Best Tree Distribution Approximators

1 Why

Which is the best tree to use for tree distribution approximation?

2 Definition

tes

We want to choose the tree whose corresponding approximator for the given distribution achieves minimum relative entropy with the given distribution among all tree distribution approximators. We call such a distribution an *optimal tree approximator* of the given distribution. We call a tree according to which an optimal tree approximator factors and optimal approximator tree.

Result

Proposition 1. *Let $A_1, \dots, A_n$ be finite non-empty sets. Define $A = \prod_{i=1}^{n} A_i$. Let $q : A \to [0, 1]$ a distribution. A tree T on $\{1, \dots, n\}$ is an optimal approximator tree if and only if it is a maximal spanning tree of the mutual information graph of q.*

Proof. First, denote the optimal approximator of q for tree T by p^*_T. Recall

$$p^*_T = q_1 \prod_{i \neq 1} q_{i|\mathsf{pa}_i}$$

Second, recall $d(q, p) = H(q, p) - H(q)$. Since $H(q)$ does not depend on p, p is a minimizer of $d(q, p)$ if and only if it is a minimizer of $H(q, p)$.

Third, express the cross entropy of p_T^* relative to q as

$$
\begin{aligned}
H(q, p_T^*) &= H(q_1) - \sum_{i \neq 1} \sum_{a \in A} q(a) \log q_{i|pai}(a_i, a_{\mathbf{pa}_i}) \\
&= H(q_1) - \sum_{i \neq 1} \sum_{a \in A} q(a) (\log q_{i,\mathbf{pa}_i}(a_i, a_{\mathbf{pa}_i}) - \log q_{\mathbf{pa}_i}(a_{\mathbf{pa}_i})) \\
&= H(q_1) - \sum_{i \neq 1} \sum_{a \in A} q(a) (\log q_{i,\mathbf{pa}_i}(a_i, a_{\mathbf{pa}_i}) - \log q_{\mathbf{pa}_i}(a_{\mathbf{pa}_i}) - \log q_i(a_i) + \log q_i(a_i)) \\
&= \sum_{i=1}^{n} H(q_i) - \sum_{i \neq 1} I(q_i, q_{\mathbf{pa}_i}) \\
&= \sum_{i=1}^{n} H(q_i) - \sum_{\{i,j\} \in T} I(q_i, q_j)
\end{aligned}
$$

where $\mathbf{pa}_i$ denotes the parent of vertex i in T $(i = 2, \ldots, n)$. $H(q_i)$ does not depend on the choice of tree. Choosing a tree to minimize the second term in the final expression above is equivalent to choosing a maximal spanning tree from the weighted graph with mutual information edge weights; namely, the mutual information graph of q.

$\square$