Bidirectionalizing ATL within Echo

Nuno Macedo Alcino Cunha





Universidade do Minho

FATBIT Workshop 2013 October 3, Braga, Portugal

ATL

- ATL is a widely used model transformation language;
- Has a de facto operational semantics implemented as an Eclipse plugin (denoted by t);
- However, it is unidirectional
 - Specification only states how to create a model N from a model M.

ATL and Echo

- The prescribed method to achieve bidirectionality is to define two transformations
 - Correctness and maintainability problem;
- ATL was created to answer the original QVT RFP
 - It shares some characteristics with the QVT standard;
- We propose the bidirectionalization of (the declarative subset of) ATL transformations with Echo.

ATL Language

• ATL main constituents are rules:

```
[[unique] lazy] rule R {
	from a: A (\pi_M)
	to b: B (\phi)
}
```

- Induce implicit traceability links.
- Bindings may rely on implicit rule calls.

Example

Person name : String Employee name : String salary : Int

```
module employ;
create OUT : Company from IN : World;
//PersonToEmployee
rule P2E {
  from
    p : World!Person ()
  to
    e : Company!Employee (
        name <- p.name
    )
}</pre>
```

ATL Semantics

- Disregards existing target model;
- Matched rules:
 - ullet For every source a that matched π_M create a b with bindings ϕ
- Unique laze rules:
 - Only run if called by other rules.

Bidirectional framework

- Lenses $(\overrightarrow{\mathbf{t}}(\overleftarrow{\mathsf{T}}(m,n)) = n \text{ and } \overleftarrow{\mathsf{T}}(m,\overrightarrow{\mathbf{t}}(m)) = m)$?
 - Infer relation such that $T(m, n) \equiv n = \overrightarrow{t}(m)$;
 - Would fail for *non-surjective* transformations;
 - Could relax the laws (to those of GRoundTram);
 - Applying \overrightarrow{t} would still disregard updates on n.
- Maintainers?
 - Both \overrightarrow{T} and \overrightarrow{T} derived from the inferred T;
 - f t would be run once to generate the initial target model..
 - ... and and would maintain them consistent;
 - Behavior of \overrightarrow{t} and \overrightarrow{T} converge for surjective transformations.

Consistency relation

- Source patterns are domain patterns (OCL predicates);
- Bindings are where constraints (post-conditions);
- Forall-there-exists tests do not suffice:
 - there must exist exactly one match;
- Requires traceability links.

Example

• E.g., **P2E**:

```
P2E<sub>▶</sub> (in: World, out: Company) ≡
\exists \ P2E_{\diamondsuit} \subseteq \mathsf{Person} \times \mathsf{Employee} \mid
\forall \ p: \mathsf{Person} \mid \exists \ ! \ e: \mathsf{Employee} \mid
\mathsf{P2E}_{\diamondsuit} \ (p,e) \land p.\mathsf{name} = e.\mathsf{name} \quad \land
\forall \ e: \mathsf{Employee} \mid \exists \ ! \ p: \mathsf{Person} \mid
\mathsf{P2E}_{\diamondsuit} \ (p,e) \land p.\mathsf{name} = e.\mathsf{name}
```

Unique lazy rules

- Are one-to-one but only if called;
- Existence is relaxed from the definition...

$$R_{\blacktriangleright} (m: M, n: N) \equiv \exists R_{\Leftrightarrow} \subseteq A \times B \mid A : A, b : B \mid R_{\Leftrightarrow} (a, b) \land \pi_M \Rightarrow \phi$$

... and is rather enforced by the rule call

$$\exists ! b : B \mid R_{\triangleleft \triangleright} (e, b) \land \forall a : A \mid R_{\triangleleft \triangleright} (a, b) \Rightarrow a = e$$

Implicit calls

- Our implementation also relies on traceability links R_⋄;
- ATL states the domain of rules must be disjoint;
- Matched target element can be retrieved.

Echo

- Higher-order existential quantifications are implemented in Alloy through skolemization;
- ATL is already deployed as an Eclipse plugin:
 - Like Echo, it is built over EMF;
- An embedding of the technique in Echo is under way.