



## Why

Suppose we want to fill up a backpack by selecting from various objects which gives us differing amounts of comfort. We only have so much space in the backpack.

## Definition

Number the objects  $1, \dots, n$ . Model the amount of space needed for a subset  $H \subset \{1, \dots, n\}$  of the  $n$  items by  $s(H)$ ; here  $s : \mathcal{P}(S) \rightarrow \mathbf{R}_+$ . Model the comfort (or value) they provide by  $v(H)$ ; here  $v : \mathcal{P}(\{1, \dots, n\}) \rightarrow \mathbf{R}$ . Given a space constraint  $c$ , we want to find  $H \subset P$  to

$$\begin{aligned} & \text{maximize} && v(H) \\ & \text{subject to} && s(H) \leq c \end{aligned}$$

In other words, find the subset of items which will fit in the bag and maximize the value. Such problems are called *knapsack problems*.

## Linear formulation

It is natural to expect the space constraint to be additive. In other words  $H \cap G = \emptyset$  means  $s(H \cup G) = s(H) + s(G)$ , from which we conclude that the function is monotonic (using the fact that it is nonnegative). I.e., given  $G \subset H$ , we have  $s(G) \leq s(H)$ . Suppose we also model the value function as additive. Given  $H \cap G = \emptyset$ , then  $v(H \cup G) = v(H) + v(G)$ .

It turns out that additivity is sufficient to have a linear representation for both  $s$  and  $v$ . For  $H \subset P$ , denote by  $\chi_H : P \rightarrow \{0, 1\}$  the characteristic function of  $H$ . In other words,  $\chi_H$  is defined by

$$x_H(i) = \begin{cases} 1 & \text{if } i \in H \\ 0 & \text{otherwise} \end{cases}$$

Then there exists  $p : \{1, \dots, n\} \rightarrow \mathbf{R}$  and  $w : \{1, \dots, n\} \rightarrow \mathbf{R}_+$  such that

$$v(H) = \sum_{i=1}^n p(i) \chi_H(i) \quad \text{for all } H \subset \{1, \dots, n\}$$

and

$$s(H) = \sum_{i=1}^n w(i)\chi_H(i) \quad \text{for all } H \subset \{1, \dots, n\}$$

We can formulate the following problem: given  $c \in \mathbf{R}_+$ ,  $w : \{0, 1\} \rightarrow \mathbf{R}_+$  and  $p : \{0, 1\}^n \rightarrow \mathbf{R}$ , find  $H \subset \{1, \dots, n\}$  to

$$\begin{aligned} & \text{maximize} && \sum_{i=1}^n p(i)\chi_H(i) \\ & \text{subject to} && \sum_{i=1}^n w_i\chi_H(i) \leq c \end{aligned}$$

It is common to identify  $\chi_H$  with a list  $x \in \{0, 1\}^n$  and to find  $x$  to

$$\begin{aligned} & \text{maximize} && \sum_i p_i x_i \\ & \text{subject to} && \sum_i w_i x_i \leq c \\ & && x_i \in \{0, 1\} \text{ for all } i = 1, \dots, n \end{aligned}$$

This problem is often called the *zero-one knapsack problem* (or *0-1 knapsack problem*). The *problem data* is the triple  $(p, w, c)$ .

### Alternative perspectives

Suppose instead the  $n$  objects are investments, the  $i$ th investment requiring  $w_i$  investment, returning  $p_i$ . Given that we have  $c$  dollars in capital, how should we allocate the funds to the investments to maximize return.

Other areas in which these problems are used as models include capital budgeting, cargo loading



## Terminology

We generally refer to the  $n$  *items*, the  $i$ th such item having *weight*  $w_i$  and profit  $p_i$ . Sometimes such problems are called *single knapsack problems* (one container), in contrast with *multiple knapsack problems* (in which there are multiple containers).

