Semantic Theory Week 7 – Event semantics

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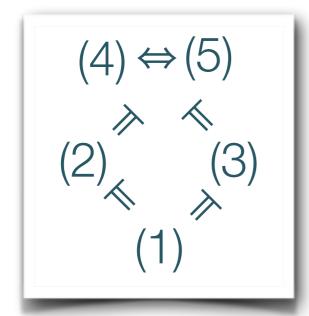
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A problem with verbs and adjuncts

(1) The gardener killed the baron

- $\rightarrow kill_1(g',b')$ $kill_1::\langle e,\langle e,t\rangle \rangle$
- (2) The gardener killed the baron in the park $\mapsto \text{kill}_2(g',b',p')$ $\text{kill}_2::\langle e,\langle e,\langle e,t\rangle\rangle$
- (3) The gardener killed the baron at midnight $\mapsto kill_3(g',b',m')$ $kill_3::\langle e,\langle e,\langle e,t\rangle \rangle$
- (4) The gardener killed the baron at midnight in the park → kill₄(g',b',m',p') kill₄ ::
- (5) The gardener killed the baron in the park at midnight → kill₅(g',b',p',m') kill₅ ::

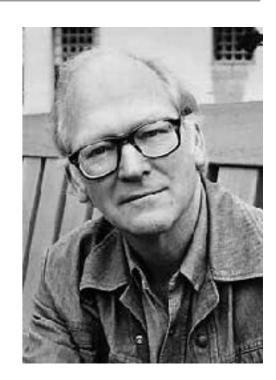


Q: How to explain the systematic logical entailment relations between the different uses of "kill"?

Davidson's solution: verbs introduce events.

Verbs expressing events have an additional event argument, which is not realised at linguistic surface:

• kill $\mapsto \lambda y \lambda x \lambda e(kill'(e,x,y)) :: \langle e,\langle e,\langle e,t \rangle \rangle \rangle$ arity = n+1



Sentences denote sets of events:

• $\lambda y \lambda x \lambda e(kill'(e,x,y))(b')(g') \Rightarrow^{\beta} \lambda e(kill'(e,g',b')) :: \langle e,t \rangle$

Existential closure turns sets of events into truth conditions

- $\lambda P \exists e(P(e)) :: \langle \langle e, t \rangle, t \rangle$
- $\lambda P \exists e(P(e))(\lambda e(kill'(e,g',b'))) \Rightarrow^{\beta} \exists e(kill'(e,g',b')) :: t$

Davisonian events and adjuncts

Adjuncts express two-place relations between events and the respective "circumstantial information": time, location, ...

- at midnight $\mapsto \lambda P \lambda e(P(e) \land time(e,m')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$
- in the park $\mapsto \lambda P \lambda e(P(e) \land Iocation(e,p')) :: \langle \langle e,t \rangle, \langle e,t \rangle \rangle$

The gardener killed the baron at midnight in the park

```
\Rightarrow \exists e \; (kill(e, g', b') \land time(e, m) \land location(e, p')) \} \vDash \exists e \; (kill(e, g', b') \land time(e, m')) \\ \Leftrightarrow \exists e \; (kill(e, g', b') \land location(e, p) \land time(e, m')) \} \vDash \exists e \; (kill(e, g', b') \land location(e, p')) \\ \vDash \exists e \; (kill(e, g', b') \land location(e, p')) \}
```

Compositional derivation of event-semantic representations

the gardener killed the baron

```
\lambda x_e \lambda y_e \lambda e_e [\text{kill}(e, y, x)](b')(g') \Rightarrow^{\beta} \lambda e [\text{kill}(e, g', b')]
... at midnight
```

 $\lambda F_{(e,t)}\lambda e_e$ [F(e) \wedge time(e, m')](λe_1 [kill(e₁, g', b')]) $\Rightarrow^{\beta} \lambda e$ [kill(e, g, b) \wedge time(e, m')]

... in the park

```
\lambda F_{\langle e,t\rangle} \lambda e_e [F(e) \wedge location(e, p')] (\lambda e_2 [kill(e<sub>2</sub>, g', b')\wedgetime(e<sub>2</sub>, m')]) \Rightarrow^{\beta} \lambda e [kill(e, g', b') \wedge time(e, m') \wedge location(e, p')]
```

Existential closure

 $\lambda P_{(e,t)} \exists e(P(e))(\lambda e'(K \land T \land L) \Rightarrow \beta \exists e [kill(e, g', b') \land time(e, m') \land location(e, p')]$

Model structures with events

To interpret events, we need enriched ontological information

Ontology: The area of philosophy identifying and describing the basic "categories of being" and their relations.

A model structure with events is a triple $M = \langle U, E, V \rangle$, where

- U is a set of "standard individuals" or "objects"
- E is a set of events
- U \cap E = \emptyset ,
- V is an interpretation function like in first order logic

Sorted (first-order) logic

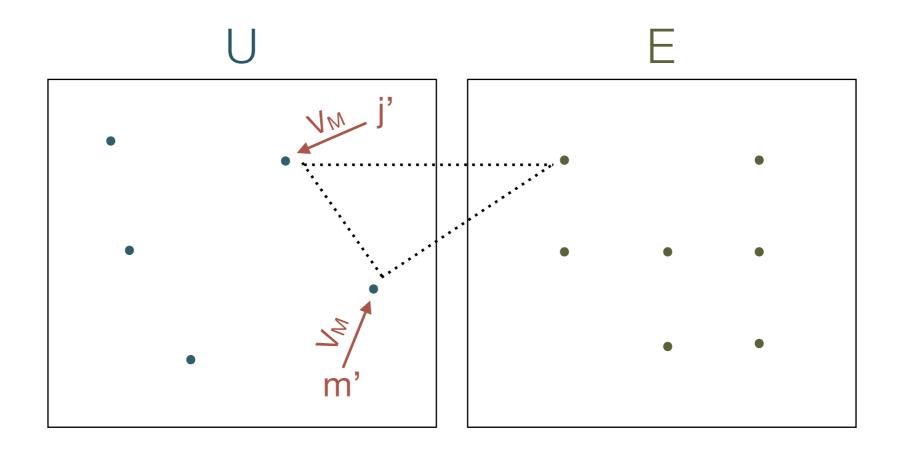
A variable assignment g assigns individuals (of the correct sortspecific domain) to variables:

- $g(x) \in U$ for $x \in VAR_U$ $VAR_U = \{x, y, z, ..., x_1, x_2, ...\}$ (Object variables)
- $g(e) \in E$ for $e \in VAR_E$ $VAR_E = \{ e, e', e'', ..., e_1, e_2, ... \}$ (Event variables)

Quantification ranges over sort-specific domains:

- $[\![\exists x \Phi]\!]^{M,g} = 1$ iff there is an $a \in U$ such that $[\![\Phi]\!]^{M,g[x/a]} = 1$
- $[\exists e \ \Phi]^{M,g} = 1$ iff there is an $a \in E$ such that $[\![\Phi]^{M,g[e/a]} = 1$
- (universal quantification analogous)

Interpreting events



Advantages of Davidsonian events

- ☑ Intuitive representation and semantic construction for adjuncts
- Uniform treatment of verb complements
- Uniform treatment of adjuncts and post-nominal modifiers
- Coherent treatment of tense information
- Highly compatible with analysis of semantic roles

Uniform treatment of verb complements

(1) Bill saw an elephant

$$\rightarrow$$
 3e 3x (see(e, b', x) \land elephant(x))

see ::
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

(2) Bill saw an accident

see ::
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

(3) Bill saw the children play

$$\rightarrow$$
 3e 3e' (see(e, b, e') \land play(e', the-children))

see ::
$$\langle e, \langle e, \langle e, t \rangle \rangle$$

Uniform treatment of adjuncts and post-nominal modifiers

Treatment of adjuncts as predicate modifiers, analogous to attributive adjectives:

```
• red \mapsto \lambda F \lambda x [F(x) \land red^*(x)] \langle \langle e, t \rangle, \langle e, t \rangle \rangle
```

- in the park $\mapsto \lambda F \lambda e [F(e) \land location(e, park)] \langle \langle e, t \rangle, \langle e, t \rangle \rangle$
- (1) The murder in the park...
- $\rightarrow \lambda F\lambda e[F(e) \land location(e, park)] (\lambda e_1 [murder(e_1)])$
- (2) The fountain in the park
- $\rightarrow \lambda F \lambda x [F(x) \land location(x, park)] (\lambda y [fountain(y)])$

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Classical Tense Logic

John walks walk(john)

John walked P(walk(john))

John will walk F(walk(john))

Syntax like in first-order logic, plus

Φ has always been the case

Φ is always going to be the case

 if Φ is a well-formed formula, then PΦ, FΦ, HΦ, GΦ are also well-formed formulae.

Φ happened in the past

Φ will happen in the future

Classical Tense Logic (cont.)

Tense model structures are quadruples $M = \langle U, T, \langle V \rangle$ where

- U is a non-empty set of individuals (the "universe")
- T is a non-empty sets of points in time
- $U \cap T = \emptyset$
- < is a linear order on T
- V is a value assignment function, which assigns to every non-logical constant α a function from T to appropriate denotations of α

 $[P\Phi]^{M, t, g} = 1$ iff there is a t' < t such that $[\Phi]^{M, t', g} = 1$

 $\llbracket F\Phi \rrbracket^{M, t, g} = 1$ iff there is a t' > t such that $\llbracket \Phi \rrbracket^{M, t', g} = 1$

Temporal Relations and Events

- (1) The door opened, and Mary entered the room.
- (2) John arrived. Then Mary left.
- (3) Mary left, before John arrived.
- (4) John arrived. Mary had left already.

Q: How to formalize temporal relations between events?

Temporal Event Structure

A model structure with events and temporal precedence is defined as $M = \langle U, E, \langle e_u, V \rangle$, where

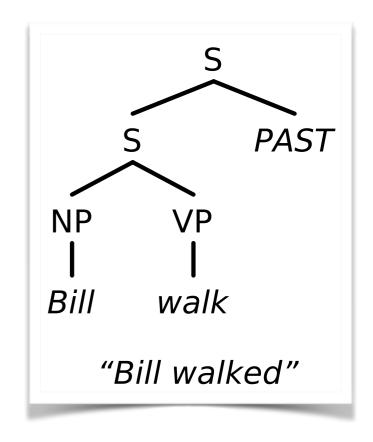
- Un $E = \emptyset$,
- < ⊆ E×E is an asymmetric relation (temporal precedence)
- $e_u \in E$ is the utterance event
- V is an interpretation function like in standard FOL
- Overlapping events: e e' iff neither e < e' nor e' < e

Tense in Semantic Construction

We can represent inflection as an abstract tense operator reflecting the temporal location of the reported event relative to the utterance event.

PAST
$$\mapsto \lambda P.\exists e [P(e) \land e < e_u] : \langle\langle e, t \rangle, t \rangle$$

PRES
$$\mapsto \lambda P. \exists e [P(e) \land e \cdot e_u] : \langle \langle e, t \rangle, t \rangle$$



Tense in Semantic Construction

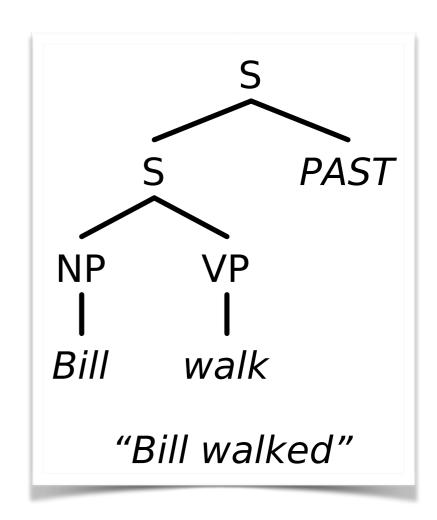
Standard function application results in integration of temporal information and binding of the event variable (i.e., replacing E-CLOS):

- walk $\mapsto \lambda x \lambda e [walk(e, x)]$
- Bill walk $\mapsto \lambda x \lambda e$ [walk(e, x)](b') $\Rightarrow^{\beta} \lambda e$ [walk(e, b')]
- Bill walk PAST

 → λΕ ∃e [E(e) ∧ e < e_u](λe' [walk(e', b)])

 ⇒β ∃e [λe' [walk(e', b)](e) ∧ e < e_u]

 ⇒β ∃e [walk(e, b) ∧ e < e_u]



Advantages of Davidsonian events

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Verbal arguments; a related problem?

- (1) John **broke** the window with a rock.
- (2) A rock broke the window.
- (3) The window broke.

And we're back to the same entailment issue:

 $\exists e(break_3(e, j, w, r)) \nvDash \exists e(break_2(e, r, w)) \nvDash \exists e(break_1(e, w))$

Semantic/Thematic roles

```
agent patient instrument

(1) John broke the window with a rock

→ ∃e [break(e) ∧ agent(e, j) ∧ patient(e, w) ∧ instrument(e, r)]

(2) A rock broke the window.

→ ∃e [break(e) ∧ patient(e, w) ∧ instrument(e, r)]

(3) The window broke.

→ ∃e [break(e) ∧ patient(e, w)]
```

In standard FOL: Thematic roles are implicitly represented by the canonical order of the arguments

In Davidsonian event semantics: Thematic roles are two-place relations between the event denoted by the verb, and an argument role filler.

Interpretation of events with thematic roles

```
John kisses Mary \rightarrow \exists e \text{ (kiss(e)} \land agent(e, j')} \land patient(e,m'))
 \llbracket \exists e \text{ (kiss(e)} \land agent(e, j')} \land patient(e, m')) \rrbracket^{M,g} = 1
iff there is an s \in E such that [kiss(e)]^{M,g[e/s]} = 1 and [agent(e, j')]^{M,g[e/s]} = 1
       and [patient(e,m')]^{M,g[e/s]} = 1
iff there is an s \in E such that s \in V_M(kiss) and \langle s, V_M(j') \rangle \in V_M(agent)
       and \langle s, V_M(m') \rangle \in V_M(patient)
                                                                                     VM kiss
                                                          agent
```

Thematic roles & verbal differences/similarities

Different verbs allow different thematic role configurations

- (1) a. John broke the window with a rock ———— agent, patient, instrument
 - b. John **smiled** at Mary ——— agent, recipient
- (2) a. The window **broke** allows inanimate subject
 - b. *The bread **cut** does not allow inanimate subject

Thematic roles capture equivalences and entailment relations between different predicates

- (3) a. Mary gave Peter the book $\forall e[give(e) \leftrightarrow receive(e)] \models (3a) \leftrightarrow (3b)$
 - b. Peter **received** the book from Mary

Determining the role inventory

Fillmore (1968): "thematic roles form a small, closed, and universally applicable inventory conceptual argument types."

A typical role inventory might consist of the roles:

 Agent, Patient, Theme, Recipient, Instrument, Source, Goal, Beneficiary, Experiencer.

But... there are some difficult cases:

- (1) Lufthansa is replacing its 737s with Airbus 320
- (2) John sold the car to Bill for 3,000€
- (3) Bill bought the car from John for 3,000€

Semantic corpora with thematic roles

- PropBank: includes a separate role inventory for every lemma
- FrameNet: "Frame-based" role inventories

Frames are structured schemata representing complex prototypical situations, events, and actions

- (1) [Agent Lufthansa] is replacing Frame: REPLACING [Old its 737s] [New With Airbus A320s]
- (2) [Agent Lufthansa] is substituting Frame: REPLACING [New Airbus A320s] [Old for its 737s]

Semantic corpora with thematic roles (cont.)

PropBank (Palmer et al. 2005): Annotation of Penn TreeBank with predicate-argument structure.

- (1) [Arg0 Lufthansa] is replacing [Arg1 its 737s] [Arg2 with Airbus A320s]
- (2) [Arg0 Lufthansa] is substituting [Arg1 Airbus A320s] [Arg2 for its 737s]

FrameNet (Baker et al. 1998): A database of frames and a lexicon with frame information

- (3) [Agent Lufthansa] is replacing_{Frame: REPLACING} [Old its 737s] [New with Airbus A320s]
- (4) [Agent Lufthansa] is substituting Frame: REPLACING [New Airbus A320s] [Old for its 737s]

Pred	replace
Arg0	Lufthansa
Argl	its737s
Arg2	AirbusA320s

Pred	substitute
Arg0	Lufthansa
Argl	AirbusA320s
Arg0 Arg1 Arg2	its737s

Frame	REPLACING
Agent Old	Lufthansa
Old	its737s
New	AirbusA320s
1	

Advantages of Davidsonian events

- ☑ Intuitive representation and semantic construction for adjuncts
- Uniform treatment of verb complements
- Uniform treatment of adjuncts and post-nominal modifiers
- Plausible treatment of tense information
- Compatible with analysis of semantic roles

... but how does it combine with other semantic constructs?

A problem with events and quantification

```
John kissed Mary
\rightarrow \lambda P.P(j') \left[ \lambda P.P(m')(\lambda y \lambda x \lambda e \left[ kiss(e) \wedge agent(e,x) \wedge patient(e,y) \right] \right]
\Rightarrow^{\beta} \lambda e \text{ [kiss(e)} \wedge \text{agent(e,i')} \wedge \text{patient(e,m')]}
\RightarrowE-CLOS \existse [kiss(e) \land agent(e,j') \land patient(e,m')]
John kissed every girl
\rightarrow \lambda P.P(j') \left[ \lambda P. \forall x(girl'(x) \rightarrow P(x))(\lambda y \lambda x \lambda e \left[ kiss(e) \land agent(e,x) \land patient(e,y) \right] \right]
\Rightarrow^{\beta} \lambda e \left[ \forall x (girl'(x) \rightarrow kiss(e) \land agent(e,j') \land patient(e,x) \right]
\RightarrowE-CLOS \existse [\forall x(girl'(x) \rightarrow kiss(e) \land agent(e,j') \land patient(e,x)]
```

Two solutions to the event quantification problem

Solution I

Interpret sentences as generalized quantifiers over events: $\langle\langle e,t \rangle,t \rangle$ instead of $\langle e,t \rangle$ (E-CLOS part of lexical semantics) (Champollion, 2010; 2015)

kiss
$$\mapsto \lambda F_{(v,t)}$$
. $\exists e \text{ (kiss(e) } \wedge F(e)) :: \langle \langle v,t \rangle, t \rangle \approx \{ F \mid F \cap KISS \neq \emptyset \}$

Solution II

Solution II

Introduce separate types for regular NPs and quantified NPs, and restrict existential closure to regular NPs (Winter & Zwarts, 2011; de Groote & Winter, 2014)

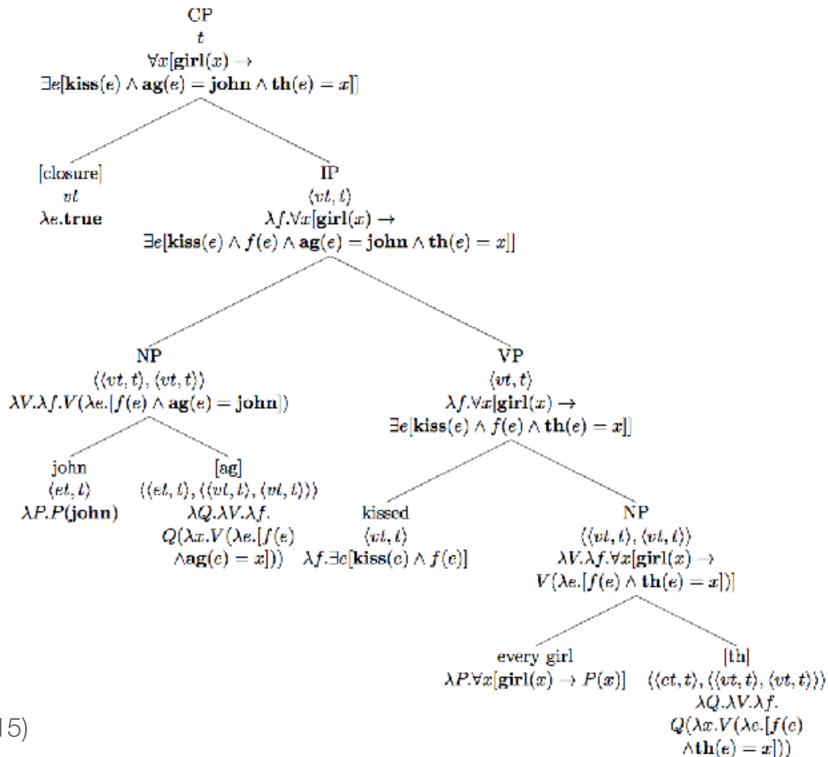
```
john \mapsto j :: e

every girl \mapsto \lambda P\lambda Q. \forall x (girl(x) \rightarrow Q(x)) :: \langle \langle e, t \rangle, \langle \langle e, t \rangle, t \rangle \rangle

kiss \mapsto \lambda x\lambda y\lambda e. kiss(e, x, y) :: \langle e, \langle e, \langle v, t \rangle \rangle \rangle

e-clos \mapsto \lambda P. \exists e(P(e)) :: \langle \langle v, t \rangle, t \rangle separate type for events!
```

Solution I: Sentences as GQs over events



Solution II: Type-restriction for existential closure

$$\frac{\vdash \text{EVERY} : N \to (NP \to S) \to S \qquad \vdash \text{GIRL} : N}{\vdash \text{EVERY GIRL} : (NP \to S) \to S} \tag{1}$$

$$\frac{\vdash \text{KISSED} : NP \to NP \to V \qquad x : NP \vdash x : NP}{x : NP \vdash \text{KISSED} \ x : NP \vdash \text{KISSED} \ x : NP \vdash \text{KISSED} \ x \text{ JOHN} : V} \tag{2}}{x : NP \vdash \text{E-CLOS} : V \to S \qquad x : NP \vdash \text{KISSED} \ x \text{ JOHN} : V}} \tag{2}$$

$$\frac{\vdash \text{E-CLOS} : V \to S \qquad x : NP \vdash \text{KISSED} \ x \text{ JOHN} : V}{x : NP \vdash \text{E-CLOS} (\text{KISSED} \ x \text{ JOHN}) : S} \tag{3}}{\vdash \lambda x \cdot \text{E-CLOS} (\text{KISSED} \ x \text{ JOHN}) : NP \to S}} \tag{3}$$

$$\frac{\vdash \text{EVERY GIRL} : (NP \to S) \to S \qquad \vdash \lambda x \cdot \text{E-CLOS} (\text{KISSED} \ x \text{ JOHN}) : NP \to S}{\vdash \text{EVERY GIRL} : (\lambda x \cdot \text{E-CLOS} (\text{KISSED} \ x \text{ JOHN})) : S}}$$

Links

- Overview paper: Lasersohn (2012) Event-Based Semantics: https://semanticsarchive.net/Archive/jFhNWM2M/ eventbasedsemantics.pdf
- PropBank: http://propbank.github.io/
- FrameNet: https://framenet.icsi.berkeley.edu/fndrupal/