



Why

We want language and notation involving order.¹

Comparisons

A *comparison* is a statement (see **Statements**) involving a partial (which may or may not be total) order.

Notation

Let A be a set. We tend to denote an arbitrary partial order on A by \preceq . So (A, \preceq) is a partially ordered set.

As usual (see **Relations**), we write $a \preceq b$ to mean $(a, b) \in A$. Alternatively, we write $b \succeq a$ to mean $a \preceq b$. In other words, \succeq is the inverse relation (see **Converse Relations**) of \preceq .

Predecessors and successors

If $a \preceq b$ and $a \neq b$, we write $a \prec b$ and say that a *precedes* b . In this case we call a the *predecessor* of b . Alternatively, under the same conditions, we write $b \succ a$ and we say that b *succeeds* a . In this case we call b the *successor* of a .

Induced partial orders

Of course, the object we have defined and denoted by \prec is a relation on A . It satisfies (i) for no elements x and y do $x \prec y$ and $y \prec x$ hold simultaneously and (ii) if $x \prec y$ and $y \prec z$, then $x \prec z$ (i.e., \prec is transitive). It is worthwhile to observe that if S is a relation satisfying (i) and (ii), then the relation R defined to mean $(a, b) \in S$ or $a = b$ is a partial order on A .

¹In the present edition, this sheet can be thought of as an extended notation section for **Orders**.

Strict and weak relations

This connection between \preceq and \prec can be generalized. The *strict relation* corresponding to a relation R on a set A is the relation S on A defined by $(a, b) \in S$ if $(a, b) \in R$ and $a \neq b$. The *weak relation* corresponding to a relation S' on a set A is the relation R' defined by $(a, b) \in R'$ if $(a, b) \in S'$ or $a = b$. For this reason, a relation is said to *partially order* a set if it is a partial order or if its corresponding weak relation is one.

