

Time complexity analysis for problem 1:

Given an array of size $n \times n$ and array of size n , the time complexity is $O(n^2)$. This is because we need to loop through each element in each row of A exactly once, and multiply that by each element in x . For an individual row in A , n things are multiplied together, giving $O(n)$ time complexity. You do this n times, so you get $n \cdot O(n) = O(n^2)$ time complexity. Basically, you loop through two for loops each n times. $n \cdot n = n^2$.

The derivation of the algorithm, and its time complexity analysis for problem 3:

The derivation of the algorithm for number 3 basically came from the fact that you can recursively define Hadmat multiplication of size n in terms of multiplication by a Hadmat matrix of size $n/2$ and the first half of x (also size $n/2$). This observation allows you to divide and conquer: Write $T(n)$ in terms of $T(n/2)$. We get, through this, that $T(n) = 2T(n/2) + O(n)$, which as we have seen from mergesort, gives $T(n) = O(n \log n)$.

Your comments on the trends that you observe for problem 5:

As expected, `matmult` grows at a polynomial rate, backing up the claim that it is $O(n^2)$. Likewise, `hadmatmult` grows at a slower rate, backing up the claim that it is $O(n \log n)$.