

Assessment 1

Logic,
Regular Languages,
and
Finite Automata

Norbert Logiewa
nl253

November 2017

Propositional Logic

1. Model as a propositional formal the poroperty that the inital state cannot be reached from c or b

Answer:

$$model \rightarrow \neg((c \rightarrow a') \vee (b \rightarrow a'))$$

2. Prove the following sequent:

$$a \rightarrow b'. b \rightarrow c'. c \rightarrow b' \vdash (a \vee b) \rightarrow (b' \vee c')$$

1. $a \rightarrow b'. b \rightarrow c'. c \rightarrow b'$ [premise]

2. $(a \rightarrow b') \wedge (b \rightarrow c') \wedge (c \rightarrow b')$ [\wedge_i] 1

3. $(a \rightarrow b') \wedge (c \rightarrow b') \wedge (b \rightarrow c')$ [commutativity of \wedge] 2

4. $(a \wedge c \rightarrow b') \wedge (b \rightarrow c')$ [distributivity of \rightarrow] 3

5. $a \wedge c \rightarrow b'$ [\wedge_{e2}] 4

6. $a \rightarrow b'$ [\wedge_{e2}] 5

7. $a \vee b \rightarrow b'$ [\vee_{i2}] 6

8. $(a \vee b) \rightarrow (b' \vee c')$ [\vee_i] 7 [conclusion]

3. Model the property that the automate can only be in one stat at a time via a defintion of a propositional function

Answer:

$$A \oplus B \equiv (A \vee B) \wedge (\neg A \vee \neg B) \text{ [definiton of the XOR operator]}$$

$$goodState(a, b, c) \equiv (a \rightarrow b') \oplus (b \rightarrow c') \oplus (c \rightarrow b')$$

4. Negate the formula and convert it to CNF

$$((a \rightarrow b') \wedge (a \wedge \neg b) \wedge (b' \wedge \neg a')) \rightarrow (b \rightarrow \neg a')$$

Answer:

1. Negate the formula

$$\neg(((a \rightarrow b') \wedge (a \wedge \neg b) \wedge (b' \wedge \neg a')) \rightarrow (b \rightarrow \neg a'))$$

2. Remove \rightarrow using $A \rightarrow B \equiv \neg A \vee B$ property

$$\neg(((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \rightarrow (\neg b \vee \neg a'))$$

$$\equiv$$

$$\neg(\neg((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \vee (\neg b \vee \neg a'))$$

3. Push \neg inwards using De Morgan's Laws

$$\begin{aligned}
 & ((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \wedge \neg(\neg b \vee \neg a') \\
 & \quad \equiv \\
 & ((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \wedge (b \wedge a')
 \end{aligned}$$

4. Distribute

$$\begin{aligned}
 & b \wedge ((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \wedge a' \wedge ((\neg a \vee b') \wedge (\neg a \vee \neg b) \wedge (b' \wedge \neg a')) \\
 & \quad \equiv \\
 & b \wedge (\neg a \vee b') \wedge b \wedge (\neg a \vee \neg b) \wedge b \wedge (b' \wedge \neg a') \\
 & \quad \wedge a' \wedge (\neg a \vee b') \wedge a' \wedge (\neg a \vee \neg b) \wedge a' \wedge (b' \wedge \neg a')
 \end{aligned}$$

5. Simplify (remove duplicate elements because $A \wedge A \equiv A$)

$$(\neg a \vee b') \wedge b \wedge (b' \wedge \neg a') \wedge (\neg a \vee \neg b) \wedge a'$$

5. Explain the principle of unit propagation from DPLL, and use it to show that CNF formula of the previous question is unsatisfiable.

1. **Answer:**

Unit propagation refers to a technique that is used to simplify logical formulas. When we have a clause that *only* contains a single literal such as x , in all other clauses that contain x , we can replace it with *True*.

2. **Answer:**

We can use it to show that the previous formula is unsatisfiable by demonstrating that it is a contradiction. We simplify it and continue to substitute *True* for single literal clauses. When we rearrange the formula, it becomes clear that we have $b \wedge \neg b$, this is a contradiction.

$$\begin{aligned}
 & (True \vee b') \wedge True \wedge \neg b \wedge b' \wedge \neg a' \wedge b \\
 & (True \vee True) \wedge True \wedge \neg b \wedge True \wedge \neg a' \wedge b \\
 & \quad \neg b \wedge True \wedge \neg a' \wedge b \\
 & \quad \neg b \wedge \neg a' \wedge b \\
 & \quad \neg b \wedge b \wedge \neg a'
 \end{aligned}$$

6. Comment on the validity of the original formula.

Answer:

The formula is not satisfiable and therefore not valid as satisfiability is a prerequisite for validity. In other words, because that logical statement was a contradiction (*assumed something can be False and True at the same time*), there was no way in which it could have been true and because of that it is not valid.

First Order Logic

1. Write a formula in first-order logic, using the parent relation, that states that two entities x and y are siblings if they share a parent

Answer:

$$\forall i. \forall j. \text{siblings}(i, j) \equiv (\text{parent}(x, i) = \text{parent}(x, j))$$

2. Assuming there is an equality relation \equiv on P such that $x \equiv y$ means x and y are the same entity, write a formula in first order logic stating that every entity has two distinct parents.

Answer:

$$\begin{aligned} \exists! x. P(x) &\equiv \exists x (P(x) \wedge \forall y (P(y) \rightarrow y = x)) \text{ }^1 \\ \forall i. \exists! y. \exists! x (\text{parent}(x, i) \wedge \text{parent}(y, i) \wedge \neg(x \equiv y)) \end{aligned}$$

3. Prove that the following sequent is valid:

$$\text{parent}(p, q). \forall x. \forall y (\text{parent}(x, y) \rightarrow \exists z. \text{parent}(z, x)) \vdash \exists z. \text{parent}(z, p)$$

-
1. $\text{parent}(p, q). \forall x. \forall y (\text{parent}(x, y) \rightarrow \exists z. \text{parent}(z, x))$ [premise]
 2. —
 3. $\exists z. \text{parent}(z, p)$ [conclusion] 1 – ?
-

4. Interpret the following formula into simple English statement

$$\forall x. \exists y. \text{parent}(y, x) \rightarrow \forall a. \exists b. \text{parent}(a, b)$$

Answer:

If all children have a parent, then all parents have a child.

5. The sequent below is not valid. Show this by proving a counter-example: a concrete definition of the universe P (a set) and the relation parent. Explain why this is a counterexample.

$$\vdash \forall x. \exists y. \text{parent}(y, x) \rightarrow \forall a. \exists b. \text{parent}(a, b)$$

In English

if all children have a parent then all parents have a child.

1. $\forall x. \exists y. \text{parent}(y, x) \rightarrow \forall a. \exists b. \text{parent}(a, b)$ [premise]

-
2. $P = \{ x \mid x \}$

-
3. *False* [conclusion]
-

¹<https://math.stackexchange.com/questions/228285/how-can-i-get-the-negation-of-exists-unique-existential-quantification>

6. What is the smallest possible counterexample universe P and relation parent . Explain your reasoning.

7. Prove that the following sequent is valid.

$$\forall x.\forall y.\text{parent}(x,y) \rightarrow \neg(x \equiv y) \vdash \neg\exists x.\text{parent}(x,x)$$

ie. if someone is a parent and has a child, then that child cannot be it's parent.

1. $\forall x.\forall y.\text{parent}(x,y) \rightarrow \neg(x \equiv y)$ [premise]

2. $\text{parent}(x_i, y_j) \rightarrow \neg(x_i \equiv y_j) \quad \forall_e 1$

3. $\neg\text{parent}(x_i, y_j) \vee \neg(x_i \equiv y_j)$ because $P \rightarrow Q \equiv \neg P \vee Q$

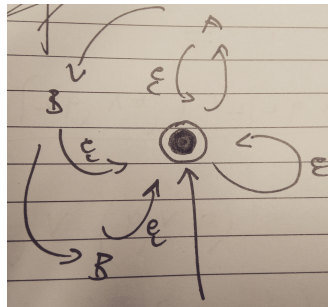
4. $\neg(\text{parent}(x_i, y_j) \wedge (x_i \equiv y_j))$ De Morgan's Law

5. $\neg\exists x.\text{parent}(x,x)$ [conclusion]

Regular Languages and Finite Automata

1. For the regular expression $(A|AB|ABB)^*$ do the following steps (and justify them):

• Answer:

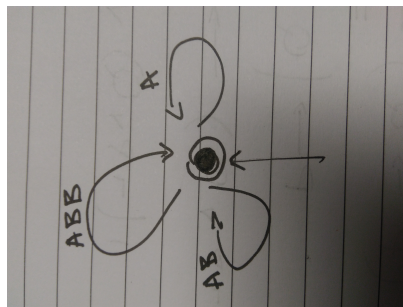


remove (semantic preserving) the ϵ transitions

Answer:

create the corresponding DFA

Answer:



2. Consider the following language:

A word is in the language if and only if it contains an even number of A s an odd number of B s and precisely one C . Design a DFA for this language. Explain your design (what information do you “store” in your states?)

Answer:

Not possible.

Answer:

Not possible.

References

- <https://en.wikipedia.org/wiki/Negation>
- https://en.wikipedia.org/wiki/Distributive_property
- https://en.wikipedia.org/wiki/Conjunctive_normal_form
- https://en.wikipedia.org/wiki/De_Morgan%27s_laws
- https://en.wikipedia.org/wiki/Modus_ponens
- https://en.wikipedia.org/wiki/Exclusive_or
- https://en.wikipedia.org/wiki/Finite-state_machine
- https://en.wikipedia.org/wiki/Uniqueness_quantification
- <https://math.stackexchange.com/questions/228285/how-can-i-get-the-negation-of-exists-unique-ex>