# Washington State University School of Electrical Engineering and Computer Science CptS 451 – Introduction to Database Systems Online

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## Homework-2 ER to Relational Translation

Name:	Solution Key
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Student Number	

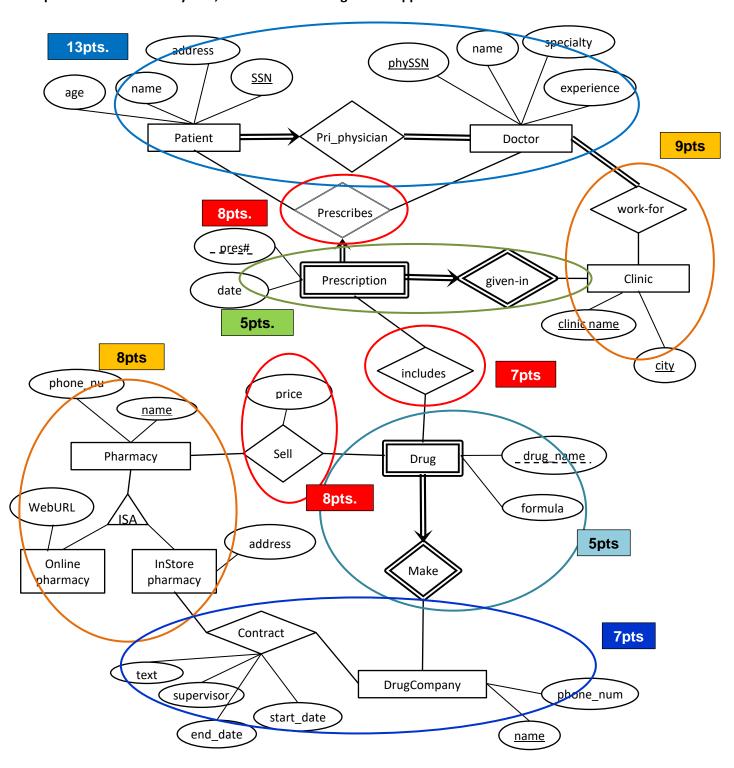
Question:	Max points:	Score:
1	70	
2	30	
Total	100	

### Question 1 (70 pts)

Consider the ER diagram illustrated in the following figure (this is similar to the ER diagram from HW1 - couple modifications have been made to make the diagram simpler.)

Convert the ER diagram to relations and write SQL DDL statements for creating the tables for those relations. Pick suitable data types for each attribute. For string attributes pick reasonable lengths. Include the appropriate constraints (domain, primary key, foreign key, UNIQUE, and NULL constraints) in your SQL DDL statements.

Note: For the one-to-many binary relations, combine the relation with the many side. For the superclass-subclass entity sets, translate them using the ER-approach.



```
CREATE TABLE Doctor (
     phySSN CHAR(9)
     name VARCHAR(30),
     speciality VARCHAR (30),
     experience integer
     PRIMARY KEY (physsn)
   );
  CREATE TABLE Patient (
     patientSSN CHAR(9),
     phySSN CHAR(9) NOT NULL, --enforces total participation
     name
               VARCHAR(30),
     age
                INTEGER,
     address VARCHAR(100),
     PRIMARY KEY (patientSSN)
     FOREIGN KEY (physsn) REFERENCES Doctor(physsn)
  );
     -- Note: Can't enforce total participation of Doctor in "Pri_physician" relationship.
  CREATE TABLE Clinic (
     clinic name VARCHAR(30),
                  VARCHAR(30),
     PRIMARY KEY(clinic name, city)
     );
  CREATE TABLE Work-for (
     clinic name VARCHAR(30),
            VARCHAR(30),
     city
              CHAR (9),
     phySSN
     PRIMARY KEY (phySSN, clinic name, city),
     FOREIGN KEY (physsn) REFERENCES Doctor (physsn),
     FOREIGN KEY(clinic name, city) REFERENCES Clinic(clinic name, city)
     -- Note: Can't enforce total participation of Work-for in "Clinic" relationship.
-- In the below solution combine Prescribes relation with Prescription
(many-one-one relationship.
-- I will also accept solutions which include separate relations for
Prescribes and Prescription.
  CREATE TABLE Prescription (
     pres num INTEGER,
     clinic name VARCHAR(30),
                  VARCHAR (30),
     city
     patientSSN CHAR(9),
phySSN CHAR(9),
     date
                  DATE,
     PRIMARY KEY (pres num, clinic name, city),
     FOREIGN KEY (clinic name, city) REFERENCES Clinic (clinic name, city)
     FOREIGN KEY (patientSSN) REFERENCES Patient(patientSSN),
     FOREIGN KEY (phySSN) REFERENCES Doctor(phySSN)
   );
```

```
CREATE TABLE DrugCompany (
   company name VARCHAR(20),
                 CHAR (10),
   phone num
   PRIMARY KEY (company name)
);
CREATE TABLE Drug (
   drug name VARCHAR(20),
   company name VARCHAR(20),
   formula VARCHAR (255),
   PRIMARY KEY(drug name, company name),
   FOREIGN KEY (company name) REFERENCES DrugCompany (company name)
);
CREATE TABLE Pharmacy (
   phar name VARCHAR(20),
   phone num CHAR(10),
   PRIMARY KEY (phar name)
);
CREATE TABLE OnlinePharmacy (
   phar name VARCHAR(20),
   Weburl Varchar (255),
   PRIMARY KEY (phar name)
   FOREIGN KEY (phar name) REFERENCES Pharmacy (phar name)
);
CREATE TABLE InStorePharmacy (
   phar name VARCHAR(20),
   address VARCHAR (100),
   PRIMARY KEY (phar name)
   FOREIGN KEY (phar name) REFERENCES Pharmacy(phar name)
 );
CREATE TABLE Sell (
   phar_name VARCHAR (20),
drug name VARCHAR (20),
   company_name VARCHAR (20), price FLOAT,
   PRIMARY KEY (phar name, drug name, company name),
   FOREIGN KEY (phar name) REFERENCES Pharmacy (phar name),
   FOREIGN KEY (drug name, company name) REFERENCES
               Drug (drug name, company name)
);
CREATE TABLE Contract (
   phar_name VARCHAR (20),
  company name VARCHAR (20),
   text
                  VARCHAR (255),
   supervisor
                VARCHAR (30),
DATE,
   start date
   end_date DATE,
```

```
PRIMARY KEY(phar_name,company_name),
FOREIGN KEY(phar_name) REFERENCES InStorePharmacy(phar_name),
FOREIGN KEY(company_name) REFERENCES DrugCompany(company_name));
```

### Question 2.

Consider the following relations:

### Relation R1

	<u>A</u>	<u>B</u>	С
1	a1	b10	s100
2	a2	b10	s300
3	a2	b20	s200
4	a3	b10	s500
5	a4	b20	s100

### Relation R2

	D	Е	F
1	d10	50	100
2	d20	125	200
3	d30	150	300
4	d40	75	400
5	d50	100	200

### Relation R3

	<u>M</u>	<u>N</u>	0	Р
1	a1	b10	d10	25
2	a1	b10	d20	5
3	a2	b10	d20	20
4	a2	b20	d20	15
5	a3	b10	d40	15
6	a4	b20	d40	5
7	a4	b20	d50	10

### Relation R4

	<u>J</u>	<u>K</u>	L
1	s200	d20	22
2	s500	d50	55

### Relation R5

	<u>S</u>	T	U
1	s100	20	555
2	s200	20	333
3	s300	30	111
4	s400	30	555
5	s500	40	444

### Primary Keys:

Relation R1: A,B
 Relation R2: D
 Relation R3: M,N,O
 Relation R4: J,K
 Relation R5: S

The following foreign key constraints are given for relations R1, R2, R3, R4 and R5:

- 1. R3(MN) references R1(AB)
- 2. R3(O) references R2(D)
- 3. R1(C) references R5(S)
- 4. R4(J) references R5(S)
- 5. R4(K) references R2(D)

### Assume

- "CASCADE" policy for delete operations, and
- "SET NULL" policy for update operations.

- a) (10pts) For the operations given below, indicate whether execution of the operation would violate some "primary key" or "integrity constraints". If your answer is yes, specify the constraints (from the above list) that would be violated.
  - i) Insert tuple (a1,b10, d20, 35) into R3. Violates the PK constraint for R3. (M,N,O) is the primary key for R3 and (M,N,O) value ('a1','b10','d20') already exists in R3. Since duplicate primary key values are not allowed, the insertion will be rejected.
  - ii) Insert tuple (s500, d20, 75) into R4. No violation.

R4 has 2 foreign keys (FKs): R4(J) references R5(S); R4(K) references R2(D) The first FK is not violated, since 's500' already appears in R5[S]. The second FK is also not violated since 'd20' already appears in R2[D]

- b) (15pts) For the operations given below, indicate whether execution of the operation would violate any "foreign key constraints". If your answer is yes, specify the constraints (from the above list) that would be violated. Apply "CASCADE" policy for delete operations, and apply "SET NULL" policy for update operations. Update the tables after applying those policies. (You may either redraw the tables or mention which tuples/attributes are deleted/updated at each table. For updates rewrite the updated tuples.) Make the changes on the original tables for each operation below.
  - i) Delete tuple (d30, 150, 300) from R2. R2 has 2 references: R3(O) references R2(D); and R4(K) references R2(D) Since 'd30' doesn't appear in R3[O] or R4[K], there won't be any violation in FKs and no tuples will be deleted in R3 and R4.
  - **ii)** Update tuple (s400,30,555) in R5 with values (6000,60,66) R5 has 2 references: R1(C) references R5(S); and R4(J) references R5(S) Since 's400' doesn't appear R1[C] and R4[J], there won't be any FK violation. No tuples in R1 and R4 will be updated.
  - iii) Update tuple (s100, 20, 555) in R5 with values (6000,60,666) R5 has 2 references: R1(C) references R5(S); and R4(J) references R5(S). Since 's100' doesn't appear R4[J], no tuples in R4 will be updated. However, 's100' appears in R1[C] on lines #1 and #5. R1 will be updated as follows:

	<u>A</u>	<u>B</u>	С
1	a1	b10	NULL
2	a2	b10	s300
3	a2	b20	s200
4	a3	b10	s500
5	a4	b20	NULL

## c) (5pts) If all tuples in R5 are deleted, what tuples will R2 and R3 contain?

If all tuples in R5 are deleted, all tuples in R1 and R4 will be deleted. Since all tuples in R1 are deleted, all tuples in R3 will be deleted as well.

R2 doesn't have any foreign keys. So no tuples will be deleted from R2.

So R3 will be empty, R2 will stay as it is.