CptS 415 Big Data

Design of Relational Data Models

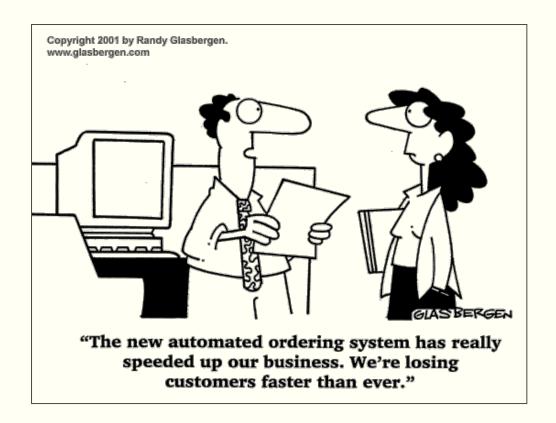
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Acknowledgements: Tinghui Wang



Design of relational data models

- Functional dependencies
- The normal forms 1NF, 2NF, 3NF, BCNF



A Big Data Fallacy

- Data model design in the era of Big Data is less important?
 - New high-volume data streams
 - Specialized hardware/software
 - Storage issues coped by hardware appliance
- Fact
 - Most data is physically located in DBMS and new special-purpose appliance
 - Data loads, extract, transform preprocessing operations continue as is
 - Database design for quality assurance

Big data a necessity at Largest Scale

A certain kind of developer at a certain kind of company

Most development still RDBMS

MySQL, Oracle, Mongo, Cassandra, some memcache, Some Hadoop...

Relational Databases Design

- Relational Database design:
 - The grouping of attributes to form "good" relation schemas
- We have assumed schema R is given
 - Universal Relation: R could have been a single relation containing all attributes that are of interest
 - Normalization breaks R into smaller relations
 - R could have been the results of some ad-hoc design of relations, which we then test/convert to normal form

The Need for Normalization

Example:

- Charges its clients by billing hours spent on each contract
- Hourly billing rate is dependent on employee's position

TABLE 5.1 A SAMPLE REPORT LAYOUT

PROJ. NUM.	PROJECT NAME	EMPLOYEE NUMBER	EMPLOYEE NAME	JOB CLASS.	CHG/HOUR	HOURS BILLED	TOTAL CHARGE
15	Evergreen	103 101 105 106 102	June E. Arbough John G. News Alice K. Johnson * William Smithfield David H. Senior	Elec. Engineer Database Designer Database Designer Programmer Systems Analyst	\$84.50 \$105.00 \$105.00 \$35.75 \$96.75	23.8 19.4 35.7 12.6 23.8	\$2,011.10 \$2,037.00 \$3,748.50 \$450.45 \$2,302.65
				Subtotal			\$10,549.70
18	Amber Wave	114 118 104 112	Annelise Jones James J. Frommer Anne K. Ramoras * Darlene M. Smithson	Applications Designer General Support Systems Analyst DSS Analyst	\$48.10 \$18.36 \$96.75 \$45.95	24.6 45.3 32.4 44.0	\$1,183.26 \$831.71 \$3,135.70 \$2,021.80
				Subtotal			\$7,171.47
22	Rolling Tide	105 104 113 111 106	Alice K. Johnson Anne K. Ramoras Delbert K. Joenbrood * Geoff B. Wabash William Smithfield	Database Designer Systems Analyst Applications Designer Clerical Support Programmer	\$105.00 \$96.75 \$48.10 \$26.87 \$35.75	64.7 48.4 23.6 22.0 12.8	\$6,793.50 \$4,682.70 \$1,135.16 \$591.14 \$457.60
				Subtotal			\$13,660.10
25	Starflight	107 115 101 114 108 118 112	Maria D. Alonzo Travis B. Bawangi John G. News * Annelise Jones Ralph B. Washington James J. Frommer Darlene M. Smithson	Programmer Systems Analyst Database Designer Applications Designer Systems Analyst General Support DSS Analyst	\$35.75 \$96.75 \$105.00 \$48.10 \$96.75 \$18.36 \$45.95	24.6 45.8 56.3 33.1 23.6 30.5 41.4	\$879.45 \$4,431.15 \$5,911.50 \$1,592.11 \$2,283.30 \$559.98 \$1,902.33
				Subtotal			\$17,559.82
				Total			\$48,941.09
* Indicates	oroject leader						

A Table in the Report Format

	PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
Þ	15	Evergreen	103	June E. Arbough	Elect. Engineer	\$84.50	23.6
			101	John G. News	Database Designer	\$105.00	1.9.4
			105	Alice K. Johnson *	Database Designer	\$105.00	35.
			106	William Smithfield	Programmer	\$35.75	12.0
			102	David H. Senior	Systems Analyst	\$96.75	23.
	18	Amber Wave	114	Annelise Jones	Applications Designer	\$48.10	24,0
			115	James J. Frommer	General Support	\$118,36	45.
			104	Anne K. Ramoras *	Systems Analyst	\$96.75	32.
			112	Darlene M. Smithson	DSS Analyst	\$45.95	44.
	22	Rolling Tide	105	Alice K. Johnson	Database Designer	\$105,00	64.
			104	Anne K. Ramoras	Systems Analyst	\$96.75	48.
			11:3:	Delbert K. Joenbrood *	Applications Designer	\$48:10	23.0
			1111	Geoff B. Wabash	Clerical Support	\$26.87	22.
			106	William Smithfield	Programmer	\$35.75	12,
	25	Starflight	107	Maria D. A 00 so	Programmer	\$35.75	24.0
			44:5	Travis B. Bawangi	Systems Analyst	\$96.75	45.
			104	John G. Nej vs *	Database Designer	\$105:00	56
			114	Annelise Jones	Applications Designer	\$48.10	30,
			1,08:	Ralph B. Washington	Systems Analyst	\$96.75	23,
			118	James: J. Frommer	General Support	\$118:36:	30.
			11:2	Darlene M. Smithson	DSS Analyst	\$45.95	431

The Need for Normalization

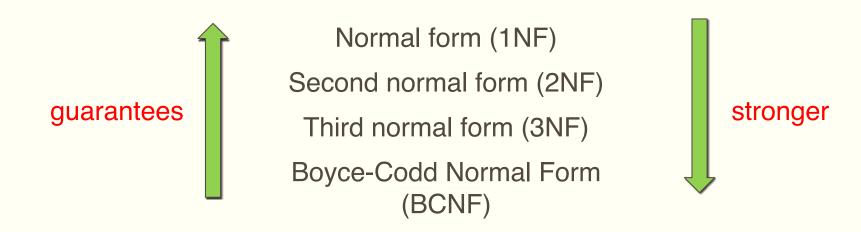
- Mixing attributes of multiple entities may cause problems
 - Information is stored redundantly wasting storage
 - Problems with update anomalies
 - Insertion anomalies
 - Deletion anomalies
 - Modification anomalies
- The report may yield different results depending on data anomaly
 - Primary keys
 - Redundancy
 - Possible data inconsistencies
 - E.g. JOB_CLASS: Elect Engineer, Elect. Eng., E. E., El. Eng.

Normalization

 Process of decomposing unsatisfactory "bad" relations by breaking up their attributes into smaller relations

Normal Form:

- Condition using keys and functional dependencies (FDs) of a relation to certify whether a relation schema is in a particular normal form
- 2NF, 3NF, BCNF based on keys and FDs of a relation schema
- 4NF based on keys, multi-valued dependencies



Functional Dependencies (FDs)

- FDs are used to specify formal measures of the "goodness" of relational designs
- lacksquare A set of attributes X functionally determines a set of attributes Y if the value of X determines a unique value Y
- $X \to Y$ holds if whenever two tuples have the same value for X, they must have the same value for Y on all relation instances

$$\forall t_1, t_2 \in r(R), \ t_1[X] = t_2[X] \implies t_1[Y] = t_2[Y]$$

- FDs are derived from the real-world constraints on the attributes
 - SSN → ENAME
 - PNUBER → {PNAME, PLOCATION}
 - {SSN, PNUMBER} → HOURS
 - What if LHS is a key?

Function Key (Cont'd)

$$A \rightarrow C$$

$$\bullet$$
 $A, B > \rightarrow C$

$$A \rightarrow D$$

$$C \rightarrow D$$

А	В	С	D
1	1	3	2
1	2	2	3
1	3	3	2

Full/Partial FD

- If removal of any attribute from X means the FD does not hold any more, it is a full dependency; otherwise it's a Partial dependency
- Transitive FD: $X \rightarrow Y$
 - If there is a set of attributes Z that are neither a primary or candidate key and both $X \to Z$ and $Z \to Y$ holds.

Inference Rules for FDs

- Given a set of FDs F, we can infer additional FDs that hold whenever the FDs in F holds.
- Armstrong's inference rules
 - $\begin{tabular}{ll} \blacksquare & {\sf Reflexive} \\ & {\sf if} \ Y \subseteq X, \ {\sf then} \ X \to Y \\ \end{tabular}$
 - Argumentation if $X \to Y$, then $X \cup Z \to Y \cup Z$ (or sometime written as $XZ \to YZ$)
 - Transitive if $X \to Y$ and $Y \to Z$, then $X \to Z$

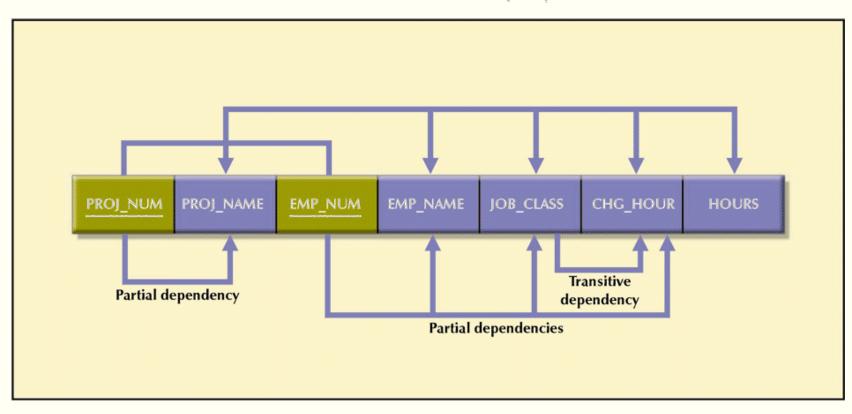
First Normal Form

- Tabular format in which
 - All key attributes are defined
 - There are no repeating groups in the table
 - All attributes are dependent on primary key
- Disallow composite attributes, multi-valued attributes and nested relations
- 1NF deals with the "shape" of the tables

1NF Normalization

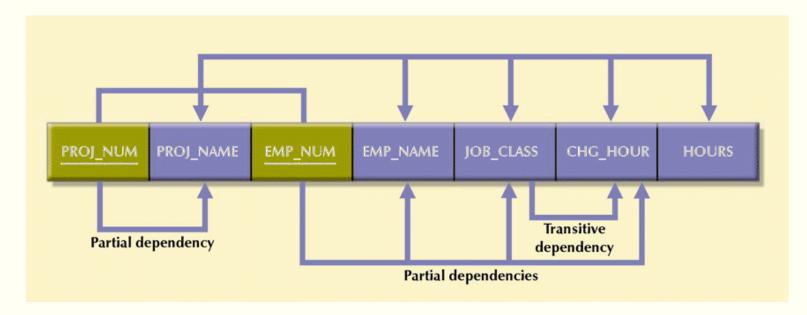
	PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
1	15	Evergreen	1.03	June E. Arbough	Elect. Engineer	\$84.50	23.
	15	Evergreen	101	Jahn G. News	Database Designer	\$105.00	19.
	15	Evergreen	105	Alice K. Johnson *	Database Designer	\$105.00	35
	15	Evergreen	106	William Smithfield	Programmer	\$35.75	12
	15	Evergreen	102	David H. Senior	Systems Analyst	\$96.75	23
	18	Amber Wave	114	Annelise Jones	Applications Designer	\$48.10	24
	18	Amber Wave	118	James J. Frommer	General Support	\$18.36	45
	18	Amber Wave	1.04	Anne: K. Ramoras *	Systems Analyst	\$96.75	32
	18	Amber Wave	112	Darlene M. Smithson	DSS Analyst	\$45.95	44
	22	Rolling Tide	1.05	Alice K. Johnson	Database Designer	\$105.00	64
	22	Rolling Tide	104	Anne: K. Ramoras	Systems Analyst	\$96.75	48
	22	Rolling Tide	113	Delbert K. Joenbrood *	Applications Designer	\$48.10	23
	22	Rolling Tide	111	Geoff B. Wabash	Clerical Support	\$26.87	22
	22	Rolling Tide	1:06	William Smithfield	Programmer	\$35.75	12
	25	Starflight	1:07	Maria D. Alonzo	Programmer	\$35.75	24
	25	Starflight	115	Travis B. Bawangi	Systems Analyst	\$96.75	45
	25	Starflight	1.01	John G. News *	Database Designer	\$105.00	56
	25	Starflight	114	Annelise Jones	Applications Designer	\$48.10	33
	25	Starflight	105	Ralph 5. Washington	Systems Analyst	\$96.75	23
	25	Starflight	118	James J. Frommer	General Support	\$18.36	30
	25	Starflight	112	Darlene M. Smithson	DSS Analyst	\$45,95	41

FIGURE 5.3 A DEPENDENCY DIAGRAM: FIRST NORMAL FORM (1NF)



Second Norm Form (2NF)

- Table is in 2NF if
 - It is in 1NF
 - It includes no partial dependencies
 - No attribute is dependent on only a portion of the primary key
 - Every attribute A not in PK is fully functionally dependent on PK
- 2NF deals with the relationship between non-key and key attributes



Conversion 1NF to 2NF

- Step 1: Write each key attribute on separate line and then write the original (composite) key on the last line; Each component will become the key in a new table.
- Step 2: Determine which attributes are dependent on which other attributes (remove anomalies)

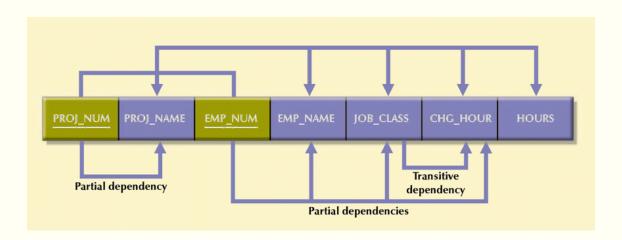
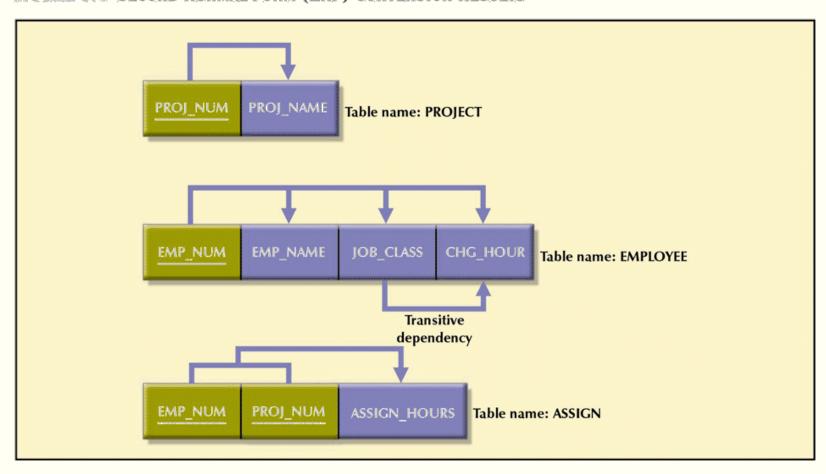


FIGURE 5.4 SECOND NORMAL FORM (2NF) CONVERSION RESULTS



Third Normal Form (3NF)

- A table is in third normal form (3NF) if
 - It is in 2NF
 - It contains no transitive dependencies
- 3NF removes transitive dependencies

Conversion to 3NF

- Step 1: Find new Fact
 - For every transitive dependency $X \to Y$, write fact Z as a PK for a new table where $X \to Z$ and $Z \to Y$
- Step 2: Identify the dependent attributes
 - lacktriangleright Identify the attributes dependent on each Z identified in previous step and find the dependency
 - Name the table to reflect its contents and function
- Step 3: Remove $X \to Y$ from original table

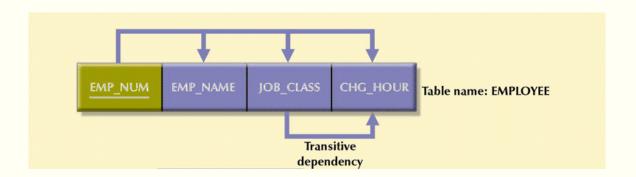
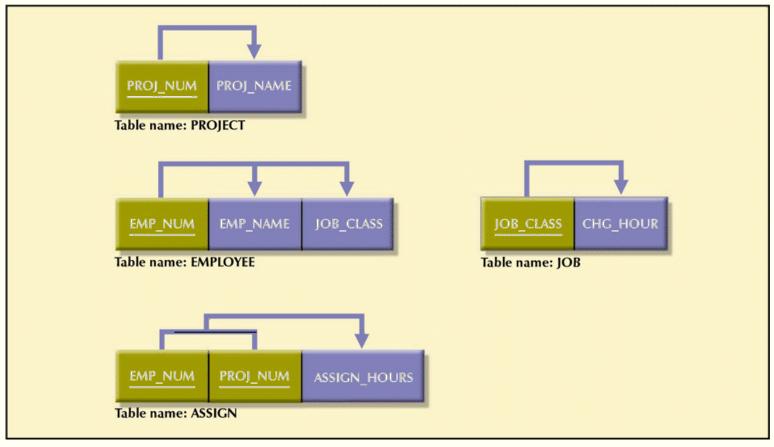


FIGURE 5.5 THIRD NORMAL FORM (3NF) CONVERSION RESULTS



Boyce-Codd Normal Form (BCNF)

■ A relation schema R is in BCNF, a.k.a. 3.5NF, if whenever an FD $X \to A$ holds in R, then X is a super-key of R

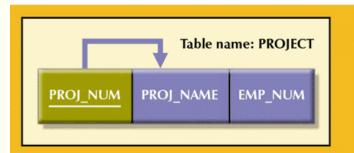
FIGURE 5.8 DECOMPOSITION TO BCNF 3NF, but not BCNF 1NF Partial dependency 3NF and BCNF 3NF and BCNF

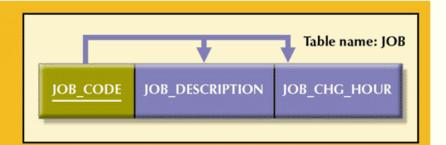
Normalization

- Each normal form is strictly stronger than the previous one
 - Every 2NF relation is in 1NF
 - Every 3NF relation is in 2NF
 - Every BCNF relation is in 3NF
- There exist relations that in 3NF but not in BCNF
- The goal is to have each relation in BCNF (or 3NF)

The Completed Database

FIGURE 5.6 THE COMPLETED DATABASE





Database name: Ch05_ConstructCo

Table name: PROJECT

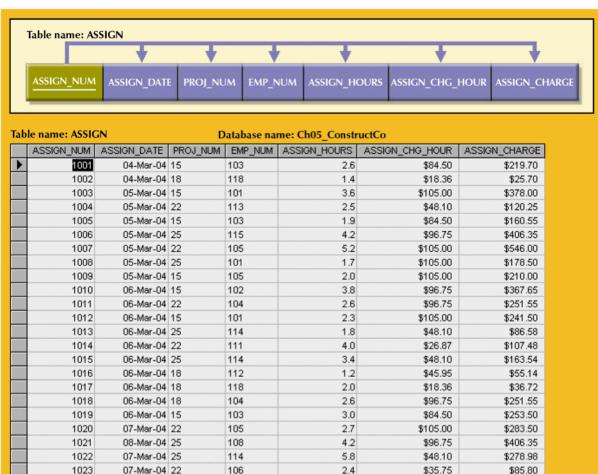
		PROJ_NUM	PROJ_NAME	EMP_NUM
•	+	15	Evergreen	105
	+	18	Amber Wave	104
	+	22	Rolling Tide	113
	+	25	Starflight	101

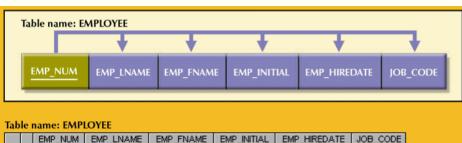
Table name: JOB

		JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
•	+	500	Programmer	\$35.75
	+	501	Systems Analyst	\$96.75
	+	502	Database Designer	\$105.00
	+	503	Electrical Engineer	\$84.50
	+	504	Mechanical Engineer	\$67.90
	+	505	Civil Engineer	\$55.78
	+	506	Clerical Support	\$26.87
	+	507	DSS Analyst	\$45.95
	+	508	Applications Designer	\$48.10
	+	509	Bio Technician	\$34.55
	+	510	General Support	\$18.36

The Completed Database

FIGURE 5.6 THE COMPLETED DATABASE (CONTINUED)





		EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE
M	+	101	News	John	G	08-Nov-98	502
	+	102	Senior	David	Н	12-Jul-87	501
	+	103	Arbough	June	E	01-Dec-94	503
	+	104	Ramoras	Anne	K	15-Nov-85	501
	+	105	Johnson	Alice	K	01-Feb-91	502
	+	106	Smithfield	∨Villiam		22-Jun-03	500
	+	107	Alonzo	Maria	D	10-Oct-91	500
	+	108	√Vashington	Ralph	В	22-Aug-89	501
	+	109	Smith	Larry	W	18-Jul-95	501
	+	110	Olenko	Gerald	A	11-Dec-93	505
	+	111	√Vabash	Geoff	В	04-Apr-89	506
	+	112	Smithson	Darlene	M	23-Oct-92	507
	+	113	Joenbrood	Delbert	K	15-Nov-94	508
	+	114	Jones	Annelise		20-Aug-91	508
	+	115	Bawangi	Travis	В	25-Jan-90	501
	+	116	Pratt	Gerald	L	05-Mar-95	510
	+	117	∨Villiamson	Angie	Н	19-Jun-94	509
	+	118	Frommer	James	J	04-Jan-04	510