## CS161 Homework 3 Problem 3-6

(a)

$$X_i = \text{ the number of probes required for the ith insertion} \\ A_j = \begin{cases} 1 & \text{if } j \text{th slot is occupied} \\ 0 & \text{otherwise} \end{cases}$$

$$\Pr(A_j = 1) = \frac{i-1}{m}$$

$$\Pr(X_i = k) = \left(\frac{i-1}{m}\right)^{k-1} \left(1 - \frac{i-1}{m}\right)$$

$$\Pr(X_i > k) = 1 - \Pr(X_i \le k)$$

$$= 1 - \sum_{j=1}^k \left(\frac{i-1}{m}\right)^{j-1} \left(1 - \frac{i-1}{m}\right)$$

$$= 1 - \left(1 - \left(\frac{i-1}{m}\right)^k\right)$$

$$= \left(\frac{i-1}{m}\right)^k$$

$$\le \left(\frac{m-1}{m}\right)^k$$

$$\le \left(\frac{m-2}{2m}\right)^k$$

$$\le \left(\frac{m}{2m}\right)^k$$

$$= \left(\frac{1}{2}\right)^k$$

Thus,  $\Pr(X_i > k) \le \left(\frac{1}{2}\right)^k$ .

(b) From part a, we have  $P(X_i > k) \leq \left(\frac{1}{2}\right)^k$ . Plug in  $k = 2 \log n$  and we get:

$$\Pr(X_i > k) \leq \left(\frac{1}{2}\right)^k$$

$$\Pr(X_i > 2\log n) \leq \left(\frac{1}{2}\right)^{2\log n}$$

$$= 2^{-2\log n}$$

$$= 2^{\log n^{-2}}$$

$$= n^{-2}$$

Thus, we have that  $\frac{1}{n^2}$  is an upper bound for  $P(X_i > 2 \log n)$ , so  $P(X_i > 2 \log n) = O\left(\frac{1}{n^2}\right)$ 

(c)

$$Pr(X > k) = Pr(X_1 > k) \text{ or } Pr(X_2 > k) \text{ or } Pr(X_3 > k) \text{ or ... or } Pr(X_n > k)$$

$$= Pr(X_1 > k) + Pr(X_2 > k) + Pr(X_3 > k) + ... + Pr(X_n > k)$$

$$= \sum_{i=1}^{n} Pr(X_i > k)$$

In this case,  $k = 2 \log n$ , so we have:

$$\Pr(X > 2\log n) = \sum_{i=1}^{n} \Pr(X_i > 2\log n)$$

$$\leq \sum_{i=1}^{n} \frac{1}{n^2}$$

$$= n\left(\frac{1}{n^2}\right)$$

$$= \frac{1}{n}$$

Thus,  $\frac{1}{n}$  is an upper bound for  $\Pr(X > 2 \log n)$ , so  $\Pr(X > 2 \log n) = O\left(\frac{1}{n}\right)$ .

(d) We know  $\Pr(X > 2 \log n) = \frac{c}{n}$ , where c is some constant. That means  $\Pr(X \leq n)$ 

$$2\log n) = 1 - \frac{c}{n}.$$

$$\begin{split} E[X] &= \sum_{x} x \Pr(X = x) \\ &\leq \left( \max_{x > 2\log n} x \right) \frac{c}{n} + \sum_{x \leq 2\log n} x \Pr(X = x) \\ &\leq \left( \max_{x > 2\log n} x \right) \frac{c}{n} + \left( \max_{x \leq 2\log n} x \right) \left( 1 - \frac{c}{n} \right) \\ &= n \frac{c}{n} + (2\log n) \left( 1 - \frac{c}{n} \right) \\ &= n \frac{c}{n} + (2\log n) \left( \frac{n-c}{n} \right) \\ &= c + (2\log n) \left( \frac{n-c}{n} \right) \\ &\leq c + (2\log n) \end{split}$$

Thus,  $E[X] = O(\log n)$ .