

## salt1: Immutable References

CS392-M1: *Rust, In Practice and in Theory*

# Syntax

$x$	(variables, $\mathcal{V}$ )
$n$	(integers, $\mathbb{Z}$ )
$w ::= x \mid * w$	(place expression, $\mathcal{W}$ )
$e ::= () \mid n \mid w \mid \& w \mid x = e$	(expressions, $\mathcal{E}$ )
$s ::= \text{let } x = e \mid \text{let mut } x = e \mid e$	(statements, $\mathcal{S}$ )
$p ::= e \mid s ; p$	(programs, $\mathcal{P}$ )

# Typing

$t ::= \textcolor{red}{()} \mid \textcolor{red}{i32} \mid \& x$	(types, $\mathcal{T}$ )
$m ::= \text{imm} \mid \text{mut}$	(mutability)
$u ::= \langle t \rangle^m$	(slot types, $\mathbb{S}_{\mathcal{T}}$ )
$\Gamma \in \mathcal{V} \mapsto \mathbb{S}_{\mathcal{T}}$	(contexts)
$\Gamma \vdash \text{writable}(x)$	(writability)
$\Gamma \vdash w : u$	(place expressions)
$\Gamma \vdash t \approx t$	(type compatibility)
$\Gamma \vdash e : t \dashv \Gamma$	(expressions)
$\Gamma \vdash s \dashv \Gamma$	(statements)
$\Gamma \vdash p : t \dashv \Gamma$	(programs)

$$\begin{array}{c}
\frac{(x \mapsto \langle t \rangle^m) \in \Gamma}{\Gamma \vdash x : \langle t \rangle^m} \text{ VAR} \quad \frac{\Gamma \vdash w : \langle \& x \rangle^{m_1} \quad \Gamma \vdash x : \langle t \rangle^{m_2}}{\Gamma \vdash * w : \langle t \rangle^{m_2}} \text{ DEREF} \quad \frac{}{\Gamma \vdash () : ()} \text{ UNIT} \\[10pt]
\frac{n \in \mathbb{Z}}{\Gamma \vdash n : \text{i32} \dashv \Gamma} \text{ INT} \quad \frac{\Gamma \vdash w : \langle t \rangle^m}{\Gamma \vdash w : t \dashv \Gamma} \text{ PLACE} \quad \frac{\Gamma \vdash x : \langle t \rangle^m}{\Gamma \vdash \& x : \& x \dashv \Gamma} \text{ REFVAR} \\[10pt]
\frac{\Gamma \vdash w : \langle \& x \rangle^m}{\Gamma \vdash \& * w : \& x \dashv \Gamma} \text{ REFDEREF} \quad \frac{}{\Gamma \vdash () \approx ()} \text{ } \approx\text{-UNIT} \quad \frac{}{\Gamma \vdash \text{i32} \approx \text{i32}} \text{ } \approx\text{-INT} \\[10pt]
\frac{\Gamma \vdash x_1 : \langle t_1 \rangle^{m_1} \quad \Gamma \vdash x_2 : \langle t_2 \rangle^{m_2} \quad \Gamma \vdash t_1 \approx t_2}{\Gamma \vdash \& x_1 \approx \& x_2} \text{ } \approx\text{-REF} \quad \frac{\nexists y. (y \mapsto \langle \& x \rangle^m) \in \Gamma}{\Gamma \vdash \text{writable}(x)} \text{ WRITABLE} \\[10pt]
\frac{\Gamma_1 \vdash x : \langle t_1 \rangle^{\text{mut}} \quad \Gamma_1 \vdash e : t_2 \dashv \Gamma_2 \quad \Gamma_2 \vdash t_1 \approx t_2 \quad \Gamma_2 \vdash \text{writable}(x)}{\Gamma_1 \vdash x = e : () \dashv \Gamma_2[x \mapsto \langle t_2 \rangle^{\text{mut}}]} \text{ ASSIGN} \\[10pt]
\frac{\Gamma_1 \vdash e : t \dashv \Gamma_2 \quad x \notin \text{dom}(\Gamma_2)}{\Gamma_1 \vdash \text{let } x = e \dashv \Gamma_2[x \mapsto \langle t \rangle^{\text{imm}}]} \text{ LET} \quad \frac{\Gamma_1 \vdash e : t \dashv \Gamma_2 \quad x \notin \text{dom}(\Gamma_2)}{\Gamma_1 \vdash \text{let mut } x = e \dashv \Gamma_2[x \mapsto \langle t \rangle^{\text{mut}}]} \text{ LETMUT} \\[10pt]
\frac{\Gamma_1 \vdash e : t \dashv \Gamma_2}{\Gamma_1 \vdash e \dashv \Gamma_2} \text{ EXPRSTMT} \quad \frac{\Gamma_1 \vdash s \dashv \Gamma_2 \quad \Gamma_2 \vdash p : t \dashv \Gamma_3}{\Gamma_1 \vdash s ; p : t \dashv \Gamma_3} \text{ PROG}
\end{array}$$

# Evaluation

$$\begin{array}{ll}
\ell ::= \ell_x & (\text{locations}, \mathcal{L}) \\
v ::= \textcolor{red}{()} \mid n \mid \ell & (\text{values}, \mathbb{V}) \\
\\
S \in \mathcal{L} \mapsto \mathbb{V} & (\text{store}) \\
\\
S \vdash w \rightsquigarrow \ell & (\text{place locations}) \\
S \vdash e \Downarrow v \dashv S & (\text{expressions}) \\
S \vdash s \dashv S & (\text{statements}) \\
S \vdash p \Downarrow v \dashv S & (\text{programs}) \\
\\
\frac{}{S \vdash x \rightsquigarrow \ell_x} \text{LOCVAR} & \frac{S \vdash w \rightsquigarrow \ell_x \quad (\ell_x \mapsto \ell_y) \in S}{S \vdash *w \rightsquigarrow \ell_y} \text{LOCDEREF} \quad \frac{}{S \vdash \textcolor{red}{()} \Downarrow \textcolor{red}{()} \dashv S} \text{UNIT} \\
\frac{n \in \mathbb{Z}}{S \vdash n \Downarrow n \dashv S} \text{INT} & \frac{S \vdash w \rightsquigarrow \ell_x \quad (\ell_x \mapsto v) \in S}{S \vdash w \Downarrow v \dashv S} \text{PLACE} \quad \frac{S \vdash w \rightsquigarrow \ell_x}{S \vdash \& w \Downarrow \ell_x \dashv S} \text{REF} \\
\frac{\ell_x \in \text{dom}(S) \quad S_1 \vdash e \Downarrow v \dashv S_2}{S_1 \vdash x = e \Downarrow \textcolor{red}{()} \dashv S_2[\ell_x \mapsto v]} \text{ASSIGN} & \frac{S_1 \vdash e \Downarrow v \dashv S_2}{S_1 \vdash \text{let } x = e \dashv S_2[\ell_x \mapsto v]} \text{LET} \\
\frac{S_1 \vdash e \Downarrow v \dashv S_2}{S_1 \vdash \text{let mut } x = e \dashv S_2[\ell_x \mapsto v]} \text{LETMUT} & \frac{S_1 \vdash e \Downarrow v \dashv S_2}{S_1 \vdash e \dashv S_2} \text{EXPRSTMT} \\
\frac{S_1 \vdash s \dashv S_2 \quad S_2 \vdash p \Downarrow v \dashv S_3}{S_1 \vdash s ; p \Downarrow v \dashv S_3} \text{PROC}
\end{array}$$