

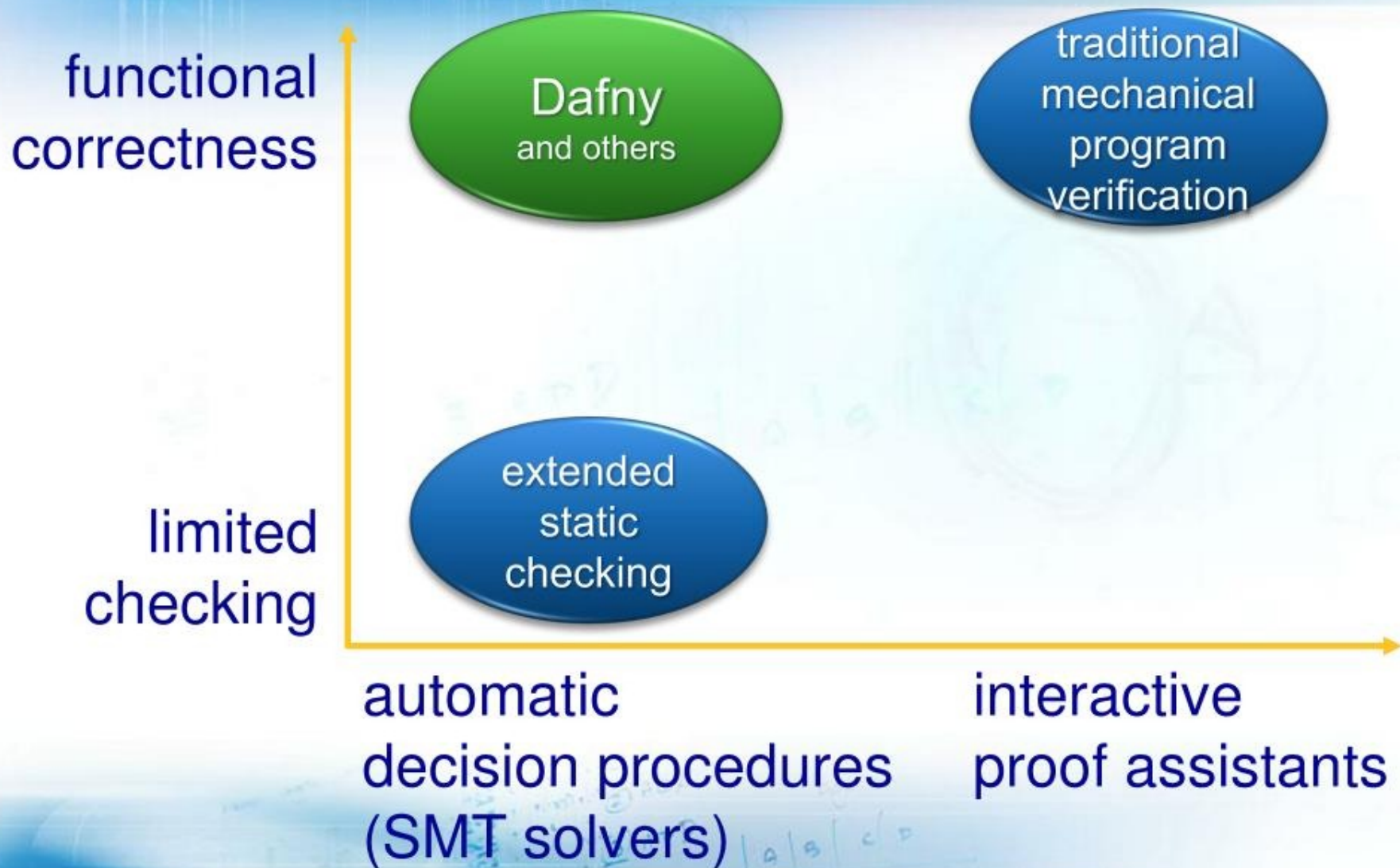
Dafny

An automatic program verifier
for functional correctness

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Program verification



User interaction

Program

Specification

Program oriented:
invariants,
assertions, ...

Formula

Formula oriented:
theorem-prover
commands, tactics

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Binary search

Dafny

- Object-based language
 - generic classes, no subclassing
 - object references, dynamic allocation
 - sequential control
- Built-in specifications
 - pre- and postconditions
 - framing
 - loop invariants, inline assertions
 - termination
- Specification support
 - Sets, sequences, algebraic datatypes
 - User-defined functions
 - Ghost variables

Top-level grammar

- Program ::= Type*
- Type ::= Class | Datatype
- Class ::= **class** Name { Member* }
- Member ::= Field | Method | Function
- Datatype ::= **datatype** Name { Constructor* }
- Generic (that is, accepts type parameters)

Types

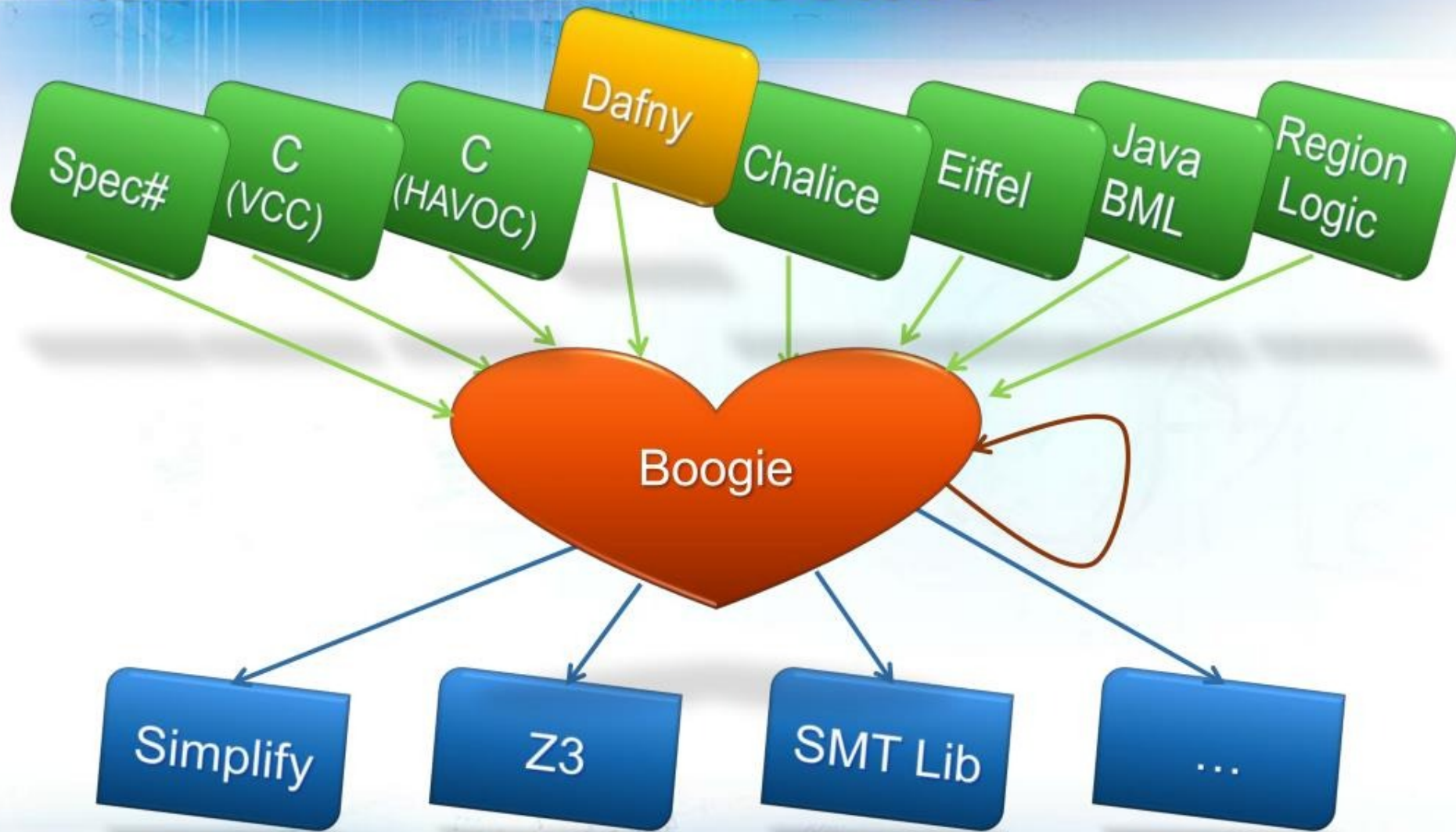
- Booleans
- Mathematical integers
- Finite sets
- Sequences
- Class types
- Algebraic datatypes

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Calculator

Verification architecture



Dafny, Boogie, VC

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From Dafny to verification-condition formulas

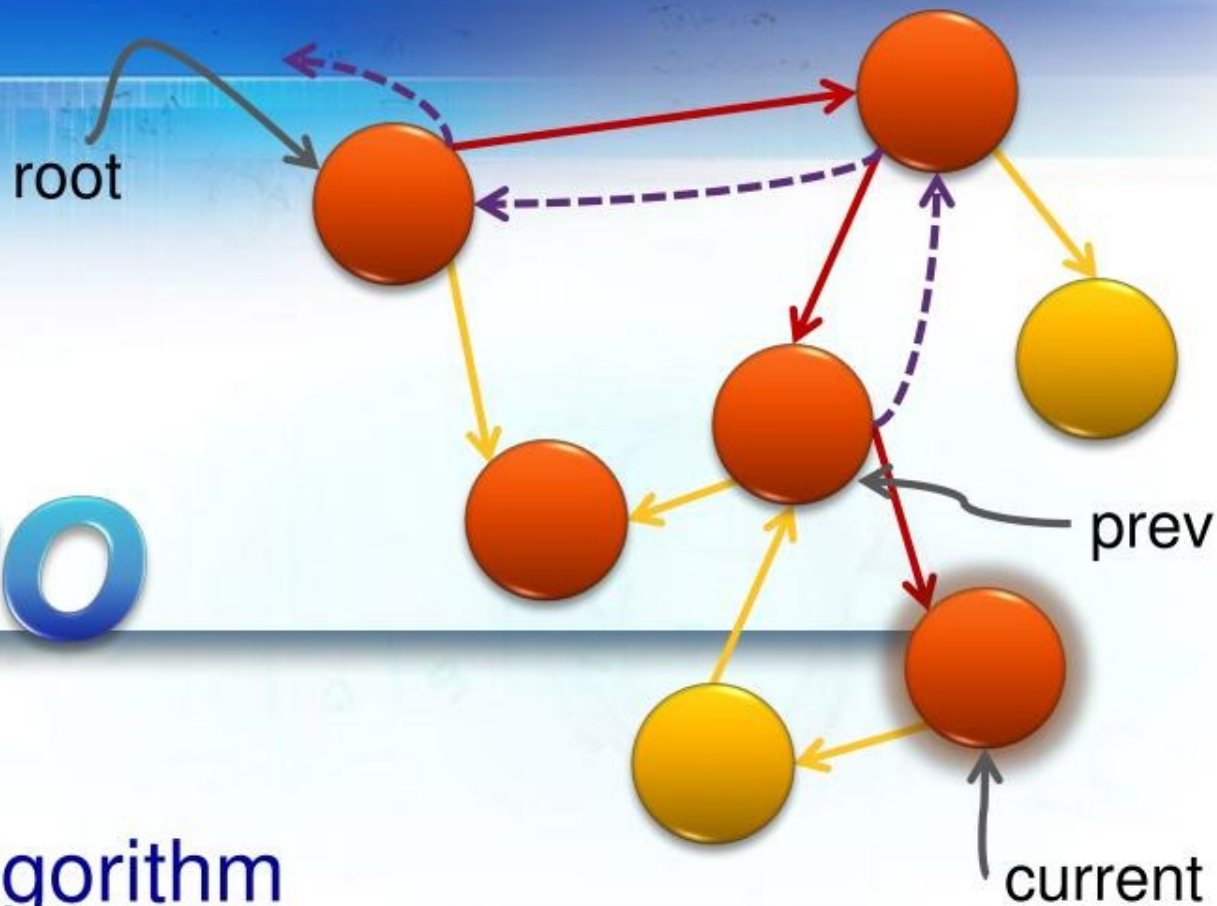
Axiomatizing functions

- function $F(x: T): U \dots \{ \text{Body} \}$
 - $(\forall x \bullet F(x) = \text{Body})$
- datatype $\text{Tree} \{ \text{Leaf}(\text{int}); \text{Split}(\text{Tree}, \text{Tree}); \}$
function $G(x: \text{Tree}): U \dots$
 $\{ \text{match } x$
 $\text{case Leaf}(n) \Rightarrow n$
 $\text{case Split}(a, b) \Rightarrow G(a) + G(b) \}$
 - $(\forall t \bullet G(t) = \text{if } \dots \text{else } G(\text{left}(t)) + G(\text{right}(t)))$
 - $(\forall n \bullet G(\text{Leaf}(n)) = n)$
 - $(\forall a, b \bullet G(\text{Split}(a, b)) = G(a) + G(b))$

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Schorr-Waite algorithm



Verifying termination

- Functions
- Loops
- Methods
- **decreases** clause
 - lexicographic tuple
 - components of tuple can be of any types
 - to compare, consider longest commonly typed prefix of the lexicographic tuple

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Using a program to prover a theorem

Conclusions

- Full functional-correctness verification is becoming more automatic
- Interaction is moving closer to the problem domain
- A well-designed language and verifier, plus a great SMT solver, go a long way

Dafny (and Boogie) open source: boogie.codeplex.com