

Linked Lists

Wednesday, February 12, 2020 10:05 PM

Components

- many items
- head, tail, cursor/pre
- node
 - data

→ next pointer
*nested static class

- Static: cannot see anything in the class if it is nested
- unlike iterator, which is a nested non-static class
- can see collection

*What should next point to if it is the last in the list?

→ null

have ^{sorted} linked list w/
head & tail

```
private static node<int> L;
```

Invariant?

- *many Nodes is correct
- *tail must be last in list
- *start(ed) @ head
- *if head is null, tail must be null
- *no cycles
- *pre cursor is null or points to a node in this list

```

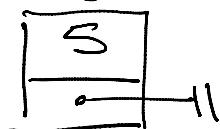
int data,
Node next;
public Node(int d, Node n){
    data=d;
    next=n;
}

```

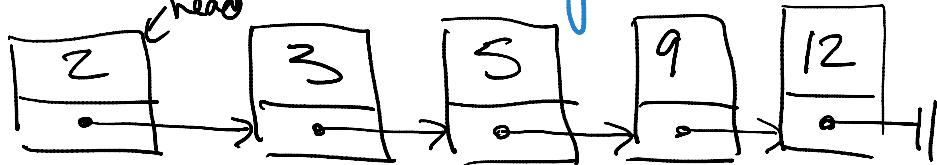
\approx
head = null; tail = null
cursor = null;

add a 5 to empty list

tail ←
head = new Node(5, null);
 \downarrow head

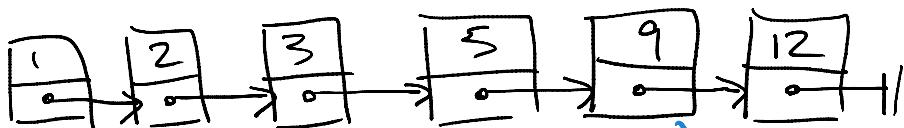


add to a longer list



add 1 (at beginning)

head = new Node(1, head);



add 15 (at end)

w/ tail:

tail = tail.next = new Node(15, null)

w/out tail:

Start @ head $\xrightarrow{\text{find}} \text{tail}$

```
for(Node n = head; n.next != null; n = n.next);  
n.next = new Node(15, null);
```

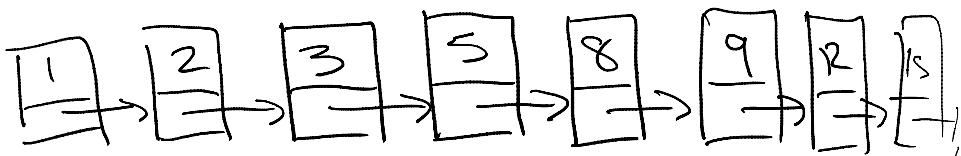
* add in between

\rightarrow add 8

* what do we need?

\rightarrow ptr to node that comes before 8.

```
for(Node n = head; n.next != null;  
n = n.next)  
if(n.next.data > 8) break;  
n.next = new Node(8, n.next);
```



* What is the runtime if we are not guaranteed to add at the head?

$O(n)$

*what if we always add
at the head?

$O(1)$

*think about the ball
Sequence. What does append
do?

→ if current elem, add
after
→ else, at head

*what was the runtime
w/ a dynamic array?

$O(n)$ *shift

*what is runtime w/
linked list ↳ cursor?

$O(1)$

*removal from beginning

head = head.next;

*why is this sufficient?

• n • .. → no pointer

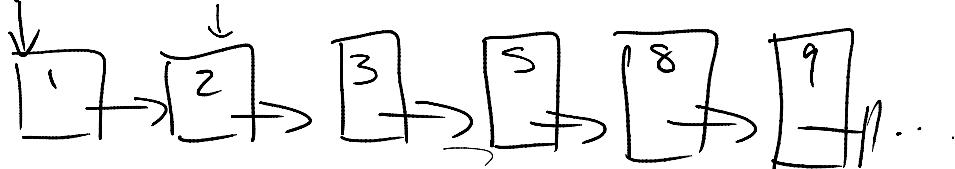
* Why is this sufficient?

- if there is no reference to an object, the garbage collector gets rid of it.
- the old head will be removed

* Remove from end or middle?

- need to find the prev node

head remove(3)



for (Node n = head; n.next != null;

 n = n.next)

 if (n.next.data == target) {

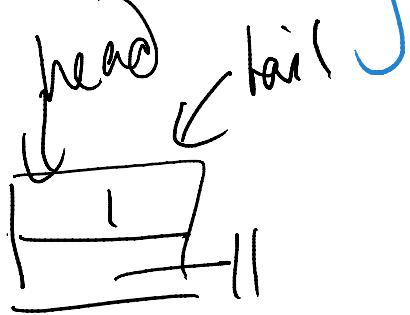
 n.next = n.next.next;
 return true;

}

 return false;

* Remove only item?

remove only from



$\text{head} = \text{tail} (= n \parallel)$

* What is the runtime?

$O(n)$

* What if only remove
at head (stack) or
w/ a cursor?

$O(1)$

array w/cursor

linked list w/cursor

add	$O(n)$	$O(1)$
remove	$O(n)$	$O(1)$

remove	$O(n)$	$O(1)$
getCurrent	$O(1)$	$O(1)$
get i-th	$O(1)$	$O(n)$

* pros of linked lists?

→ faster insert/remove

→ no need to resize

* pro of array?

→ fast random access

• watch out for null

pointers!

→ if list is empty, all
pointers are null

→ last element's next

is null

→ make sure to
check

→ if cursor is at the
head → precursor is null!
- make sure you

new → previous ...
→ make sure you
handle that.

* how do you handle clone?

→ clone ~~creates~~ fields, but
doesn't create a new list.