

Optimal Tree Density Approximators

1 Why

Which is the optimal tree to use for tree density approximation?

2 Definition

We want to choose the tree whose corresponding approximator for the given density achieves minimum relative entropy with the given density among all tree density approximators. We call such a density an *optimal tree approximator* of the given density. We call a tree according to which an optimal tree approximator factors and optimal approximator tree.

3 Result

Proposition 1. Let $g : \mathbb{R}^n \to [0,1]$ be a density. A tree T on $\{1,\ldots,n\}$ is an optimal approximator tree if and only if it is a maximal spanning tree of the differential mutual information graph of q.

Proof. First, denote the optimal approximator of g for tree T by f_T^* . Recall

$$f_T^* = f_1 \prod_{i
eq 1} f_{i|\mathsf{pa}_i}$$

Second, recall d(g, f) = H(g, f) - H(g). Since H(g) does not depend on f, f is a minimizer of d(g, f) if and only if it is a minimizer of H(g, f).

Third, express the cross entropy of f_T^* relative to g as

$$\begin{split} H(q,p_T^*) &= h(q_1) - \sum_{j \neq i} \left(\int_{\mathbf{R}^d} g(x) \log g_{i|pai}(x_i, x_{\mathbf{pa}_i}) dx \right) \\ &= H(q_1) - \sum_{i \neq 1} \sum_{a \in A} q(a) (\log q_{i,\mathbf{pa}_i}(a_i, a_{\mathbf{pa}_i}) - \log q_{\mathbf{pa}_i}(a_{\mathbf{pa}_i})) \\ &= H(q_1) - \sum_{i \neq 1} \sum_{a \in A} q(a) (\log q_{i,\mathbf{pa}_i}(a_i, a_{\mathbf{pa}_i}) - \log q_{\mathbf{pa}_i}(a_{\mathbf{pa}_i}) - \log q_i(a_i) + \log q_i(a_i)) \\ &= \sum_{i = 1}^n H(q_i) - \sum_{i \neq 1} I(q_i, q_{\mathbf{pa}_i}) \\ &= \sum_{i = 1}^n H(q_i) - \sum_{\{i, j\} \in T} I(q_i, q_j) \end{split}$$

where \mathbf{pa}_i denotes the parent of vertex i in T (i = 2, ..., n). $H(g_i)$ does not depend on the choice of tree. Choosing a tree to minimize the second term in the final expression above is equivalent to choosing a maximal spanning tree from the weighted graph with differential mutual information edge weights; namely, the mutual information graph of q.