mary entered the room.
- (2) John arrived. Then Mary left.
- (3) Mary left, before John arrived.
- (4) John arrived. Mary had left already.

Q: How to formalize temporal relations between events?

# Temporal Event Structure

A model structure with events and temporal precedence is defined as  $M = \langle U, E, \langle e_u, V \rangle$ , where

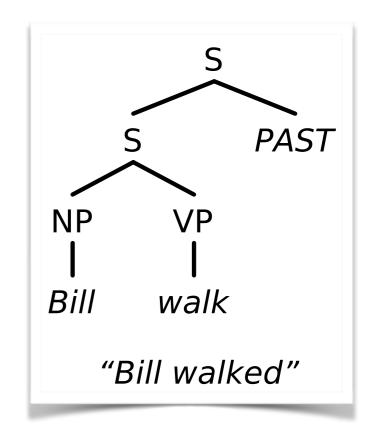
- Un  $E = \emptyset$ ,
- < ⊆ E×E is an asymmetric relation (temporal precedence)</li>
- $e_u \in E$  is the utterance event
- V is an interpretation function like in standard FOL
- · Overlapping events: e · e' iff neither e < e' nor e' < e

#### Tense in Semantic Construction

We can represent inflection as an abstract tense operator reflecting the temporal location of the reported event relative to the utterance event.

PAST 
$$\mapsto \lambda P.\exists e [P(e) \land e < e_u] : \langle\langle e, t \rangle, t \rangle$$

PRES 
$$\mapsto \lambda P. \exists e [P(e) \land e \cdot e_u] : \langle \langle e, t \rangle, t \rangle$$



#### Tense in Semantic Construction

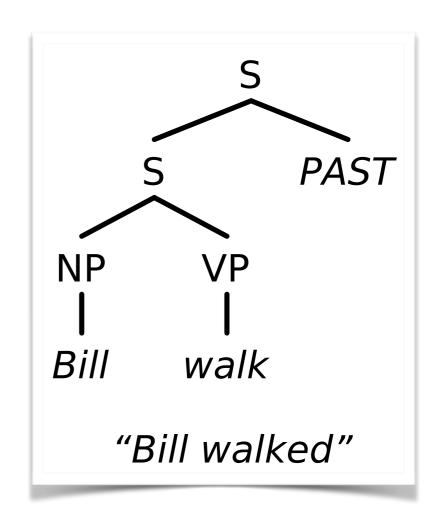
Standard function application results in integration of temporal information and binding of the event variable (i.e., replacing E-CLOS):

- walk  $\mapsto \lambda x \lambda e$  [walk(e, x)]
- Bill walk  $\mapsto \lambda x \lambda e$  [walk(e, x)](b')  $\Rightarrow^{\beta} \lambda e$  [walk(e, b')]
- Bill walk PAST

   → λΕ ∃e [E(e) ∧ e < e<sub>u</sub>](λe' [walk(e', b)])

   ⇒β ∃e [λe' [walk(e', b)](e) ∧ e < e<sub>u</sub>]

   ⇒β ∃e [walk(e, b) ∧ e < e<sub>u</sub>]

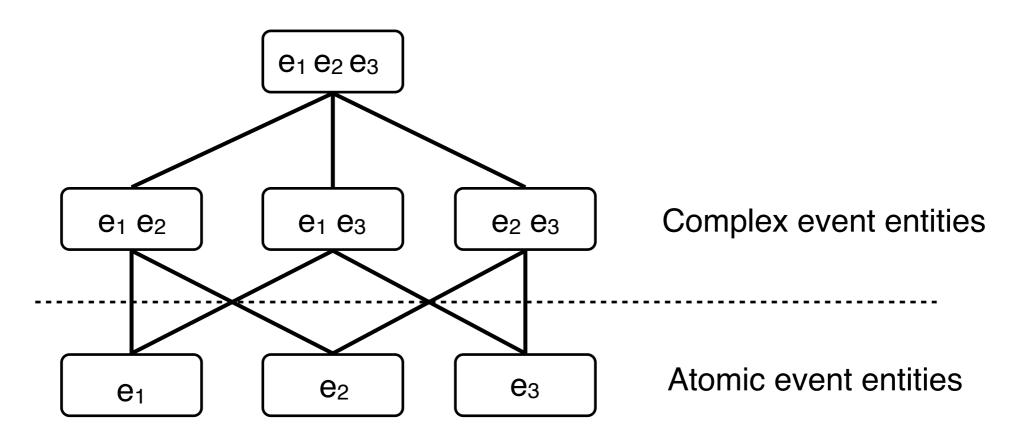


#### **Event Structure**

Observation: Events are generally constructs that consist of various (temporally ordered) sub-events

• E.g., "scripts": visit a restaurant or shopping in the supermarket

Idea: Induce structure into events universe



#### Lattices and Semi-lattices

A **partial order** is a structure  $\langle A, \leq \rangle$  where  $\leq$  is a reflexive, transitive, and antisymmetric relation over A.

- The **join** of a and  $b \in A$  (Notation:  $a \sqcup b$ ) is the lowest upper bound for a and b.
- The **meet** of a and  $b \in A$  (Notation:  $a \sqcap b$ ) is the highest lower bound for a and b.

A **lattice** is a partial order  $\langle A, \leq \rangle$  that is closed under meet and join.

A join semi-lattice is a partial order  $\langle A, \leq \rangle$  that is closed under join

#### Model Structure with Sub-Events

We can change the structure of the events universe to represent sub-event relations:  $M = \langle U, \langle E, \leq_e \rangle, \langle e_u, V \rangle$ , where:

- U  $\cap$  E =  $\emptyset$ ,
- < ⊆ E×E is an asymmetric relation (temporal precedence)</li>
- $e_u \in E$  is the utterance event
- ⟨E, ≤e⟩ is a join semi-lattice
- V is an interpretation function

# Model Structure with Sub-Events (cont.)

The model structure  $M = \langle U, \langle E, \leq_e \rangle, <, e_u, V \rangle$  must observe some additional constraints on < and  $\leq_e$ , for instance:

- If  $e_1 < e_2$  and  $e_1' \le_e e_1$  and  $e_2' \le_e e_2$ , then  $e_1' < e_2'$
- If  $e_1' \circ e_2'$  and  $e_1' \leq_e e_1$  and  $e_2' \leq_e e_2$ , then  $e_1 \circ e_2$

Sidenote: We could introduce a similar structuring of the universe of entities in order to capture *plurality* and other *composite entities* 

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#### Links

- Overview paper: Lasersohn (2012) Event-Based Semantics: <a href="https://semanticsarchive.net/Archive/jFhNWM2M/">https://semanticsarchive.net/Archive/jFhNWM2M/</a> eventbasedsemantics.pdf
- PropBank: <a href="http://propbank.github.io/">http://propbank.github.io/</a>
- FrameNet: <a href="https://framenet.icsi.berkeley.edu/fndrupal/">https://framenet.icsi.berkeley.edu/fndrupal/</a>