

Linked ListsDoubly linked

- next & previous ptrs
- keep in mind extra links to change
- prev & next are opposites:
 $x.\text{prev}.\text{next} = x.\text{next}.\text{prev} = x$

* must make sure correctly linked in well-formed

Adding

- at beginning: make sure to point old head's prev to new
- at end: make sure to point new tail's prev to old
- in middle: make sure updating in correct order!

`Node p = new Node(v);`

* add after current

`p.next = current.next;`

`current.next = p;`

`p.prev = current;`

* `current.next.prev = p,`
current.next is p.

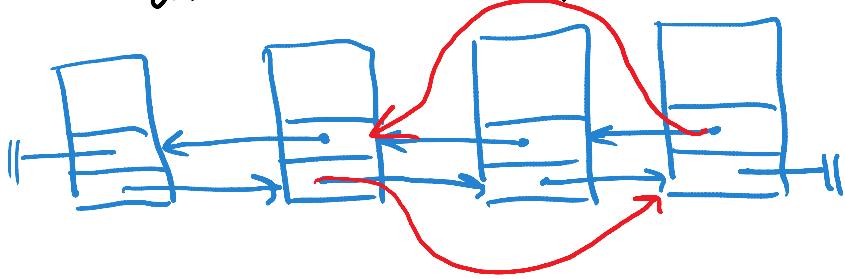
* Current. next · prev = null

current.next is p.

→ create a cycle

Removal

- at beginning: make new head's prev null
- at end: make new tail's next null
- in middle: keep track of prevs



$$n.\text{next}.\text{prev} = n.\text{prev}.$$

$$n.\text{prev}.\text{next} = n.\text{next}$$

Dummy Nodes

- extra node typically at front of list w/ null data
- precursor never null
- simpler - fewer special cases

Cyclic Lists

- head & tail is connected
- no null pointers
- precursor never null

- no null pointers
 - precursor never null
 - fewer special cases, simpler code
- adding first or removing last are special cases if there is no dummy node
 - no special cases w/ a dummy

Endogenous

- links are part of the data
- no nodes
- in regular lists, a node acts as a box that data is put into
- endogenous lists have no nodes & remain unboxed
- objects point to each other.
- does not require additional objects - space efficient
- simplifies code
- an object cannot be in 2

- an object cannot be in 2 lists simultaneously
- when removing, must set next & prev to null
- when adding, next $\&$ prev must be null prior
 - else throw exception

Sorting

Insertion Sort

- "sorted" section \nexists "unsorted"
- take each element from "unsorted" section
- traverse "sorted" section backwards \nexists find the spot the current element should be inserted into

Ex: 5 2 7 3 6 1

* Start w/ 2
* go through sorted

* start w/
 * go through sorted

5 | 2 7 3 6 1
↑ ↑ ; ↑

insert here

2 5 | [7] 3 6 1
S : U

2 5 ; 7 3 6 1
↑ ; ↑

belongs here

2 5 7 | [3] 6 1
S : U

2 5 7 ; 3 6 1
↑ ↑ ; ↑ ; ↑

after no not here

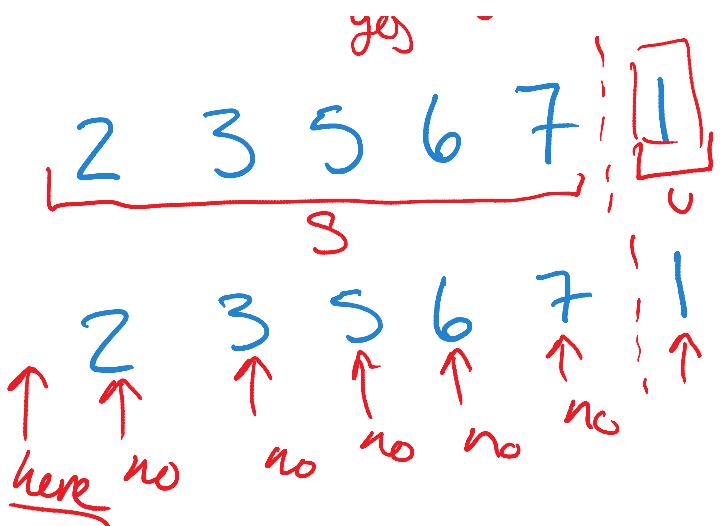
2 3 5 7 | [6] 1
S : U

2 3 5 7 ; 6 1
↑ ; ↑ ; ↑

yes

no

- - , - 1 1



1 2 3 5 6 7

* runtime? $O(n^2)$

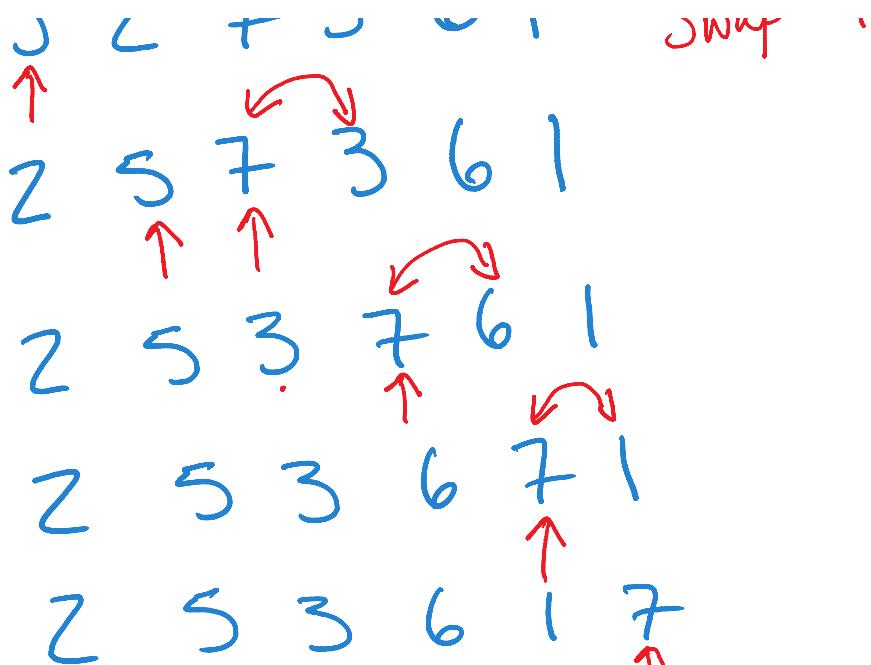
→ for each element, look
@ rest

* if already sorted $O(n)$

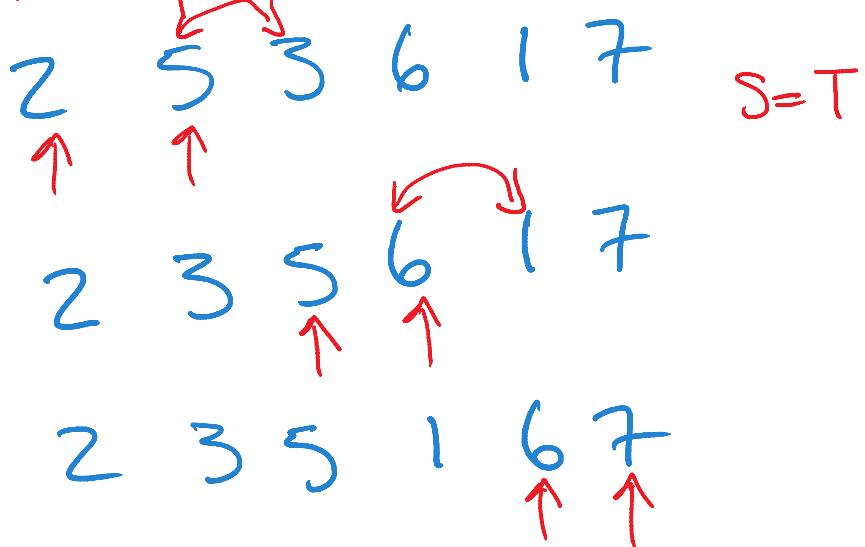
Bubble Sort

- very simple \Rightarrow inefficient
- start at beginning
- if 2 next to each other are out of order, swap
- keep swapping until no more

5 2 7 3 6 | swap: false
↑ ↗ ↘ Swap: T

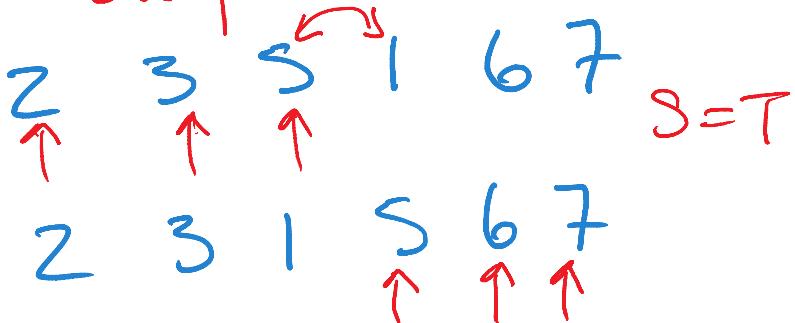


*Swap is true - restart $S=F$

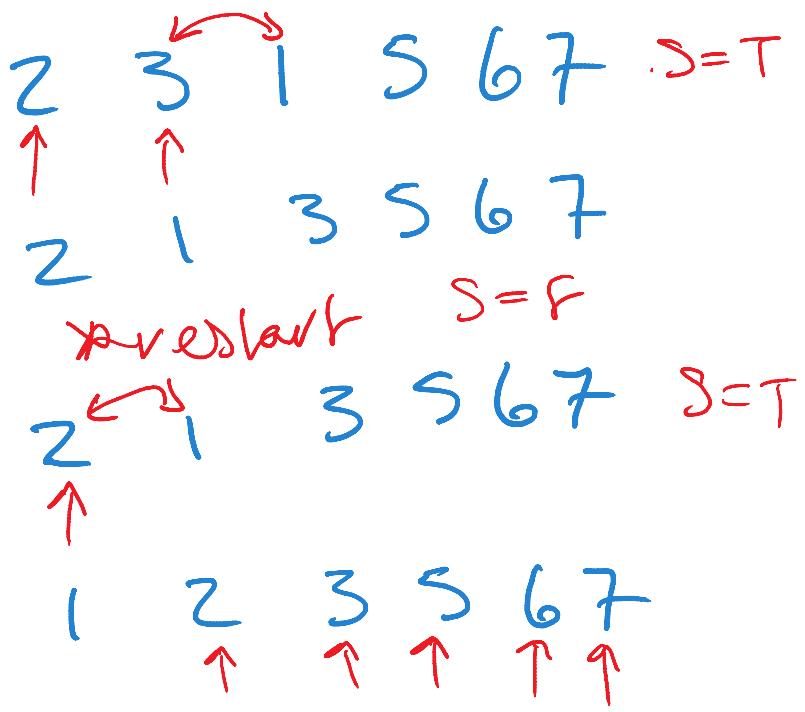


*Swap == T : Start over

Swap $\rightarrow F$

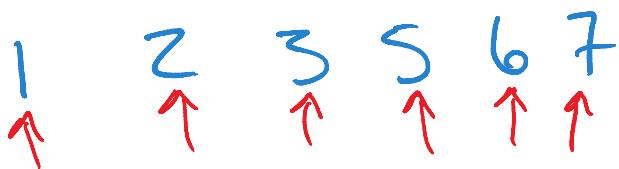


*restart $S=F$



* $S = T$, restart

$S=F$



* done.

* run time: $O(n^2)$

best: $O(n)$

Comparator

- generic interface to compare objects
- create classes that implement ... comparator

- create classes "new" T
- comparator
- single method: compare(o1, o2)
- returns an int
 - < 0 if o1 < o2
 - = 0 if o1 = o2
 - > 0 if o1 > o2

Anonymous Class

- can create an anonymous class - one that only exists in a specific context
- instantiate when used
- has no name - cannot be used elsewhere
- common to make some objects, particularly comparators, as anonymous classes

Ex:

Comparator<String> c = new
Comparato<String> c = new
Comparato<String> c = new
Comparato<String> c = new

```
Comparator<String> c = new
```

```
    Comparator<String>() {
```

```
        public int compare(String s1, String s2) {
```

```
            //etc
```

```
        }
```

```
    };
```

"Pass the Buck" Recursion

→ if it isn't your problem,
pass it on to someone else

* don't update arguments

→ pass entire problem to
someone else

→ if you can solve the
problem, solve it.

* make sure someone can solve
it & you don't pass it back
to the person who originally
had the problem.