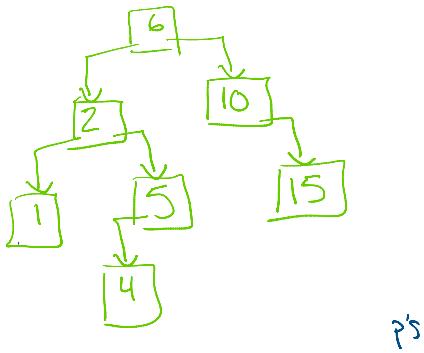
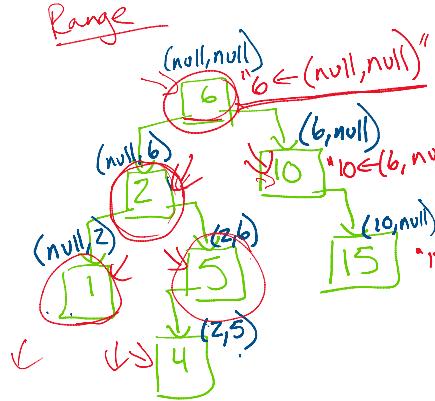


- root
- nodes
 - left
 - right
- * 2 rules:
 - left always less than parent
 - right always greater than parent



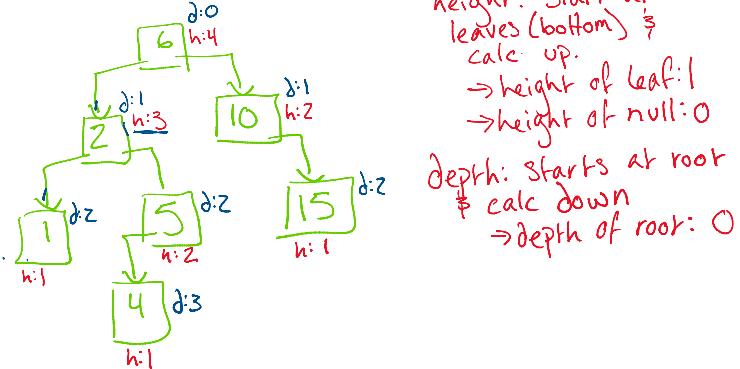
p's



left's range:
 → low: parent's low
 → high: parent
 right's range:
 → low: parent
 → high: parent's high

do TreeRange(n, left,
 left's low,
 left's hi,
 set)

$$\text{height} = \max(\text{children's height}) + 1$$



height: Start at leaves (bottom) & calc up.
 → height of leaf: 1
 → height of null: 0

depth: Starts at root & calc down
 → depth of root: 0

Recursion:
 → call method inside itself

int sum(int n) → return sum of 0 to n

if (n == 0) return 0;

return n + sum(n-1);

→ if passing info down:
 → put in parameter

→ if passing info up:
 → put in return

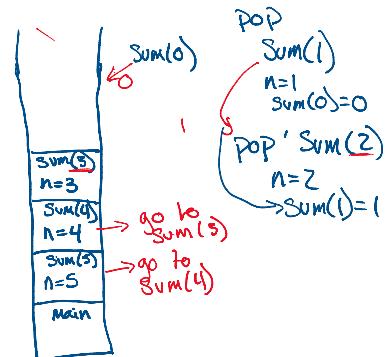
→ USE recursion to solve problems w/ subproblems
 → Subproblems almost identical

Recursion tree

$$\begin{aligned} \text{Sum}(5) &\rightarrow 5 + \text{sum}(4) = 5 + 10 + 15 \\ \downarrow \\ \text{Sum}(4) &\rightarrow 4 + \text{sum}(3) = 4 + 6 + 10 \\ \downarrow \\ \text{Sum}(3) &\rightarrow 3 + \text{sum}(2) = 3 + 3 + 6 \\ \downarrow \\ \text{Sum}(2) &\rightarrow 2 + \text{sum}(1) = 2 + 3 \\ \downarrow \\ \text{Sum}(1) &\rightarrow 1 + \text{sum}(0) = 1 \\ \downarrow \\ \text{Sum}(0) &\rightarrow 0 \end{aligned}$$

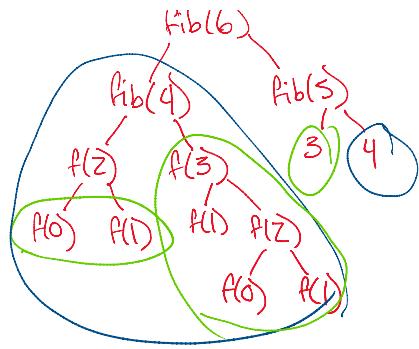
$$\text{fib}(n) = \underline{\text{fib}(n-2)} + \underline{\text{fib}(n-1)}$$

0, 1, 1, 2, 3, 5, 8 ...



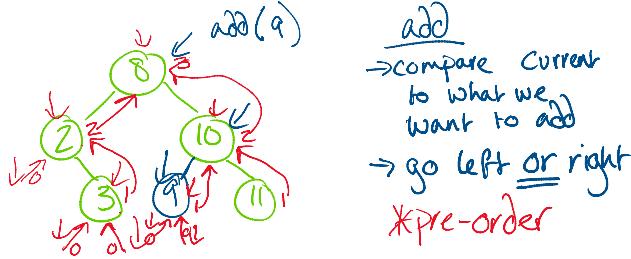
Call Stack
 * code
 * "state" / variables

→ use recursion
problems w/ subproblems
→ subproblems almost identical
* don't use if recalculating
something you already know



Traversals

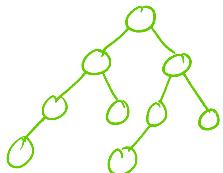
- pre-order:
 - process current node
 - go left
 - go right
- * children need info from parent
* check in Range
- in-order:
 - go left
 - process current
 - go right
- * post-order:
 - go left
 - go right
 - process current
- * need info from children
* heights



add
→ compare current to what we want to add
→ go left or right
* pre-order

Runtime

best:



if going down
↓
1 subtree
(add, contains)
→ $O(\log n)$
* dividing # of nodes in half

worst:

$O(n)$

* if doing operation that looks at every node

→ doHeight
→ checkInRange

$O(n)$