$$\begin{split} \overline{\mathbf{Xi}: \{^m_i\}} & \to \mathbf{E1} & \overline{\mathbf{e}: \{^w_i | \, \mathbf{Xi} \in \mathrm{var}(\mathbf{e})\}} & \to \mathbf{E2} & \star \in \{+, -\} & \frac{\mathbf{Xi}: V_1 \quad \mathbf{Xj}: V_2}{\mathbf{Xi} \star \mathbf{Xj}: pV_1 \oplus V_2} & \to \mathbf{E3} \\ \\ \star \in \{+, -\} & \frac{\mathbf{Xi}: V_1 \quad \mathbf{Xj}: V_2}{\mathbf{Xi} \star \mathbf{Xj}: V_1 \oplus pV_2} & \to \mathbf{E4} & \frac{\mathbf{e}: V}{\mathbf{Xj} = \mathbf{e}: \mathbf{1} \overset{\mathbf{j}}{\leftarrow} V} & \mathbf{A} & \frac{\mathbf{C1}: M_1 \quad \mathbf{C2}: M_2}{\mathbf{C1}; \quad \mathbf{C2}: M_1 \otimes M_2} & \mathbf{C} \\ \\ & \frac{\mathbf{C1}: M_1 \quad \mathbf{C2}: M_2}{\mathbf{if} \quad \mathbf{b} \quad \mathbf{then} \quad \mathbf{C1} \quad \mathbf{else} \quad \mathbf{C2}: M_1 \oplus M_2} & \mathbf{I} \\ \\ & \forall i, M^*_{ii} = m \quad \frac{\mathbf{C}: M}{\mathbf{loop} \quad \mathbf{X}_{\ell} \{\mathbf{C}\}: M^* \oplus \{^p_{\ell} \to j \mid \exists i, M^*_{ij} = p\}} & \mathbf{L} \\ \\ & \forall i, M^*_{ii} = m \quad \mathrm{and} \quad \forall i, j, M^*_{ij} \neq p \quad \frac{\mathbf{C}: M}{\mathbf{while} \quad \mathbf{b} \quad \mathbf{do} \quad \{\mathbf{C}\}: M^*} & \mathbf{W} \end{split}$$

Figure 1: mwp-bounds flow analysis inference rules

I Introduction

I.1 mwp-Bounds Analysis

Blah [1]

II References

[1] Neil D. Jones and Lars Kristiansen. "A flow calculus of *mwp*-bounds for complexity analysis". In: *ACM Transactions on Computational Logic* 10.4 (Aug. 2009), 28:1–28:41. DOI: 10.1145/1555746.1555752.