Limited Direct Execution

How does the OS keep control?

Limited Direct Execution

Time-sharing is the idea that multiple processes can run together on a single machine as though they are in sole control of the machine

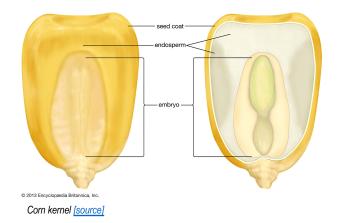
How do processes share time?

multiprogramming – when a process waits for I/O, the OS can have another take over the CPU

multitasking – each process gets a time-slice, a time limit before the next process gets to execute on the CPU

How does the OS keep control?

Hardware provides *interrupts*, *kernel mode* and *user mode*We will see how the OS uses these three mechanisms to implement CPU virtualization



Direct Execution

OS

- 1. Create entry for process
- 2. Allocate memory for program
- 3. Load program into memory
- 4. Set up stack (argc and argv parameters for main
- 5. Clear registers
- 6. Call main()

- 1. Free memory of process
- 2. Remove from process list

Program

1. Run main()

2. Execute return from main

What is wrong?

Process has unrestricted access to memory and resources.

OS has no way to switch another process, must wait for program to finish.

Restricting Access with Processor Mode

CPU has bit that indicates if in user mode or kernel mode

When in user mode, "privileged" instructions not allowed and memory boundaries are enforced

Cannot read/write outside of address space bounds Cannot read/write I/O devices

When in kernel mode, all instructions are allowed

Normal processes only execute in user mode, the OS executes in kernel mode

Processor Mode

OS Kernel Mode

Program User Mode

- 1. Create entry for process
- 2. Allocate memory for program
- 3. Load program into memory
- 4. Set up stack (argc and argv parameters for main
- 5. Clear registers
- 6. Call main()

enter user mode

enter kernel mode

- 1. Run main()
- 2. Execute return from main

- 1. Free memory of process
- 2. Remove from process list

How does mode change from user to kernel?

Can program change mode to kernel?

System Call

Problem: How can the program perform privileged operations, for example read from a file?

If no modes, user could just make a functional call to kernel libraries, easy!

A system call is different from a normal function call in that it transfers control to the OS.

A program initiates a system call by causing a *trap* (a software generated interrupt)

For example, the RISC V instruction that causes a trap is *ecall*

S Kernel Mode

System Call

Hardware

Program User Mode

Initialize process and return-from-trap into main()

Return-from-trap changes mode to user

Trap changes mode to kernel

- 1. Run main()
- 2. System call causes trap

- 1. Handle trap
 - ...
- 2. return-from-trap

Return-from-trap changes mode to user

Trap changes mode to kernel

- Call exit() system call
- 2. System call traps into OS

Free memory and remove process

System Call Details

OS

Kernel Mode

Hardware

User Mode Program

- Create entry for process 1.
- 2. Allocate memory for program
- 3. Load program into memory
- 4. Set up stack (argc and argv parameters for main)
- Clear registers 5.
- return-from-trap

1. Handle trap

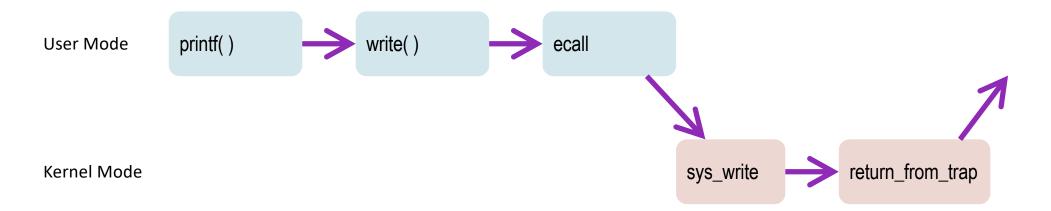
2. return-from-trap

- Restore regs from kernel stack
- Move to user mode
- Jump to main
- Save regs to kernel stack
- Move to kernel mode
- Jump to trap handler
- Run main()
- 2. Call system call
- System call traps into OS
- Restore regs from kernel stack
- Move to user mode
- Jump to PC after trap
- Run main()
- Call exit() system call 2.
- System call traps into OS

- Free memory of process 1.
- 2. Remove from process list

System Call Example

Library functions like printf typically have some code that executes in user mode, it then makes one or more system calls.



```
swtch.S
    # Context switch
 2
    #
    #
         void swtch(struct context *old, struct context *new);
 3
                                                                    load new context to registers
    #
 4
                                                                    Id = load doubleword command
    # Save current registers in old. Load from new.
 5
 6
 7
                                                            ld ra, 0(a1)
                                                   24
 8
     .globl swtch
                          store registers to old context
                                                            ld sp, 8(a1)
                                                   25
     swtch:
                          sd = store doubleword command<sub>2.6</sub>
                                                            ld s0, 16(a1)
10
         sd ra, 0(a0)
                                                            ld s1, 24(a1)
                                                   27
         sd sp, 8(a0)
11
                                                            ld s2, 32(a1)
                                                   28
         sd s0, 16(a0)
12
                                                            ld s3, 40(a1)
                                                   29
         sd s1, 24(a0)
13
                                                            ld s4, 48(a1)
                                                   30
14
         sd s2, 32(a0)
                                                            ld s5, 56(a1)
                                                   31
15
         sd s3, 40(a0)
                                                            ld s6, 64(a1)
                                                   32
         sd s4, 48(a0)
16
                                                            ld s7, 72(a1)
                                                   33
17
         sd s5, 56(a0)
                                                            ld s8, 80(a1)
                                                   34
18
         sd s6, 64(a0)
                                                            ld s9, 88(a1)
                                                   35
19
         sd s7, 72(a0)
                                                            ld s10, 96(a1)
                                                   36
         sd s8, 80(a0)
20
                                                            ld s11, 104(a1)
                                                   37
         sd s9, 88(a0)
21
                                                   38
         sd s10, 96(a0)
22
                                                            ret
                                                   39
23
         sd s11, 104(a0)
```