



Definition

A *nondeterministic finite automata* $N = (Q, \Sigma, \delta, q_0, F)$ is a list where Q and Σ are finite sets, $\delta : Q \times (\Sigma \cup \{\emptyset\}) \rightarrow \mathcal{P}(Q)$, $q_0 \in Q$ and $F \subset Q$.

As with finite automata, we call Q the *states*, Σ the *alphabet*, δ the *transition function*, q_0 the *start state*, and F the *accept states* (or *final states*). An input $u \in \mathbf{str}(\Sigma)$ results in a state sequence $x \in \mathbf{str}(Q)$ with $x_1 = q_0$ and $x_{i+1} = \delta(x_i, u_i)$ for $i = 1, \dots, |u|$.

Main result

For any automata M , there exists a nondeterministic finite automata N such that N accepts the same languages as M .

