```
title: "coursework2_nkatz01_dsa"
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date: "12 December 2018"

1. Fast computation

I've found this solution here https://stackoverflow.com/questions/10195252/calculate-x-y-in-olog-n but after thoroughly studying it, I think I now understand how it works and could not think of a better solution myself.

```
public static int computeAtoPowN(int a, int n) {
  if (n == 0) {
    return 1;
}
  int results = computeAtoPowN(a, n / 2);
  if (n % 2 != 0) {
    return a * results * results;
} else {
    return results * results;
}
```

I believe the reason why this solution works is due to the realisation that a exponentiation operation can be broken down in the following. For an odd exponent n:

```
a^n=a(a^2)^(n-1)/2
```

(Let the $a(a^2)$ part of the equation, be called operation A.)

```
and for an even exponent n: a^n=(a^2)^n/2 (Let the (a^2) part of the equation be called operation B.)
```

This is what's happening in this algorithm.

- The base case is n==0, returning 1 to the last caller, always resulting in producing a value == to the base, for the second return.
- if/when the caller's n is odd, operation A is carried out (the ${\tt IF}$).
- if/when the caller's n is even, operation B is carried out (the second condition Else).
- and it happens n/2 [for even numbers] because, n is fed to itself n/2 every time
- it happens (n-1)/2 [for odd numbers] because n is an int and so the fraction is discarded everytime the 'division operations' is being performed when the function calls itself.
- yet, whilst in the case where if one before the last call to itself, n was even, coming back from the base case, the sequence of operations will be odd, even (first condition [if], second condition [else] eg 1,2 where n begins with 2); for a case where one before the last call to itself, n was odd, coming back from the base case, the sequence of operations will always be odd, odd one after the other (first condition [if] twice, eg 1,3 where n begins with 3), making up that additional multiplication operation (by the base non-square)

it needs due to its odd, odd integer.

2. Missing Number problem with a Sorted array public static int missingNumberlogn() { int $A[] = \{1, 2, 3, 4, 5, 6, 8, 9, 10\}; //eq$ int start = 0;int end = A.length - 1;int middle; while (end - start > 1) { middle = (start + end) / 2;if (A[middle] == (middle + 1))start = middle + 1; else end = middle - 1; if (start != end && A[end] - A[start] > 1) $\{// \text{ array is size } 2$ and difference >1 return (A[start] + 1); else if (start != 0) {//array is of size 1 but not the first element if (A[end] - A[end - 1] > 1)// no worries of subtracting previous element return A[end] - 1; else return A[end] + 1;//number missing is at the end boundary of sequence. program will reach this if array starts at 1 else if (A[start + 1] - A[start] > 1)//array is of size 1 and is the first element of array and difference to second element >1 return A[start] + 1; else return A[start] - 1; //number missing is at the beginning boundary of sequence. program will reach this if array starts at any

3. Subarray problem

3.1 Iterative algorithm

number after but not including 1

```
int[] A = \{1,2,3,4\}; //eg
    int[] B = \{1, 2, 5\}; //eq
    int count=0;
    for (int i=0; i<B.length; i++) {
      for (int j=0; j<A.length; j++)</pre>
        if (B[i]==A[i])
       {count++;
       break; }
    if (count==B.length)
    System.out.println("B is sub array of A");
   else
        System.out.println("B is not a sub array of A");
  }
  I believe this solution works because if not all values are found
also in A, then the variable count will not be the same at B.length.
  I believe this is a O(n^2) runtime because it's a loop within a
loop.
    3.2 Recursive algorithm
    To query of B is subarray of A, call the following method with
method call recursiveSol(0, A.length - 1, A, B); (The first param
refers to first element in B, the second, to the last element in A.)
   public static int recursiveSol(int j, int i, int[] A, int[] B) {
    int found = 0;
    if (B.length > A.length) {
      System.out.println("Array B is not sub of Array A");
      return B.length - 1;
    }
    if (i == -1) \{// base case - Array A's finished traversing back to
front
     return 0;
    }
    found += recursiveSol(j, i - 1, A, B); // increase indexes in A -
equivalent of inner loop - pass 1 (if matching
                                            // value found in A) or 0
(if matching value not found in A) through the
```

```
// recursive calls back to
the first caller.
    if (B[j] != A[i])
      found = found + 0;
   else
      found = found + 1;
   j++;
    if (i == A.length - 1 && j < B.length) // increase index in B -
equivalent of outer loop
      found += recursiveSol(j, i, A, B); // sum up the 1s gained (or
not) for each element in B - from each "whole loop
                                          // cycle" through A.
    if (j == 1 \&\& i == A.length - 1) {// it has finished reversing}
from all the recursion calls
      if (found == B.length)
        System.out.println("B is a subarray of A");
        System.out.println("B is not subarray of A");
    }
    return found; // added advantage: if one wants evaluation above to
be done by outside caller -
                 // that's possible.
  }
    3.3 Array Sorted (SortedAndOofN() uses helper method
compareArrays( int[] A, int j, int[] B, int i))
        public static void SortedAndOofN() {
    int[] A = \{1, 2, 3, 4\}; //eg used
    int[] B = {5,1,2,3};//eq used
    if (B.length>A.length)//we can do this as we know values aren't
repeated
   { System.out.println("Array B not subof array A");
   return; }
    int i=0;//index counter for B
    int j=0; //index counter for A
    int flag=0;
   while (i<B.length)
        if (flag==1 || j==A.length) //if value of corresponding index
in A is more than value of corresponding index in B, or if all indexes
in A have been visited - terminate the loop
        break; }
```

```
flag=compareArrays(A, j, B, i); //send array ref and index no
to compareArrays func
        System.out.println(flag);
        if (flag==-1)//if value of corresponding index in A is less
than value of corresponding index in B,
         i--;//pass again the current index no of Array B.
        j++;//whilst increasing the index of A, regardless.
      }
    if (flag==0)
    System.out.println("Array B is subbary of A");
   System.out.println("Array B is not subbary of A");
  public static int compareArrays( int[] A, int j, int[] B, int i){
    if (A[j]>B[i])
    return 1;//because Arrays are sorted, we know the chance for the
value to be found in A is gone.
    if (A[j] == B[i])
   return 0;
   else return -1; //we may still find the value further in A.
  }
```