

# Hash Tables

- Components
  - keys/elements
  - array
- hashCode
  - hashCode ideally creates a unique integer for each key based on the key's fields
    - hashCode is based on the fields of the object
    - two objects that are equal--have the same fields--produce the same hashCode
    - if overriding equals or hashCode, need to override the other as well
- hash function maps keys to array indices
  - hash function uses key's hashCode and the array length in computation, typically `key.hashCode() % array.length`
  - another hash function squares the hashCode and returns middle digits
  - another hash function multiplies the key's hashCode by a number less than 1 and returns the first few digits of the decimal
- Open-Address Hashing
  - dealing with collisions
    - collision: when there is already an item where a new element belongs
    - linear probing: search next indices one at a time until you find an empty space
      - causes clustering: when many items end up next to each other
      - avoid clustering with double hashing:
        - a second hash function determines how the next index is found
        - `data.length` and `data.length-2` must be prime numbers
          - If you want to know \*why\*, see Donald Knuth's The Art of Programming, Vol 3
          - Donald Knuth is like a CS theory god
        - second hash function returns a number between 1 and `data.length-2`
      - quadratic probing:  $h+1, h+2^2, h+3^2, h+4^2, \dots$
    - Worst-case runtime:  $O(n)$ 
      - Have  $n$  collisions before finding the element or an available location
    - Typically adding/searching/removing takes  $O(1)$  with open address and a good hash function
    - What happens if we remove an element from the table?
      - What if we search for another element that had previously collided with that element?
      - One way to deal with this is a boolean array to record whether there used to be something there

- Another way is to use “placeholders” or “zombies”--put a dummy element or a sentinel value in the spot to replace a removed item
      - When searching, pretend the zombie is an element
      - When adding, replace the zombie
- Chained Hashing
  - each array index has a linked list
  - if there is a collision, just add to the list
    - speeds up adding  $O(1)$
    - searching and removing can still be slow,  $O(n)$ 
      - If our hashCode and hash functions were really bad and everything ended up at the same index, we'd be searching through a linked list with  $n$  elements
- Efficiency
  - depends on the load factor, or a percentage of the array's usage
  - To keep things efficient, we have a max capacity; in your homework this is  $\frac{3}{4}$ 
    - as soon as your hashtable is 75% full, you make a bigger array and rehash--DO NOT COPY--the elements
      - the results of the hash function depend on the array length, so the hash values will no longer be the same!
  - depends on the type of hashing (open-address vs chained) and on the collision handling technique with open-address
  -