

salt3: Blocks

CS392-M1: *Rust, In Practice and in Theory*

Syntax

x	(variables, \mathcal{V})
n	(integers, \mathbb{Z})
$w ::= x \mid * w$	(place expression, \mathcal{W})
$e ::= () \mid n \mid w \mid \& w \mid \&\text{mut} w \mid \text{copy } w \mid w = e \mid \{ p \}$	(expressions, \mathcal{E})
$s ::= e \mid \text{let } x = e \mid \text{let mut } x = e$	(statements, \mathcal{S})
$p ::= e \mid s ; p$	(programs, \mathcal{P})

Typing

l	(lifetimes, Lt)
$t ::= (\textcolor{red}{()}) \mid \textcolor{red}{i32} \mid \& w \mid \&\text{mut } w$	(types, \mathcal{T})
$\tilde{t} ::= \lfloor t \rfloor \mid t$	(partial types, $\tilde{\mathcal{T}}$)
$m ::= \text{imm} \mid \text{mut}$	(mutability)
$u ::= \langle \tilde{t} \rangle_l^m$	(slot types, $\mathbb{S}_{\mathcal{T}}$)
$\Gamma \in \mathcal{V} \mapsto \mathbb{S}_{\mathcal{T}}$	(contexts)
$\text{copyable}(t)$	(copyability)
$\Gamma \vdash w : u$	(place expressions)
$\Gamma \vdash \text{readable}(w)$	(readability)
$\Gamma \vdash \text{writable}(w)$	(writability)
$\Gamma \vdash \tilde{t} \approx \tilde{t}$	(type compatibility)
$\Gamma \vdash \langle e : t \rangle_l \dashv \Gamma$	(expressions)
$\Gamma \vdash \langle s \rangle_l \dashv \Gamma$	(statements)
$\Gamma \vdash \langle p : t \rangle_l \dashv \Gamma$	(programs)
$\frac{(x \mapsto \langle \tilde{t} \rangle_l^m) \in \Gamma}{\Gamma \vdash x : \langle \tilde{t} \rangle_l^m} \text{ (var)}$	
$\frac{\Gamma \vdash w_1 : \langle \& w_2 \rangle_{l_1}^{m_1} \quad \Gamma \vdash w_2 : \langle t \rangle_{l_2}^{m_2}}{\Gamma \vdash * w_1 : \langle t \rangle_{l_2}^{m_2}} \text{ (deref)}$	
$\frac{}{\Gamma \vdash \langle \textcolor{red}{()} : \textcolor{red}{()} \rangle_l \dashv \Gamma} \text{ (unit)}$	
$\frac{n \in \mathbb{Z}}{\Gamma \vdash \langle n : \textcolor{red}{i32} \rangle_l \dashv \Gamma} \text{ (int)}$	
$\frac{\#y, l. (y \mapsto \&\text{mut } *^l x) \in \Gamma}{\Gamma \vdash \text{readable}(*^k x)} \text{ (readable)}$	

$$\begin{array}{c}
\frac{}{\text{copyable}(\textcolor{red}{()})} \text{ (copy-unit)} \\
\frac{}{\text{copyable}(\textcolor{red}{i32})} \text{ (copy-int)} \\
\frac{}{\text{copyable}(\& w)} \text{ (copy-brw)} \\
\frac{\Gamma \vdash w : \langle t \rangle_{l_1}^m \quad \Gamma \vdash \text{readable}(w) \quad \text{copyable}(t)}{\Gamma \vdash \langle \text{copy } w : t \rangle_{l_2} \dashv \Gamma} \text{ (place-copy)} \\
\frac{\Gamma \vdash \text{readable}(*^k x) \quad \nexists y, l. (y \mapsto \& *^l x) \in \Gamma}{\Gamma \vdash \text{writable}(*^k x)} \text{ (writable)} \\
\frac{\Gamma \vdash x : \langle \&\text{mut } w \rangle_{l_1}^m \quad \Gamma \vdash \text{writable}(x)}{\Gamma \vdash \langle x : \&\text{mut } w \rangle_{l_2} \dashv \Gamma[x \mapsto \lfloor \&\text{mut } w \rfloor]} \text{ (place-move)} \\
\frac{\Gamma \vdash w : \langle t \rangle_{l_1}^m \quad \Gamma \vdash \text{readable}(w)}{\Gamma \vdash \langle \& w : \& w \rangle_{l_2} \dashv \Gamma} \text{ (brw)} \\
\frac{\Gamma \vdash w : \langle t \rangle_{l_1}^{\text{mut}} \quad \Gamma \vdash \text{writable}(w) \quad \Gamma \vdash \text{mutable}(w)}{\Gamma \vdash \langle \&\text{mut } w : \&\text{mut } w \rangle_{l_2} \dashv \Gamma} \text{ (mut-brw)} \\
\frac{}{\Gamma \vdash \textcolor{red}{i32} \approx \textcolor{red}{i32}} \text{ (\approx-int)} \\
\frac{}{\Gamma \vdash \textcolor{red}{()} \approx \textcolor{red}{()} } \text{ (\approx-unit)} \\
\frac{\Gamma \vdash w_1 : \langle \tilde{t}_1 \rangle_{l_1}^{m_1} \quad \Gamma \vdash w_2 : \langle \tilde{t}_2 \rangle_{l_2}^{m_2} \quad \Gamma \vdash \tilde{t}_1 \approx \tilde{t}_2}{\Gamma \vdash \& w_1 \approx \& w_2} \text{ (\approx-brw)} \\
\frac{\Gamma \vdash w_1 : \langle \tilde{t}_1 \rangle_{l_1}^{m_1} \quad \Gamma \vdash w_2 : \langle \tilde{t}_2 \rangle_{l_2}^{m_2} \quad \Gamma \vdash \tilde{t}_1 \approx \tilde{t}_2}{\Gamma \vdash \&\text{mut } w_1 \approx \&\text{mut } w_2} \text{ (\approx-mbrw)} \\
\frac{\Gamma \vdash t_1 \approx \tilde{t}_2}{\Gamma \vdash \lfloor t_1 \rfloor \approx \tilde{t}_2} \text{ (\approx-partial}_1\text{)} \\
\frac{\Gamma \vdash \tilde{t}_1 \approx t_2}{\Gamma \vdash \tilde{t}_1 \approx \lfloor t_2 \rfloor} \text{ (\approx-partial}_2\text{)}
\end{array}$$

$$\begin{aligned}
\text{write}(\Gamma, x, t) &= \Gamma[x \mapsto t] \\
\text{write}(\Gamma, *^{k+1} x, t) &= \text{write}(\Gamma, *^k w, t) \quad \text{where} \quad (x \mapsto \&\text{mut } w^m) \in \Gamma
\end{aligned}$$

$$\begin{aligned}
\text{replace}(\&*^{k+1}w_1, w_1, \&[\text{mut}]w_2) &= \&*^k w_2 \\
\text{replace}(\&\text{mut } *^{k+1}w_1, w_1, \&[\text{mut}]w_2) &= \&\text{mut } *^k w_2 \\
\text{replace}(t_1, w, t_2) &= t_1 \\
\text{replace}(\lfloor t_1 \rfloor, w, t_2) &= \lfloor \text{replace}(t_1, w, t_2) \rfloor \\
\text{replace}(\Gamma, w, t_2) &= \{x \mapsto \text{replace}(\tilde{t}_1, w, t_2) : (x \mapsto \tilde{t}_1) \in \Gamma\}
\end{aligned}$$

$$\text{update}(\Gamma, w, t_2) = \text{replace}(\text{write}(\Gamma, w, t_2), w, t_2)$$

$$\frac{}{\Gamma \vdash \text{i32} : l} \text{ (int-lftm)}$$

$$\frac{\Gamma \vdash w : \langle t \rangle_{l_1}^m \quad l_1 \leq l_2}{\Gamma \vdash \& w : l_2} \text{ (brw-lftm)}$$

$$\frac{\Gamma \vdash w : \langle t \rangle_{l_1}^m \quad l_1 \leq l_2}{\Gamma \vdash \&\text{mut } w : l_2} \text{ (mut-brw-lftm)}$$

$$\frac{\Gamma_1 \vdash w : \langle \tilde{t}_1 \rangle_{l_1}^{\text{mut}} \quad \Gamma_1 \vdash e : t_2 \dashv \Gamma_2 \quad \Gamma_2 \vdash \tilde{t}_1 \approx t_2}{\Gamma_2 \vdash t_2 : l_1 \quad \Gamma_3 = \text{update}(\Gamma, w, t_2) \quad \Gamma_3 \vdash \text{writable}(w)} \frac{}{\Gamma_1 \vdash \langle w = e : \textcolor{red}{()} \rangle_{l_2} \dashv \Gamma_3} \text{ (assign)}$$

$$\text{drop}(\Gamma, l) = \Gamma \setminus \{x \mapsto \langle \tilde{t} \rangle_l^m : (x \mapsto \langle \tilde{t} \rangle_l^m) \in \Gamma\}$$

$$\frac{\Gamma_1 \vdash \langle p : t \rangle_{l+1} \dashv \Gamma_2 \quad \Gamma_2 \vdash t : l}{\Gamma_1 \vdash \langle \{ p \} : t \rangle_l \dashv \text{drop}(\Gamma_2, l+1)} \text{ (block)}$$

$$\frac{\Gamma_1 \vdash \langle e : t \rangle_l \dashv \Gamma_2}{\Gamma_1 \vdash \langle e \rangle_l \dashv \Gamma_2} \text{ (expr-stmt)}$$

$$\frac{\Gamma_1 \vdash \langle e : t \rangle_l \dashv \Gamma_2 \quad x \notin \text{dom}(\Gamma_2)}{\Gamma_1 \vdash \langle \text{let } x = e \rangle_l \dashv \Gamma_2[x \mapsto \langle t \rangle_l^{\text{imm}}]} \text{ (let)}$$

$$\frac{\Gamma_1 \vdash \langle e : t \rangle_l \dashv \Gamma_2 \quad x \notin \text{dom}(\Gamma_2)}{\Gamma_1 \vdash \langle \text{let mut } x = e \rangle_l \dashv \Gamma_2[x \mapsto \langle t \rangle_l^{\text{mut}}]} \text{ (let-mut)}$$

$$\frac{\Gamma_1 \vdash \langle s \rangle_l \dashv \Gamma_2 \quad \Gamma_2 \vdash \langle p : t \rangle_l \dashv \Gamma_3}{\Gamma_1 \vdash \langle s ; p : t \rangle_l \dashv \Gamma_3} \text{ (prog)}$$

Evaluation

$$\begin{array}{ll}
\ell ::= \ell_x & (\text{locations}, \mathcal{L}) \\
v ::= \textcolor{red}{()} \mid n \mid \ell & (\text{values}, \mathbb{V}) \\
\tilde{v} ::= \perp \mid v & (\text{partial values}, \tilde{\mathbb{V}}) \\
r ::= \langle \tilde{v} \rangle_l & (\text{slot types}, \mathbb{S}_{\mathbb{V}}) \\
\\
S \in \mathcal{L} \mapsto \mathbb{S}_{\mathbb{V}} & (\text{store}) \\
\\
\langle S, e \Downarrow S, v \rangle_l & (\text{expressions}) \\
\langle S, s \Downarrow S \rangle_l & (\text{statements}) \\
\langle S, p \Downarrow S, v \rangle_l & (\text{programs}) \\
\\
\frac{}{x \rightsquigarrow \ell_x} (\text{loc-var}) \\
\\
\frac{w \rightsquigarrow \ell_x \quad (\ell_x \mapsto \ell_y) \in S}{* w \rightsquigarrow \ell_y} (\text{loc-drf}) \\
\\
\frac{}{\langle S, \textcolor{red}{()} \Downarrow S, \textcolor{red}{()} \rangle_l} (\text{unit}) \\
\\
\frac{n \in \mathbb{Z}}{\langle S, n \Downarrow S, n \rangle_l} (\text{int}) \\
\\
\frac{w \rightsquigarrow \ell_x \quad (\ell_x \mapsto \langle v \rangle_{l_1}) \in S}{\langle S, \text{copy } w \Downarrow S, v \rangle_{l_2}} (\text{place-copy}) \\
\\
\frac{w \rightsquigarrow \ell_x \quad (\ell_x \mapsto \langle v \rangle_{l_1}) \in S}{\langle S, w \Downarrow S[\ell_x \mapsto \langle \perp \rangle_{l_1}], v \rangle_{l_2}} (\text{move}) \\
\\
\frac{w \rightsquigarrow \ell}{\langle S, \& w \Downarrow S, \ell \rangle_l} (\text{brw}) \\
\\
\frac{w \rightsquigarrow \ell}{\langle S, \&\text{mut } w \Downarrow S, \ell \rangle_l} (\text{mut-brw})
\end{array}$$

$$\frac{w \rightsquigarrow \ell_x \quad (\ell_x \mapsto \langle \tilde{v} \rangle_{l_1}) \in S_1 \quad \langle S_1, e \Downarrow S_2, v \rangle_{l_2}}{\langle S_1, w = e \Downarrow S_2[\ell_x \mapsto \langle v \rangle_{l_1}], \textcolor{red}{(\textcolor{red}{O})} \rangle_{l_2}} \text{(assign)}$$

$$\text{drop}(S, l) = S \setminus \{\ell \mapsto \langle \tilde{v} \rangle_l : (\ell \mapsto \langle \tilde{v} \rangle_l) \in S\}$$

$$\frac{\langle S_1, p \Downarrow S_2, v \rangle_{l+1}}{\langle S_1, \{p\} \Downarrow \text{drop}(S_2, l+1), v \rangle_l} \text{(block)}$$

$$\frac{\langle S_1, e \Downarrow S_2, v \rangle_l}{\langle S_1, e \Downarrow S_2 \rangle_l} \text{(expr-stmt)}$$

$$\frac{\langle S_1, e \Downarrow S_2, v \rangle_l}{\langle S_1, \text{let } x = e \Downarrow S_2[\ell_x \mapsto \langle v \rangle_l] \rangle_l} \text{(let)}$$

$$\frac{\langle S_1, e \Downarrow S_2, v \rangle_l}{\langle S_1, \text{let mut } x = e \Downarrow S_2[\ell_x \mapsto \langle v \rangle_l] \rangle_l} \text{(let-mut)}$$

$$\frac{\langle S_1, e \Downarrow S_2, v \rangle_l}{\langle S_1, e \Downarrow S_2 \rangle_l} \text{(expr-stmt)}$$

$$\frac{\langle S_1, s \Downarrow S_2 \rangle_l \quad \langle S_2, p \Downarrow S_3, v \rangle_l}{\langle S_1, s ; p \Downarrow S_3, v \rangle_l} \text{(prog)}$$