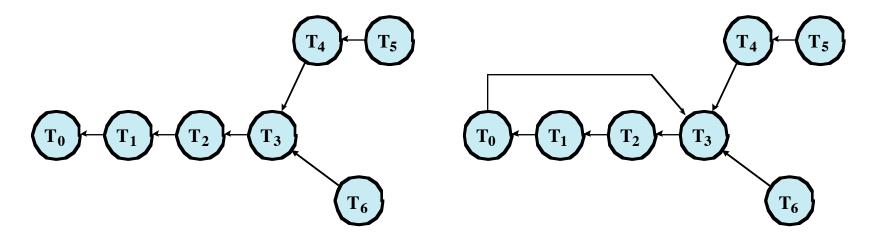
```
/* Data object D; receiving a lock_request(T;) */
                                                                 /* Transaction T<sub>i</sub> receiving an update message */
   if (Locked_by(Dj) == null)
                                                                    if (Wait_for(Ti) != Wait_for(Ti))
       send(granted);
   else
                                                                      Wait_for(Ti) = Wait_for(Ti);
                                                                    if (intersect(Wait_for(Ti)), Request_Q(Ti)) = null)
       send not granted to Ti;
                                                                      update(Wait for(Ti), Request Q(Tj);
       send Locked_by(Dj) to Ti
   }
                                                                        DECLARE DEADLOCK;
                                                                        /* initiate deadlock resolution as follows */
/* Transaction T_i makes a lock request for data object D_i
                                                                        /* T<sub>i</sub> is chosen as the transaction to be aborted */
                                                                        /* Ti releases all the data objects it holds */
   send lock_request(Ti) to Di;
                                                                        send_clear(Ti, Held_by(Ti));
   wait for granted/not granted;
                                                                        allocate each data object D_i held by T_j to the
   if (granted)
                                                                 first
                                                                               requester T_k in Request_Q(T_1);
      Locked_by(D_i) = T_i;
                                                                        for (every transaction T_n in Request_Q(Tj)
       Held by(T_i) = f;
                                                                 requesting
                                                                               data object Di held by Ti)
   else /* suppose D_{\dot{1}} is being used by transaction T_{\dot{1}} */
                                                                               Enqueue(T_n, Request_Q(T_k));
       Held_by(T_i) = T_i;
       Enqueue(Ti, Request_Q(Ti));
                                                                    }
       if (Wait for(Tj) == null)
          Wait_for(T_i) = T_i;
      else
                                                                 /* Transaction T_k receiving a clear(T_1, T_k) message */
          Wait for(T_i) = Wait for(T_j);
                                                                    purge the tuple having Tj as the requesting transaction
       update(Wait_for(Ti), Request_Q(Ti));
                                                                    from Request_Q(Tk);
```

Figure 15.14 A Distributed Deadlock Detection Algorithm



Transaction	Wait_for	Held_by	Request_Q
T_0	nil	nil	T_1
T_1	T_0	T_0	T_2
T_2	T_0	T_1	T_3
T_3	T_0	T_2	T_4, T_6
T_4	T_0	T_3	T_5
T_5	T_0	T_4	nil
T_6	T_0	T_3	nil

Transaction	Wait_for	Held_by	Request_Q
T_0	T_0	T_3	T_1
T_1	T_0	T_0	T_2
T_2	T_0	T_1	T_3
T_3	T_0	T_2	T_4, T_6, T_0
T_4	T_0	T_3	T_5
T_5	T_0	T_4	nil
T_6	T_0	T_3	nil

(a) State of system before request

(b) State of system after T_0 makes a request to T_3

Figure 15.15 Example of Distributed Deadlock Detection Algorithm of Figure 15.14