

CSCI 6461 Computer Systems Architecture

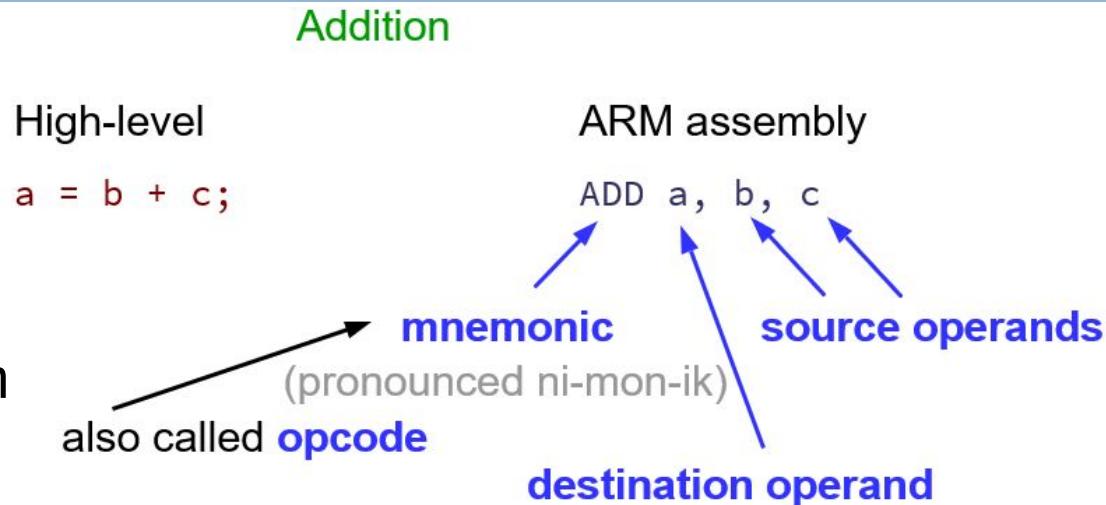
Lecture 3

ISA: Basic Instructions

Addition and Subtraction Ops

Note the structure of these instructions:

instr to, from, from



The **mnemonic** indicates what operation to perform

Subtraction

High-level
 $a = b - c;$

ARM assembly
SUB a, b, c

More complex example

More complex high-level code translates into multiple ARM instructions

```
a = b + c - d; //comment
```

ARM Assembly code:

```
ADD t, b, c ; t = b + c  
SUB a, t, d ; a = t - d
```

**we are not naming our registers correctly yet - this is just an example

Operands

- An instruction operates on **operands (a and b here)** .
 - ADD a, b, c
- Operands can be stored in registers or memory, or they may be constants stored in the instruction itself.
- Operands stored as constants or in registers are accessed quickly, but they hold only a small amount of data.
- Additional data must be accessed from memory, which is large but slow.

Register Naming Convention

ARM register names are preceded by the letter R

```
; R0 = a, R1 = b, R2 = c  
ADD R0, R1, R2 ; a = b + c
```

The following ARM assembly code uses a register, R4,
to store the intermediate calculation:

```
ADD R4, R1, R2 ; t = b + c  
SUB R0, R4, R3 ; a = t - d
```

Example

C to Assembly

C to Assembly Example

Translate the following high-level code into ARM assembly language.

```
a = b - c;  
f = (g + h) - (i + j);
```

Solution

```
; R0 = a,    R1 = b,    R2 = c,    R3 = j  
; R4 = g,    R5 = h,    R6 = i,    R7 =  
SUB R0, R1, R2 ; a = b - c  
ADD R8, R4, R5 ; R8 = g + h  
ADD R9, R6, R7 ; R9 = i + j  
SUB R3, R8, R9 ; f = (g + h) - (i + j)
```

Note

- We have not yet discussed how R0, R1 etc were initialized with values.
- Also, we used 9 registers in the previous example.
- This is probably wasteful
- In general, we will try to reduce the number of registers used in a particular stage

Constants (Immediates)

- In addition to register operations, ARM instructions can use constant or immediate operands.
- These constants are called immediates, because their values are immediately available from the instruction and do not require a register or memory access.

High-level

```
a = a + 4;  
b = a - 12;
```

ARM assembly

```
; R7 = a, R8 = b  
ADD R7, R7, #4  
SUB R8, R7, #0xC
```

Constants

- Hexadecimal constants in ARM assembly start with `0x`, for example
 - `MOV R4, #0xFA65`
- Single character constants, for example
 - `MOV R5, #'A'`

Logical Instructions

MVN – MoVe and Not

ORR – OR

EOR – XOR

BIC – bit clear

Source registers			
R1	0100 0110	1010 0001	1111 0001
R2	1111 1111	1111 1111	0000 0000

Assembly code

AND R3, R1, R2

ORR R4, R1, R2

EOR R5, R1, R2

BIC R6, R1, R2

MVN R7, R2

Result			
R3	0100 0110	1010 0001	0000 0000
R4	1111 1111	1111 1111	1111 0001
R5	1011 1001	0101 1110	1111 0001
R6	0000 0000	0000 0000	1111 0001
R7	0000 0000	0000 0000	1111 1111

- The 1st source is always a register
- The 2nd source is either an immediate or another register
- BIC clears the bits that are asserted in R2

Shift Instructions

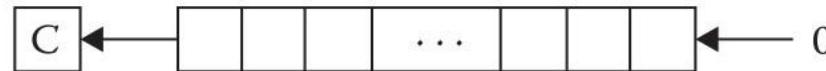
ASR – the sign bit shifts into the most significant bits.

The carry flag is updated to the last bit shifted out

LSL

Logical shift left by n bits

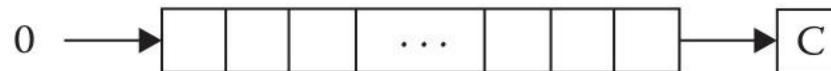
Multiplication by 2^n



LSR

Logical shift right by n bits

Unsigned division by 2^n



ASR

Arithmetic shift right by n bits

Signed division by 2^n

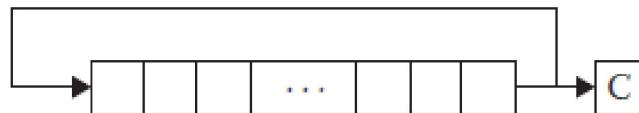


Rotate Right

ROR

Rotate right by n bits

32-bit rotate



- The amount by which to shift can be an immediate or a register.
- We can use ROR to invert the bit pattern of a register
 - $R0 = 111000$
 - $RO, R0, \#3$
 - $R0 = 000111$
- Again, the carry flag is updated to the **last** bit rotated out

Examples, 1/2

Source register

R5	1111 1111	0001 1100	0001 0000	1110 0111
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Assembly Code

LSL R0, R5, #7 R0

LSR R1, R5, #17 R1

ASR R2, R5, #3 R2

ROR R3, R5, #21 R3

Result

1000 1110	0000 1000	0111 0011	1000 0000
0000 0000	0000 0000	0111 1111	1000 1110
1111 1111	1110 0011	1000 0010	0001 1100
1110 0000	1000 0111	0011 1111	1111 1000

Examples, 2/2

Source registers				
R8	0000 1000	0001 1100	0001 0110	1110 0111
R6	0000 0000	0000 0000	0000 0000	0001 0100

Assembly code

LSL R4, R8, R6

ROR R5, R8, R6

Result

R4	0110 1110	0111 0000	0000 0000	0000 0000
R5	1100 0001	0110 1110	0111 0000	1000 0001

No ASL

- There is no ASL, or an arithmetic shift left
- You would never need such an instruction,
 - since arithmetic shifts need to preserve the sign bit,
 - and shifting signed data to the left will do so as long as the number doesn't overflow.
 - for example:

-1 is 0xFFFFFFFF

shifting it left results in 0xFFFFFFF, which is -2 and correct

LSL as an operand

- When the final argument of basic arithmetic instruction is a register
 - we can optionally add a shift distance operator
 - LSL/LSR
 - when used in this mode, the shift operator does not change the register value it operates on
 - it merely uses the value as a parameter to the operand

LSL as a parameter

```
ADD R0, R0, R1, LSL #1
```

- This adds a left-shifted version of R1 before adding it to R0.
- R1 itself remains unchanged

The shift distance can be an immediate between 1 and 32, or it can be based on a register value:

```
MOV R0, R1, ASR R2
```

Example

```
MOV R0, #10  
MOV R1, #20  
ADD R0, R0, R1, LSL #1
```

What is the state of R0 and R1 at the end of this code?

Initializing values using immediates

MOV is a useful way to initialize register values:

High-level

```
i = 0;  
x = 4080;
```

ARM assembly

```
; R4 = i, R5 = x  
MOV R4, #0  
MOV R5, #0xFF0
```

The instruction MOV R5, #0xFF1 generates an error

You can overcome this using

```
MOV R5, #0xFF0  
ADD R5, #0x1
```

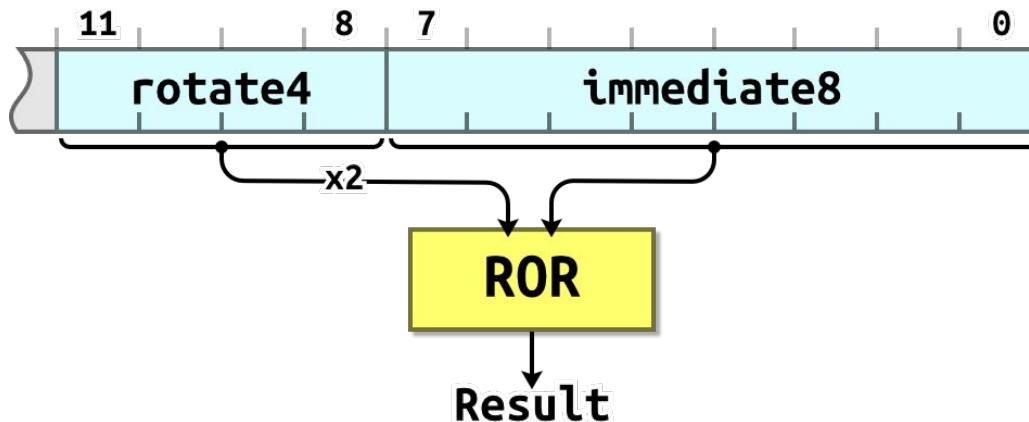
Another method:

```
MOV R5, #0xFF0  
ORR R5, R5, #0x001
```

Literal Operands -

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- You can't fit an arbitrary 32-bit value into a 32-bit instruction word.
- ARM data processing instructions have 12 bits of space for values in their instruction word. This is arranged as a four-bit rotate value and an eight-bit immediate value:



Literal Operands -

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- The 4-bit rotate value stored in bits 11-8 is multiplied by two giving a range of 0-30 in steps of two.
- Using this scheme we can express immediate constants such as:
 - 0x000000FF
 - 0x00000FF0
 - 0xFF000000
 - 0xF000000F
- But immediate constants such as:
 - 0x000001FE
 - 0xF000F000
 - 0x55550000
- ...are not possible.

Multiply Instructions

MUL R1, R2, R3

- multiplies the values in R2 and R3
- and places the least significant bits of the product in R1
- the most significant 32 bits of the product are discarded

UMULL – unsigned multiply long

SMULL – signed multiply long

UMULL/SMULL

Multiply two 32-bit numbers and produce a 64-bit product

Example:

```
UMULL R1, R2, R3, R4
```

- Performs an unsigned multiply of R3 and R4
- The least significant 32 bits of the product is placed in R1 The most significant 32 bits are placed in R2

Multiply-accumulate instructions

MLA r7, r8, r9, r3

; $r7 = r8 * r9 + r3$

the least significant 32 bits of the result

SMLAL r4, r8, r2, r3 ; $\{r8, r4\} = r2 * r3 + \{r8, r4\}$

UMLAL r5, r8, r0, r1 ; $\{r8, r5\} = r0 * r1 + \{r8, r5\}$

MS32bits LS32bits

These instructions can boost the math performance in applications such as matrix multiplication and signal processing consisting of repeated multiplies and adds.

RSB

The **RSB** (Reverse SuBtract) instruction subtracts the value in **Rn** from the value of **Operand2** . This is useful because of the wide range of options for Operand2 .

RSB Rd, Rn, Src2 ; Rd \leftarrow Src2 – Rn

Example 1

Consider the following operation:

```
SUB r0, r2, r3, LSL #2 ; r0 = r2 - r3*4
```

Suppose we want modify (shift) register r2 before the subtraction instead of register r3.

This is done using the reverse subtract operation

```
RSB r0, r3, r2, LSL #2 ; r0 = r2*4 - r3
```

Example 2

An assembly program to perform the function of absolute value with only two instructions. Solution:

```
MOVS r1,r0
```

```
RSBLT r1, r1, #0
```

- MOVS performs the same function as MOV, but also updates the N and Z flags.
- Perform a reverse subtract if r1 is negative. In other words: $r1 = 0 - r1$ if $r1$ negative

Consider the following instruction:

```
// r0 = r1 + r1*4 = r1*5  
ADD r0, r1, r1, LSL #2;
```

- Why do it this way instead of using MUL?
- Answer: the size and power usage of a multiplier array are large.
- In very low power applications, it's often necessary to play every trick in the book to save power.

MRS/MSR "Special Registers"

MRS (Move PSR to general-purpose register)
instruction to read the flags

MSR (Move general-purpose register to PSR)
instruction to write the flags

```
MRS r0, CPSR ; load the contents of the CPSR into r0  
MRS r1, SPSR ; load the contents of the SPSR into r1
```

From there, you can examine any flags that you like. You cannot use register r15 as the destination register.