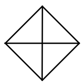


Why

Can we generalize the idea of flash codes.[1]

Definition

Let X be a set and A be an alphabet set. We denote the set of all finite sequences (strings) in A by $\mathsf{str}(A)$. We read $\mathsf{str}(A)$ aloud as "the strings in A." The length zero string is $\varnothing$.

A *code* for X *in* A is a function from X to $\mathsf{str}(A)$. In this context, we refer to the finite set A as an *alphabet* and we call $c(x)$ the *codeword* of x. The *length* of $x \in X$, with respect to a code $c : X \to \mathsf{str}(A)$, is the length of the sequence $c(x)$ (its codeword). We call a code *nonsingular* if it is injective.

Examples

Define $c : \{\alpha, \beta\} \to \{0, 1\}$ by $c(\alpha) = (0,)$ and $c(\beta) = (1,)$.[2]

Code extensions

Let $s, t \in \mathsf{str}(A)$ of length m and n respectively. The *concatenation* of s with t is the length $m+n$ string $u \in \mathsf{str}(A)$ defined by $u_1 = s_1, \dots, u_m = s_m$ and $u_{m+1} = t_1, \dots, u_{m+n} = t_n$. We denote the concatenation of s and t by st. Note, however, that $st \neq ts$, although $s(tr) = (st)r$.

Given a code $c : X \to \mathsf{str}(A)$, we can produce a code for $\mathsf{str}(X)$ in a natural way. The *extension* of c is the function $C : \mathsf{str}(X) \to \mathsf{str}(A)$ defined, for $\xi = (\xi_1, \dots, \xi_n) \in \mathsf{str}(X)$, by

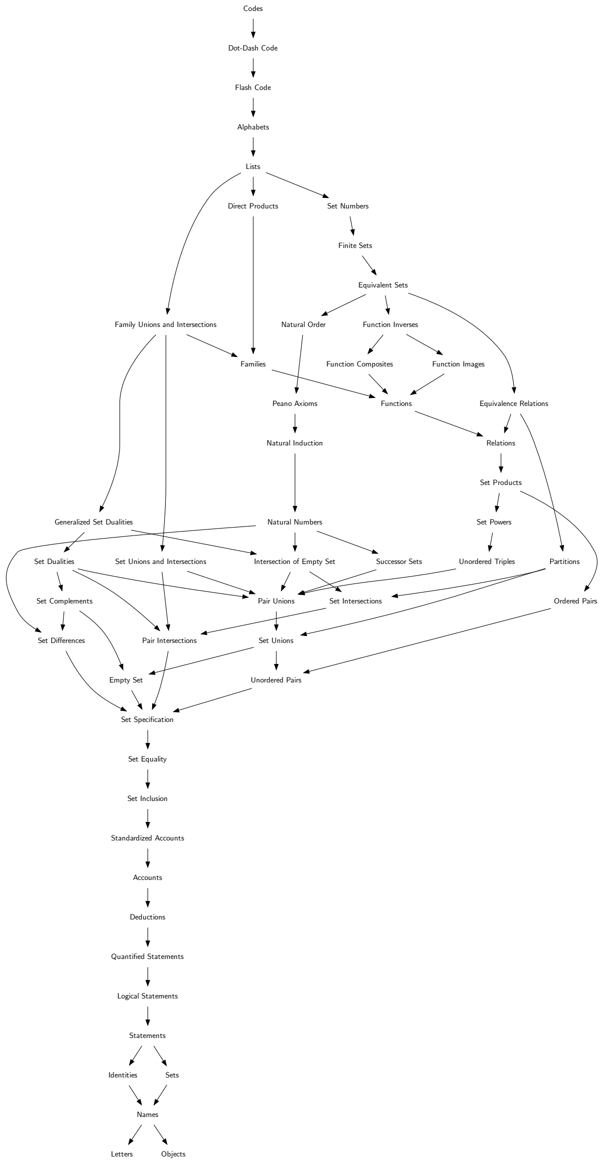
$$C(\xi) = c(\xi_1) \cdots c(\xi_n).$$

We call an code *uniquely decodable* if its extension is injective. In other

[1] The reliance of this sheet on Flash Codes and Dot-Dash Codes is for this justification, and not for any of the terms presented.

[2] Future editions will include additional examples.

words, given the code $C(\xi)$ for a sequence $\xi \in \mathsf{str}(X)$, we can recover ξ. We call $C(\xi)$ the *encoding* of ξ. We call ξ the *decoding* of $C(\xi)$.


Codes
Dot-Dash Code
Flash Code
Alphabets
Lists
Direct Products
Set Numbers
Finite Sets
Equivalent Sets
Family Unions and Intersections
Natural Order
Function Inverses
Families
Function Composites
Function Images
Peano Axioms
Functions
Equivalence Relations
Natural Induction
Relations
Set Products
Generalized Set Dualities
Natural Numbers
Set Powers
Set Dualities
Set Unions and Intersections
Intersection of Empty Set
Successor Sets
Unordered Triples
Partitions
Set Complements
Pair Unions
Set Intersections
Ordered Pairs
Set Differences
Pair Intersections
Set Unions
Empty Set
Unordered Pairs
Set Specification
Set Equality
Set Inclusion
Standardized Accounts
Accounts
Deductions
Quantified Statements
Logical Statements
Statements
Identities
Sets
Names
Letters
Objects