

Partial Orders

A relation R on a set S is a partial order if it is reflexive, anti-symmetric and transitive.

eg. \geq is a partial order on \mathbb{Z}

$a \geq a$ for all $a \in \mathbb{Z}$, so reflexive

$a \geq b \wedge b \geq a \rightarrow a = b$, so anti-symmetric.

$a \geq b \wedge b \geq c \rightarrow a \geq c$, so transitive.

eg. "divides" (denoted by $|$) is a partial order on \mathbb{Z}^+ .

$a | a$ for all $a \in \mathbb{Z}^+$, so reflexive

$a | b \wedge b | a \rightarrow a = b$, so anti-symmetric

$a | b \wedge b | c \rightarrow a | c$, so transitive.

eg. \subseteq is a partial order on $\mathcal{P}(\mathbb{N})$

$A \subseteq A$ for all $A \in \mathcal{P}(\mathbb{N})$, so reflexive

$A \subseteq B \wedge B \subseteq A \rightarrow A = B$, so anti-symmetric

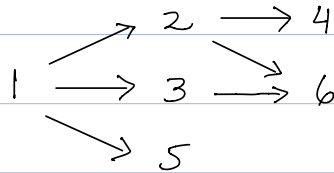
$A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$, so transitive.

Def A total order $R \subseteq A \times A$ is a partial order where for every $a, b \in A$, either $(a, b) \in R$ or $(b, a) \in R$.

Def a and b are incomparable if $(a, b) \notin R$ and $(b, a) \notin R$.

eg. \leq is a total order since $\forall a, b \in \mathbb{Z}$, $a \leq b$ or $b \leq a$.

eg. 'divides' is not a total order



Note: 2 and 3 are incomparable since $2 \nmid 3$ and $3 \nmid 2$.

eg. Let R be a relation over $\mathbb{Z} \times \mathbb{Z}$ where $(a, b) R (c, d)$ iff $a \leq c \wedge b \leq d$.

R is not total since $(1, 2)$ and $(2, 1)$ are incomparable.

eg. Let R' be a relation over $\mathbb{Z} \times \mathbb{Z}$

where $(a, b) R' (c, d)$ iff

$$a \leq c \vee (a = c \wedge b \leq d)$$

R' is a total order (lexicographic order)

$$(1, 2) R' (2, 1).$$

Def A strict partial order is
anti-reflexive and transitive.

Theorem A strict partial order is
anti-symmetric.

Proof: Suppose $a R b$ and $b R a$. By
transitivity, $a R a$. This contradicts
the fact that R is anti-reflexive.
So $\neg (a R b \wedge b R a)$. Therefore, R is
anti-symmetric.

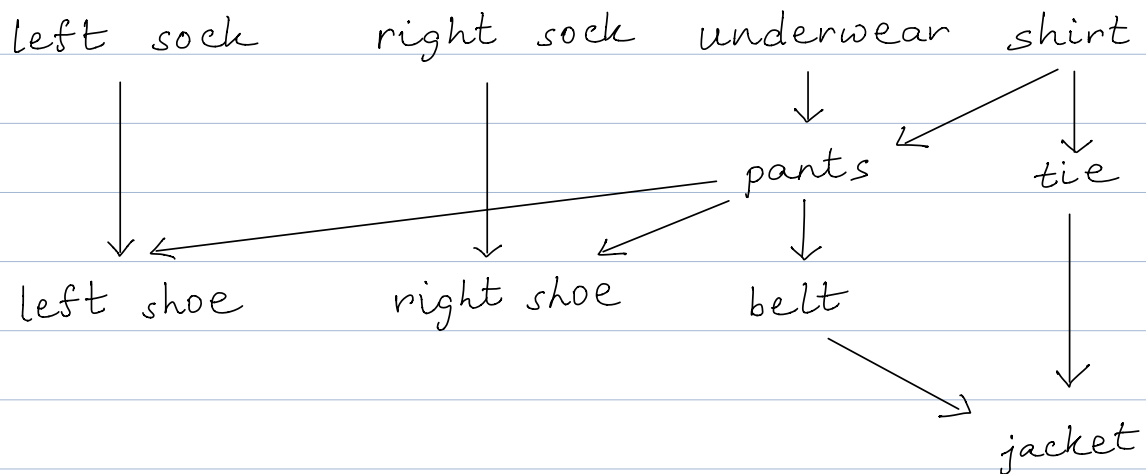
QED.

$<$ and \subset are strict partial orders.

A strict partial order can be represented
by a DAG (directed acyclic graph).

eg. a pre-requisite dependancy flowchart

eg.



We can solve scheduling problems by relying on a partial order.

Topological Sorting

come up with a total order of the items that is consistent with the partial order.

- ① underwear, shirt, pants, left sock, right sock, belt, tie, left shoe, right shoe, jacket.
- ② left sock, shirt, tie, underwear, right sock, pants, left shoe, belt, jacket, right shoe.