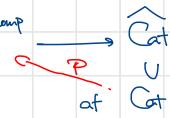


§ Cocompletion

Then $\text{LFun}(\mathcal{P}(\mathcal{C}), \mathcal{D}) \xrightarrow[\text{LKE}]{\perp} \text{Fun}(\mathcal{C}, \mathcal{D})$: equiv.

i.e. \mathcal{P} is a local adj to $\widehat{\text{Cat}}^{\text{cocomp}}$

(2)



This is a consequence of the following:

Rmk representable functors are atomic, i.e. $\text{Map}_{\text{Psh}(\mathcal{C})}(\mathcal{F}_X, -) : \text{Psh}(\mathcal{C}) \rightarrow \text{Ani-pres}$.

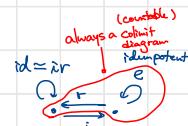
Def $\mathcal{C}_0 \hookrightarrow \mathcal{C}$ is dense if $\text{Lang}_f = \text{id}$, i.e. $\forall X \in \mathcal{C}$

Small colimits

$$\text{Colim } c \xrightarrow{\sim} X$$

$c \hookrightarrow X$
 $c \in \mathcal{C}_0$

ex $\mathbb{Q} \hookrightarrow \mathbb{R}$ is dense but $\mathbb{Z} \hookrightarrow \mathbb{R}$ is not. ($\mathbb{R} \xrightarrow{L} \mathbb{R}$)



Thm $\mathcal{C} \in \text{Cat} \rightarrow \mathcal{C} \xrightarrow{\perp} \mathcal{P}(\mathcal{C})$ is dense

exer Any atomic obj of $\mathcal{P}(\mathcal{C})$ is a retract of a representable. $\mathcal{C} \hookrightarrow \widehat{\mathcal{C}} = \widehat{\mathcal{P}(\mathcal{C})}^{\text{atomic}}$ is the idempotent completion

$$\begin{array}{ccccc} \mathcal{C} & \longrightarrow & \mathcal{P}(\mathcal{C}) & \xrightarrow{\quad} & \mathcal{D} \\ \downarrow & \nearrow \text{id} & & & \nearrow \\ \mathcal{P}(\mathcal{C}) & & \text{left Kan extended} & & \\ \downarrow & & \oplus & & \\ X & & F: \text{colim-pres} & & \end{array}$$

Rmk $\widehat{\text{Cat}}^{\text{cocomp}} \xrightleftharpoons[\perp]{\exists} \widehat{\text{Cat}}$ $\mathcal{C} \hookrightarrow \text{Fun}(\mathcal{C}^\text{op}, \widehat{\text{Ani}})$ & take the closure. If \mathcal{C} : locally small.
If it is a subset of $\text{Fun}(\mathcal{C}^\text{op}, \text{Ani})$. If moreover \mathcal{C} : accessible.

Prop $\mathcal{C} \hookrightarrow \mathcal{D}$ is dense iff the restricted Yoneda emb $\mathcal{D} \rightarrow \mathcal{P}(\mathcal{D}) \xrightarrow{\text{res}} \mathcal{P}(\mathcal{C})$ is f.f.

$$\begin{array}{ccc} \mathcal{C} & \longrightarrow & \mathcal{P}(\mathcal{C}) \xrightarrow{h_Y} \mathcal{D} \\ \downarrow & \text{RKE}_F \downarrow \text{id}_{\mathcal{P}(\mathcal{C})} & \uparrow \text{id} \\ \mathcal{C} \hookrightarrow \mathcal{D} & \longrightarrow & \mathcal{P}(\mathcal{D}) \ni h_Y \xrightarrow{\quad} \mathcal{C} \rightarrow \mathcal{D} \end{array}$$

$\mathcal{D} \xleftarrow{\quad} \mathcal{P}(\mathcal{C})$ refl. localization

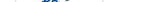
§ Reflective (aka. Bousfield) localization

Exer refl. loc \Leftrightarrow loc. functor w/ right adj.

Def. $\mathcal{C} \xrightarrow{\text{refl.}} \mathcal{D}$ is a localization if it admits a fully faithful right adjoint.

- $\text{• A subcategory } \mathbb{D} \subset \mathcal{C} \text{ is } \underline{\text{reflective}} \text{ if the incl. admits a left adj.}$

Prop(localization \Leftrightarrow idempotent monad) Let $L : \mathcal{C} \rightarrow \mathcal{C}$ be an endofunctor w/ ess. im $L\mathbb{I}$. TFAE:

- (1) \exists localization \mathcal{L}  & an equivalence $L \xrightarrow{\sim} f^* f_*$

- (2) $L: \mathcal{C} \rightarrow \mathcal{L}^{\mathbb{R}}$ is left adjoint to the inclusion $i: \mathcal{L}^{\mathbb{R}} \rightarrow \mathcal{C}$

- (3) \exists nat. tr $\eta : \text{id}_B \rightarrow L : \mathcal{C} \rightarrow \mathcal{C}$ s.t. $\forall c \in \mathcal{C} \quad \eta_c, L(\eta_c) : Lc \rightarrow LLc$ are equiv.

proof (1) \Leftrightarrow (2) replace \mathbb{D} by its essential image in \mathbb{K}

(2) \Rightarrow (3) η : unit of $L\rightarrow i$

$$\hookrightarrow \forall c_0 \in \mathcal{C}, c_i \in \mathcal{L}^c \quad \text{Map}(L_{c_0}, c_i) \xrightarrow{\eta_{c_0}^*} \text{Map}(c_0, c_i)$$

(3) \Rightarrow (2)

enough to show y is the unit

$$\text{i.e. } \text{Map}_{\mathcal{L}_B}(L_C, L_C) \xrightarrow{\sim} \text{Map}_{\mathcal{L}_C}(L_C, L_C)$$

• To-suri because $\eta_{Co} \downarrow$ $L_{Co} \xrightarrow{1f} L_{Co}$ $\xrightarrow{\forall f} L_{Co}$

$$\text{• } \underline{\text{mono}} \quad \text{Map}_g(\mathcal{L}\mathcal{C}_0, \mathcal{L}\mathcal{C}_1) \xrightarrow{\cong} \text{Map}_g(\mathcal{L}\mathcal{C}_0, \mathcal{L}\mathcal{C}_1)$$

$$\begin{array}{ccc}
 & \xrightarrow{\cong} & \xrightarrow{\cong} \\
 & \downarrow & \downarrow \\
 \text{Map}_\infty(L_{C_0}, L_{C_1}) & \longrightarrow & \text{Map}_\infty(C_0, L_{C_1}) \xrightarrow{\cong} f \\
 & \xrightarrow{\cong} \downarrow L & \downarrow L \\
 & & \cong \\
 \text{Map}_\infty(LL_{C_0}, LL_{C_1}) & \longrightarrow & \text{Map}_\infty(L_{C_0}, LL_{C_1}) \xrightarrow{\cong} Lf \\
 & \downarrow \cong & \downarrow \cong \\
 & & \downarrow \cong \\
 \text{Map}_\infty(LL_{C_0}, L_{C_1}) & \longrightarrow & \text{Map}_\infty(L_{C_0}, L_{C_1}) \xrightarrow{\cong} \eta_{L_{C_0}}^{-1} \cdot Lf
 \end{array}$$

$$\begin{array}{c} \text{Map}(Lc_0, c_1) \xrightarrow{\eta_{c_0}^*} \text{Map}(c_0, c_1) \\ \hookrightarrow \left\{ \begin{array}{l} c_0 = Lc \rightsquigarrow \eta_{Lc} : Lc \xrightarrow{\sim} LLc \text{ by Yoneda} \\ c_0 = c, c_1 = LLc \\ \text{naturality square} \end{array} \right. \end{array}$$

To- inj.
+
map is from
 $\forall k \in H$.

- If \overline{f} : colimit cone, Loft is a colimit cone of f
(because $\text{Loft}|_K = K$)

- \mathcal{C} is closed under limits in \mathbb{D}

$$\frac{d}{Ld} \rightarrow \underset{\textcircled{3}}{f}$$

$$\text{Map}_{\text{Fun}(k, \Phi)}(\underline{d}, \underline{\text{id}}) \cong \text{Map}_{\Phi}(\underline{d}, \underline{\lim^{\Phi} \text{id}})$$

$$\text{Map}_{\text{Fun}(k, e)}(\underline{\text{Id}}, f) \xrightarrow{\sim} \text{Map}_{\text{Fun}(k^Q, e)}(\underline{id}, \tilde{f})$$

limit if is L-loc

$\text{Map}_S \left(\stackrel{\text{Id}}{\sim} \underset{\text{def}}{\lim} f \right)$

ess. im of loc. can be characterized by the following:

Def Let $W \subset \text{Ar}(b)$. $c \in b$ is W -local if $\forall f: a \rightarrow b$,

$$\text{Map}_b(b, c) \xrightarrow{\sim} \text{Map}_b(a, c)$$

exer $D \stackrel{L}{\varprojlim}_i b$. $c \in \text{Im } i \Leftrightarrow$ local wrt L-equiv. Moreover, $D \cong b[L\text{-equiv}]$

Warning. $\{W\text{-loc obj}\}$ $\stackrel{\leftarrow}{\subseteq} \stackrel{\curvearrowright}{c}$ not nec. exist.
closed under lim

§ Presentable categories, compactness, AFT

Idea: presentable cats are cocomplete (\rightsquigarrow usually large) categories that is generated by small set of generators & relations.

algebra	higher category theory
(small) Set	(accessible) ∞ -category
sums	colimits
abelian groups	presentable ∞ -cat
free ab gps	presheaves
products	(finite) limits
column. ring	∞ -topos)

$$\text{Cat} \xrightarrow{\mathcal{P}} \widehat{\text{Cat}}^{\text{cocomplete}} \subset \widehat{\text{Cat}}$$

, a large cat non-full

Def \mathcal{C} is presentable if $\exists \mathcal{C}_0$ small, $R \subset \mathcal{P}(\mathcal{C}_0)$ (small set of morphisms).

$$\mathcal{C} = \mathcal{P}(\mathcal{C}_0; R) \xleftarrow{\subseteq} \mathcal{P}(\mathcal{C}_0) \quad \text{full subcat of } R\text{-local obj. i.e.}$$

\Downarrow

$$X \hookrightarrow \forall f: A \rightarrow B \in R, \text{Map}(B, X) \xrightarrow{\sim} \text{Map}(A, X).$$

$$\text{Pr}^L \subset \widehat{\text{Cat}}^{\text{cocomplete}} \quad \text{full subcat of presentable cats.}$$

fact • $\mathcal{C} \subset \mathcal{P}(\mathcal{C}_0)$ closed under lim. $\mathcal{C} \xleftarrow{\subseteq} \mathcal{P}(\mathcal{C}_0) \Rightarrow \mathcal{C}$ is also cocomplete.

$$\begin{array}{ccc} \coprod_{r \in R} \mathcal{P}(\Delta) & \xrightarrow{r} & \mathcal{P}(\mathcal{C}_0) \\ \downarrow & \lrcorner & \uparrow \lrcorner \\ \coprod_{r \in R} \mathcal{P}(\Delta) & \longrightarrow & \mathcal{P}(\mathcal{C}_0; R) \end{array}$$

in $\widehat{\text{Cat}}^{\text{cocomplete}}$ "presentation"

It turns out that $\text{Pr}^L \subset \widehat{\text{Cat}}^{\text{cocomplete}}$ closed under small colim

\rightsquigarrow it is a closure of presheaves under small colim

$\rightsquigarrow \text{Pr}^L$: category of \mathbb{D} -compact objects

For more effective we need to discuss K -filt. colim first.

K : infinite (regular) cardinal

Def • $K \in \text{sSet}$ is K -small if it has finitely many nondeg. simplices.

$\bigcup_{K \text{-small}}$ $\text{Set}_{\leq K}$ has K -small colim.

• $\mathcal{C} \in \widehat{\text{Cat}}$ is K -filtered if $\forall K \xrightarrow{f} \mathcal{C}$ ($\iff \forall (K, f) \mathcal{C}_{pf} \neq \emptyset$)

$$\begin{array}{c} \xrightarrow{f} \\ \downarrow \lrcorner \quad \lrcorner \nearrow \\ \mathcal{C}^K \subset \mathcal{C} \end{array} \quad \text{full sub}$$

* If \mathcal{C} has K -filt colim, $X \in \mathcal{C}$ is K -cpt if $\text{Map}_{\mathcal{C}}(X, -): \mathcal{C} \rightarrow \text{Ani}$ preserves K -filt colim.

• For \mathcal{C} small, $\text{Ind}_{\mathbb{K}}(\mathcal{C}) \subset \mathcal{P}(\mathcal{C})$ full sub of $F: \mathcal{C}^{\text{op}} \rightarrow \text{Ani}$ s.t. $\int F = \mathcal{C}_{/\mathbb{K}}$ is
 \uparrow
 \mathbb{K} -filtered (called the \mathbb{K} -flat functor)
 $\left(\begin{array}{l} \text{Rank } \mathcal{C} \text{ such that } F \simeq \text{colim}_{\mathbb{K} \in \mathbb{K}} \text{colim}_{\mathbb{K} \in \mathbb{K}} \text{rep} \\ \text{Prop: this is closed under } \mathbb{K}\text{-flat colimits [065Z]} \end{array} \right)$

Rem for \mathcal{C} large,
this is the def.

→ f. the inclusion pres \mathbb{K} -filt. colim

• $\mathcal{C} \hookrightarrow \text{Ind}_{\mathbb{K}}(\mathcal{C})$ factors through

\mathbb{K} -compact obj.

Prop $\forall \mathcal{D}$ with small \mathbb{K} -filt colim

$\text{Fun}^{\text{K-filt}}(\text{Ind}_{\mathbb{K}}(\mathcal{C}), \mathcal{D}) \xleftarrow[\sim]{\text{LKE}} \text{Fun}(\mathcal{C}, \mathcal{D})$.

$$\widehat{\text{Cat}} \xleftarrow[\sim]{\text{Ind}_{\mathbb{K}}} \widehat{\text{Cat}} \quad \begin{matrix} \text{Ind}_{\mathbb{K}} \\ \downarrow \text{Forget} \end{matrix}$$

Prop Consider $\mathcal{C} \xrightarrow{f} \mathcal{D}$ fully faithful, where \mathcal{D} has \mathbb{K} -filt colim

$$\xrightarrow{\sim} \text{Ind}_{\mathbb{K}}(\mathcal{C}) \xrightarrow{f} \mathcal{D}$$

(1) If $\mathcal{C} \subset \mathcal{D}^{\mathbb{K}}$, f is fully faithful

(2) If moreover \mathcal{C} generates \mathcal{D} under \mathbb{K} -filt colim (i.e. the closure of \mathcal{C} is \mathcal{D}) then $f: \mathcal{C} \rightarrow \mathcal{D}$ is fully faithful.

Proof over. (harder if \mathcal{C} : large) (reduce to when $\mathcal{D} = \text{Ani}^{\mathbb{K}}$ & use $\mathcal{P}(\mathcal{C})$)

ex many large cats we know, also flat modules $\text{rank } \mathcal{C}: (\mathbb{N}, 1)\text{-Cat} \rightarrow \text{Ind}_{\mathbb{K}}(\mathcal{C}): (\mathbb{N}, 1)\text{-Cat}$

Def \mathcal{C} : (\mathbb{K} -)accessible $\iff \exists \mathbb{K} \exists \mathcal{C}_0$ small. $\text{Ind}_{\mathbb{K}}(\mathcal{C}_0) \xrightarrow{\sim} \mathcal{C}$ (\hookrightarrow can take $\mathcal{C}_0 = \mathcal{C}^{\mathbb{K}}$)

A functor $f: \mathcal{C} \rightarrow \mathcal{D}$ is accessible if it is \mathbb{K} -contd for some \mathbb{K}

($\hookrightarrow f: \mathcal{C}, \mathcal{D}: \text{acc}, \exists \mathbb{T} f: \mathbb{T}$ -contd & pres. \mathbb{T} -cpt obj (HTT Rmk 5.4.2.13))

Warning \mathbb{K} -acc does not imply \mathbb{T} -acc for $\mathbb{T} > \mathbb{K}$. but it does if $\mathbb{T} \gg \mathbb{K}$. i.e. if $\mathbb{K} \subset \mathbb{T}$ $\mathcal{C}^{\mathbb{K}} \subset \mathcal{C}^{\mathbb{T}}$

\hookrightarrow If \mathcal{C} acc, $\mathcal{C}^{\mathbb{T}} = \text{small } \forall \mathbb{T} \quad \bigcup \mathcal{C}^{\mathbb{T}} = \mathcal{C}$

$\text{Acc}_{\mathbb{K}}$: \mathbb{K} -acc functors & \mathbb{K} -contd functors

$$\text{Acc}_{\mathbb{K}} \xleftarrow[\sim]{\text{LKE}} \widehat{\text{Cat}} \quad \begin{matrix} \text{Ind}_{\mathbb{K}} \\ \downarrow \text{idem} \\ \widehat{\text{Cat}} \end{matrix}$$

Rmk • Small cats are accessible iff idem cpt.

• $\text{Acc} = \bigcup \text{Acc}_{\mathbb{K}} \hookrightarrow \widehat{\text{Cat}}$ creates limits (in fact, Cat -weighted lim)

• adjoints between accessible cats are accessible.

Def \mathcal{C} : \mathbb{K} -compactly generated \iff cocomplete & \mathbb{K} -accessible.

$$\hookrightarrow \text{colim}_{\mathbb{K}} \mathcal{C} \hookrightarrow \text{colim}_{\mathbb{K}} \mathcal{C}$$

$\hookrightarrow \text{Ind}_{\mathbb{K}}(\mathcal{C}_0) \xrightarrow{\sim} \mathcal{C}$ for \mathcal{C}_0 : small with \mathbb{K} -small colim.

$$\hookrightarrow \mathcal{C}_0 \xrightarrow{\sim} \mathcal{P}(\mathcal{C}_0)^{\mathbb{K}} \hookrightarrow \text{Ind}_{\mathbb{K}}(\mathcal{P}(\mathcal{C}_0)^{\mathbb{K}}) = \mathcal{P}(\mathcal{C}_0)$$

$$\Pr_{\mathbb{K}}^L \cap \Pr_{\mathbb{K}}^L$$

Prop presentable \iff $\text{Acc. loc. of } \mathcal{P}(\mathcal{C}_0) \iff$ cocomplete & accessible.

- * A category has colimits iff it has small objects & preserves limits.
- I : K-filtered $\Leftrightarrow \text{colim} : \text{Fun}(I, \text{Ari}) \rightarrow \text{Ari}$ pres. K-small limits.

Cor K-compact objects are closed under K-small colimits

$$\begin{aligned} \text{Map}_\beta(\text{colim}_I X_i, \text{colim}_J Y_j) &\simeq \lim_I \text{Map}_\beta(X_i, \text{colim}_J Y_j) \\ &\stackrel{\text{K-small}}{\uparrow} \quad \stackrel{\text{K-filt}}{\downarrow} \\ &= \lim_I \text{colim}_J \text{Map}_\beta(X_i, Y_j) \\ &\leftarrow \text{colim}_J \lim_I \text{Map}_\beta(X_i, Y_j) \\ &\simeq \text{colim}_J \text{Map}_\beta(\text{colim}_I X_i, Y_j) \end{aligned}$$

Adjunct functor then: (1) $\mathcal{C} \in \text{Pr}^L$, D:locsm $\Rightarrow (\mathcal{C} \xrightarrow{f} \mathcal{D} \text{ adj} \Leftrightarrow \text{pres. small colim})$

(2) $\mathcal{C}, \mathcal{D} \in \text{Pr}^L \Rightarrow (f : \mathcal{C} \xrightarrow{f} \mathcal{D} \text{ radj} \Leftrightarrow \text{acc. & pres. small colim})$

Ex if \mathcal{C} presentable, $W \subseteq \text{Ar}(\mathcal{C})$ small
 $\rightarrow \{W\text{-loc. obj}\} \xrightarrow{\exists!} \mathcal{C}$

(In fact, Hacc. loc is of this form)

Def $\text{Pr}^R \subset \widehat{\text{Cat}}$ presentable cats & right-adjoints,

$\text{RFun}(\mathcal{C}^\text{op}, \mathcal{D})$

Prop $\text{Pr}^L \subset \widehat{\text{Cat}}$ creates small colimits
 $\text{Pr}^R \subset \widehat{\text{Cat}}$ $\left\{ \begin{array}{l} \text{creates small colimits} \\ \text{idea: the limit of } S \rightarrow \text{Pr}^{LR} \rightarrow \widehat{\text{Cat}} \\ = \text{Cartesian sections of some fib w/} \\ \text{presentable fibers.} \end{array} \right.$

Pr_K^L

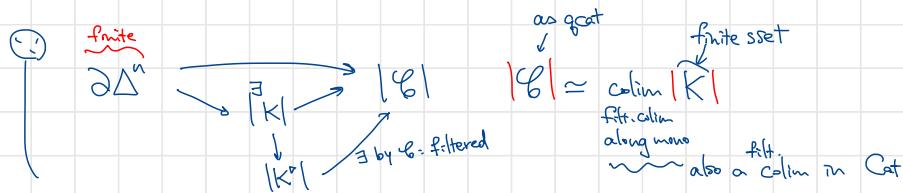
Exer? $\text{LFun}(\text{RFun}(\mathcal{C}^\text{op}, \mathcal{D}), \mathcal{E}) \simeq \text{LFun}(\mathcal{C}, \text{LFun}(\mathcal{D}, \mathcal{E}))$

$\begin{array}{c} \text{LFun}(\mathcal{C}, \mathcal{D}^\text{op})^\text{op} \\ \text{LFun}(\mathcal{D}^\text{op}, \mathcal{C}) \\ \text{RFun}(\mathcal{D}^\text{op}, \mathcal{C}) \end{array}$

ex. \mathcal{C} with final obj. is K -filt $\Leftrightarrow K$

- $X \in \text{Set}$ is K -opt $\Leftrightarrow \#X < K$.

\mathcal{C} : filtered $\Rightarrow |\mathcal{C}| = *$



Prop TFAE: (1) \mathcal{C} is K -filtered ($\forall f, \mathcal{C}_f + \phi$)

(2) $\forall K \xrightarrow{f} \mathcal{C}$ K : K -sm. $\Rightarrow \mathcal{C}_f$: K -filtered

(3) $\forall K$: K -sm, the diagonal $\mathcal{C} \rightarrow \text{Fun}(K, \mathcal{C})$ is cofinal. ($\Leftrightarrow \mathcal{C}_f$ | $\simeq *$)

(1) \Rightarrow (2) join of K -small sssets is K -small. (2) \Rightarrow (3) \Rightarrow (1) clear

Fact • $\forall \mathcal{C}$: ∞ -cat, \exists cofinal map from a poset $P \rightarrow \mathcal{C}$

If \mathcal{C} : filtered, $\exists P$: K -filtered ($K = \aleph_0 \rightarrow$ directed in classical sense)

• If $K > \aleph_0$, $\mathcal{C} \in q\text{Cat}$ or Kan is K -small (as sset)

$\Leftrightarrow \mathcal{C}$ is K -compact $\in \text{Cat}$ (or Ani). This is false for $K = \aleph_0$.

$$\text{Ani}^{\aleph_0} \not\supset \text{Ani}^{\aleph_0}$$

§ sifted colimits, projective objects, animation

Def A simplicial set K is sifted $\Leftrightarrow K \neq \emptyset, \Delta: K \rightarrow K \times K$ is cofinal
 $\Leftrightarrow \forall I: \text{fin set. } \Delta: K \rightarrow K^I$ is cofinal.

Rew ↗ invariant under cat. eq.
 \rightsquigarrow property of an ∞ -cat.

↪ Prop cofinal \Rightarrow w.h.e.

$$K \xrightarrow{\Delta} K \times K : \text{w.h.e.} \Rightarrow |K| \simeq *$$

$$\text{Quillen's thm A} \rightsquigarrow \mathcal{C} : \text{sifted} \Leftrightarrow \mathcal{C} \neq \emptyset, (\mathcal{C} \times \mathcal{C})_{(x,y)} / \mathcal{C} \times \mathcal{C} \stackrel{\mathcal{C} \text{ weakly contr}}{\simeq} \mathcal{C}_{(x,y)}/ \mathcal{C} \text{ has } \sqcup$$

Important examples: • filtered \Rightarrow sifted

• Δ^{op} is sifted.

proof NTS

$$\Delta^{\text{op}}_{[m]} / \Delta^{\text{op}}_{[n]} \simeq (\Delta_{/[m]} \times \Delta_{/[n]})^{\text{op}} : \text{weakly contr.}$$

IS \leftarrow use $\Delta \subset \text{Cat}$

$$\Delta_{/[m] \times [n]}$$

Note $\text{Cat} \xrightarrow{H} \text{Ani}$ is left-Kan extended from Δ

$$(L\text{Fun}(\text{Cat}, \text{Ani}) \hookrightarrow L\text{Fun}(Psh(\Delta), \text{Ani}) \simeq \text{Fun}(\Delta, \text{Ani}))$$

In particular, $\forall X \in \text{Cat}, |X| \simeq \text{colim } * \simeq |\Delta_X|$

$$\text{So } |\Delta_{/[m] \times [n]}| \simeq |[m] \times [n]| \simeq *.$$

\simeq has terminal obj

kerodon does more complicated argument. hope I'm not missing something...

I always get confused how to directly show

$$A \times B \simeq (A \times B)_{\text{circ}}$$

so I record it here:

$$\begin{array}{ccc} A+B & \xrightarrow{\quad} & C+B \\ \downarrow & \nearrow & \downarrow \\ A & \xrightarrow{\quad} & C \\ \downarrow & \nearrow & \downarrow \\ A-C & \xrightarrow{\quad} & C-C \\ \downarrow & \nearrow & \downarrow \\ A & \xrightarrow{\quad} & C \end{array}$$

back, front are pb
right is pb bc
 $B(C,A) \xrightarrow{\cong} B$
Similarity for bottom...
 \rightarrow so is top & left

Fact • finite product & sifted colim commute in Ani \rightsquigarrow Def. compact projective
 • sifted colimit " $=$ " filt. colim + geom real.
 • sifted colim + fin coprod " $=$ " all colim

• compact proj gen.
 $\exists \mathcal{C}_0, P_{\Sigma}(\mathcal{C}_0) \xrightarrow{\sim} \mathcal{C}$
 $\mathcal{C}_0^{\text{op}} \rightarrow \text{Ani}$ fin prod preserving

Analogous theory for 1-sifted colim: replace geom.

$$\mathcal{C}: 1\text{-cp gen} \Rightarrow \mathcal{C}^{1\text{-cp}} \Rightarrow P_{\Sigma}(\mathcal{C}^{1\text{-cp}}) \text{ animation of } \mathcal{C},$$

\uparrow localization
 $\text{Fun}(\Delta^{\text{op}}, \mathcal{C})$

Recap: What you need to know about presentable stuff

General thy: classes of diagram shapes $A, B \rightsquigarrow$

$(A, B) = \begin{cases} \emptyset \\ \text{idem. all} \end{cases}$	$\xrightarrow{\text{"small" / "filtered"}}$	$\begin{cases} A \& B - \text{colim "generate" small colim} \\ A^{\text{op}} - \text{lim \& B-colim commute in Ani} \end{cases}$	$\begin{cases} (\text{fin}, \text{filt}) \\ (\text{K-small}, \text{K-filt}) \end{cases}$	$\begin{cases} \text{(or your enriching cat)} \\ \text{nice thy of A-cocompletion} \\ \text{(of B-cocomplete cat} \rightsquigarrow \text{small cpt+}) \end{cases}$
--	---	--	--	--

locally small
Co-complete cats are
atomically gen \Rightarrow cpt proj gen \Rightarrow cpt gen ($\Rightarrow K\text{-cpt gen } \forall K: \text{inf reg}$)

(gen = closure of "cpt" obj under small colim)
(weak generation is enough!)

warning: $K \subset T$ does not imply
 $K\text{-cg} \Rightarrow T\text{-cg}$ but
 T is true for
certainly many $K \subset T$

of the form: $\mathcal{G} = P(\mathcal{C}_0)$

Can take $\mathcal{C}_0 = \mathcal{C}^{\text{atom}}$

\mathcal{C} : Cocomp translates to:

thus choice
is automatically
relm. cpt.

$P_{\Sigma}(\mathcal{C}_0)$

\mathcal{C}^{cp}

\mathcal{C}_0 has fin cocomp

$\text{Ind}(\mathcal{C}_0)$

\mathcal{C}^{Ko}

\mathcal{C}_0 : fin cocomp

$\text{Ind}_K(\mathcal{C}_0)$

\mathcal{C}^K

\mathcal{C}_0 : K-sm cpt.

Any obj of \mathcal{G} is
canonically a $\mathbb{Z}_{\geq 0}$
colim of \mathcal{C}_0

any small

sifted

filtered

K-filt.

← strong generation!

atomic pres
+ factor

P^L , idem

cp-pres. funct

P^L , cp-gen

fact: any obj of $P_{\Sigma}(\mathcal{C}_0)$ is canonically a
geom. real. of $\text{Ind}(\mathcal{C}_0)$

P^L - cpt obj pres
funct.

Ind^L -

\mathcal{C}_0 -

rex funct.

K-cpt obj pres
funct.

P^L_K -

\mathcal{C}_0 -

rex funct.

colim-pres

P^L_R = $\bigcup_K P^L_K$

$\subset \widehat{\text{Cat}}$

$\begin{array}{l} \text{ex} \\ \mathcal{G} \xleftarrow{L} R \xrightarrow{R} \mathcal{D} \\ L \text{ pres K-cpt obj} \\ \Leftrightarrow R \text{ pres K-filt colim} \end{array}$

non-full subsets:

$P \uparrow \downarrow \approx$
 $\text{Cat} \supset \text{Cat}^{\text{idem}}$

$P_{\Sigma} \uparrow \downarrow \approx$
 $\text{Cat} \xrightarrow{\text{idem}} \text{Cat}^{\text{idem}}$
↓ pres funct.

$\text{Ind} \uparrow \downarrow \approx$
 $\text{Cat}^{\text{rex}} \supset \text{Cat}^{\text{rex}}$
↓ rex funct.

$\text{Ind}_K \uparrow \downarrow \approx$
 $\text{Cat}^{\text{rex}} \supset \text{Cat}^{\text{rex}}$
↓ rex funct.

ex $\mathcal{C}_0 = \langle R^{\otimes n} |_{n \geq 0} \rangle \rightsquigarrow \mathcal{C}^{\text{idem}} \hookrightarrow \text{Ind}(\mathcal{C}_0) \hookrightarrow \text{Fun}(I_+, \text{Set})$
 $\hookrightarrow \text{Mod}_R^{\text{op}}$ proj flat

Prop \mathcal{G} : presentable. $\mathcal{C}_0 \hookrightarrow \mathcal{G}$
• dense $\Leftrightarrow \mathcal{G} \rightarrow P_{\Sigma}(\mathcal{C}_0)$ ff.
weakly dense \Leftrightarrow conservative.
(colim closure of $\mathcal{C}_0 \rightarrow \mathcal{G}$)

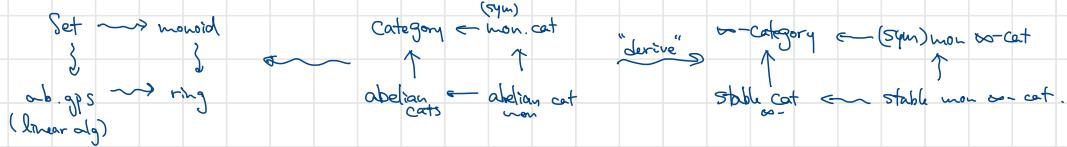
$P_{(K)}^R$: presentable cats
& right adjoints (pres. K-filt)
colim

$P_K^R \simeq P_K^L, \text{op}$

Fact $P^L, P_{(K)}^R \subset \widehat{\text{Cat}}$ creates limits

Part 2. Higher algebra

(rings & modules)



{ Stable categories }

Def • A category \mathcal{C} is pointed if $\exists 0 \in \mathcal{C}$: zero obj i.e. the object which is both initial & terminal.

$(\Leftrightarrow \exists_{\text{initial}} \xrightarrow{\phi} \exists_{\text{terminal}} \xleftarrow{*} \text{obj.}, \text{ and the canonical map } \phi \rightarrow * \text{ is an iso.})$

- A category \mathcal{C} is stable if (1) $0 \in \mathcal{C}$.

(2) finite limits & colimits exist. ($\stackrel{\exists}{\Leftarrow}$ pushouts & pullbacks)

(3) A square $\Delta^1 \times \Delta^1 \rightarrow \mathcal{C}$ is cartesian iff cocartesian

Slogan Stability = fin lim & fin colim "agree".

Def • $F: \mathcal{C} \rightarrow \mathcal{D}$ between cats w/ fin.lim is left exact if F pres. fin.lim.

• colim right exact colim

• exact if both.

exer \mathcal{C} all equivalent if \mathcal{C}, \mathcal{D} : stable. $\text{Cat}^{\text{st}} \subset \text{Cat}$ non-full of stable cats & exact functors

Rank • $\text{Cat}^{\text{st}} \subset \text{Cat}$ creates $\begin{cases} \text{Fun}(K, -) & \text{& small limits, (i.e., Cat-weighted)} \\ \text{Ind}_K(-), (-)^K & \text{of stable categories are stable.} \end{cases}$ \lim

$$\rightsquigarrow \begin{array}{c} \text{Cat}^{\text{st}} \xrightarrow{\text{Ind}} \text{Pr}_{w.st}^L \\ \Downarrow \begin{matrix} \text{Ind} \\ (-)^K \end{matrix} \end{array} \quad \begin{array}{c} \text{Cat}^{\text{st}, K\text{-rex}} \xrightarrow{\text{Ind}_K} \text{Pr}_{K,st}^L \\ \Downarrow \begin{matrix} \text{Ind}_K \\ (-)^K \end{matrix} \end{array}$$

Ren

I think $\text{Cat}^{\text{st}} \xrightarrow{\perp} \text{Cat}$ is given by the stable closure of Sp-Yoneda emb $\mathcal{C} \hookrightarrow \text{fun}(\mathbb{P}, \mathbb{P})$.

Prop • \mathcal{C} with pushouts.

$$\hookrightarrow \text{Fun}(\Delta_0^2, \mathcal{C}) \xrightleftharpoons[\sim]{\text{res}} \text{Fun}([1]^2, \mathcal{C}) \xrightleftharpoons[\sim]{\text{res}} \text{Fun}(\Delta_2^2, \mathcal{C})$$

$\{\text{po sq}\}$

$\{\text{pb sq}\}$

• $\text{res} \circ \text{LKE}$ has a radj iff \mathcal{C} has pullbacks & the radj is given by pullback
 $\text{po} = \text{pb}$ iff this is an adjoint equiv iff $\text{res} \circ \text{LKE}$ is an equivalence.

Next goal: weaken (2) & (3)

Def $\forall x, y \in \mathcal{C}. \quad * \simeq \text{Map}(x, 0) \times \text{Map}(0, y) \xrightarrow{\circ} \text{Map}(x, y)$ picks up a morphism $x \rightarrow y$. denoted by 0 .

$\mathcal{C} = \text{An}_x \rightsquigarrow x \circ y$ is a constant map at the base point.

\triangleleft zero morphism can have a nontrivial automorphism; $\forall X: \text{conh. } (S^0 \rightarrow X) \simeq (S^0 \circ X)$.
 e.g. (SLX many identifications)

Def - A triangle in \mathcal{C} is a commutative square

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow & & \downarrow g \\ 0 & \rightarrow & C \end{array}$$

often denoted by $A \xrightarrow{f} B \xrightarrow{g} C$ by abuse of notation, making the homotopy $gof \simeq 0$ implicit.
 (but it is a part of data)

• $A \xrightarrow{f} B \xrightarrow{g} C$ is a (co)fiber seq if the square is (co)cartesian

\hookrightarrow $\begin{cases} A \text{ (or } f\text{)} \text{ is a fiber of } g & (A \xrightarrow{f} f^{-1}B \xrightarrow{g} C) \\ C \text{ cotrivial } f & (cof(f) \xrightarrow{g} C) \end{cases}$

• $\begin{array}{c} A \xrightarrow{f} B \\ \downarrow & \downarrow g \\ 0 \xrightarrow{0} B \end{array} \rightsquigarrow \begin{cases} A \simeq \Omega B & \text{if it is cartesian (loop)} \\ \Sigma A \simeq B & \text{if it is cocartesian (suspension)} \end{cases}$

exer \mathcal{C} : pointed, has fibers & cotibers / loop & suspension

$\Rightarrow \begin{cases} \text{• define } \text{Fun}(\Delta^1, \mathcal{C}) \xrightleftharpoons[\sim]{\text{fib}} \text{Fun}(\Delta^1, \mathcal{C}). \text{ It has a radj iff } \mathcal{C} \text{ has fibers.} \\ \quad X \xrightarrow{f} Y \quad \longmapsto \quad Y \rightarrow \text{cof}(f) \quad (\text{then radj is given by } (Y \xrightarrow{f} Y) \hookrightarrow (f^{-1}B \xrightarrow{g} C)) \\ \text{• Similarly for } \mathcal{C} \xrightleftharpoons[\sim]{\text{fib}} \mathcal{C}. \end{cases}$

When \mathcal{C} : stable, these adjunctions are equivalences.

shift notation

$$\sum^n X := X[n]$$

$$\Sigma^n X := X[-n]$$

Prop For \mathcal{C} pointed,

$$\begin{aligned} \mathcal{C} = \text{Stable} &\iff \begin{cases} (2') \forall f: X \rightarrow Y \text{ has a fiber } (\cong \text{a cofiber}) \\ (3') \text{ A triangle } X \xrightarrow{f} Y \xrightarrow{g} Z \text{ is a fib seq iff it is a cofib seq.} \end{cases} & \text{Fun}(\Delta^1, \mathcal{C}) \xrightarrow{\text{fib}} \text{Fun}(\Delta^1, \mathcal{C}) \\ &\iff \begin{cases} (2) \mathcal{C} \text{ has finite limits (e.g. fin. colim.)} \\ (3') (\Sigma \dashv \Omega) \text{ is an equivalence.} \end{cases} \end{aligned}$$

proof sketch: (3) \Rightarrow (3') \Rightarrow (3'') is dear.

$$\begin{aligned} \bullet (2') + (3'') &\Rightarrow \text{additive} & \left[\begin{array}{c} \Omega X \rightarrow Y \rightarrow \circ \\ \downarrow \quad \downarrow \quad \downarrow \\ \circ \rightarrow X \sqcup Y \xrightarrow{\text{red}} X \\ \downarrow \quad \downarrow \quad \downarrow \\ \circ \rightarrow \Sigma Y \end{array} \right] & \text{semiadditive} \\ \rightsquigarrow \text{eq}(X \xrightarrow{f} Y) &\cong \text{fib}(X \xrightarrow{f+g} Y). & \Rightarrow (2). \\ \bullet (3') \Rightarrow (3) &\text{ exer.} & (2)(3'') \Rightarrow (3) &\text{ harder exer.} \end{aligned}$$

$$+ \forall X \simeq \underset{\overset{\circ}{\longrightarrow}}{\Sigma} \underset{\overset{\circ}{\longleftarrow}}{\Sigma} X$$

$$\text{Map}(Y, \Sigma \Sigma X) \stackrel{\cong}{=} \underset{\overset{\circ}{\longrightarrow}}{\Sigma} \underset{\overset{\circ}{\longleftarrow}}{\Sigma} \underset{\overset{\circ}{\longrightarrow}}{\Sigma} Y = \underset{\overset{\circ}{\longrightarrow}}{\Sigma} \underset{\overset{\circ}{\longleftarrow}}{\Sigma} \underset{\overset{\circ}{\longrightarrow}}{\Sigma} X$$

Def $S^0 = (*)_+ \in \text{Ani}_*$, $S^n = \Sigma^n S^0 \simeq \Sigma^n X = \text{Map}_*(S^n, X)$, $\pi_{\text{th}} X := \pi_0 \Sigma^n X$

Rem fib seq \Rightarrow long exact seq of homotopy groups. (reduces to: $\pi_0: \text{Ani}_+ \rightarrow \text{Set}_+$ preserves fibers)

$$\in \text{Gp}, \text{Ab}$$

from the cogroup str on S^1
 $\in \text{ho}(\text{Ani}_+)$