Project #1

My solution

Projection-based method

- Items are listed in lexicographic order
- Let P and DB(P) be a node's pattern and its associated projected database.
- Mining is performed by recursively calling this function:
 - TP(P, DB(P))
 - 1. Determine the frequent items in DB(P), and denote them by E(P).
 - 2. Eliminate from DB(P) any items not in E(P).
 - 3. For each item x in E(P), call TP(Px, DB(Px)).

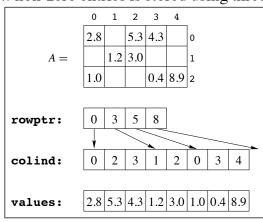
What is the right data structure for storing the projected database?

- Operations that needs to support:
 - Determining the frequency of each item.
 - Projection:
 - Extract the transactions and their subsets of items that support a given pattern.
- What should be the upper bound on the complexity of each of these operations?

Dense matrices are stored in the computer memory by using two-dimensional arrays. For example, a matrix with n rows and m columns, is stored using a $n \times m$ array of real numbers. However, using the same two-dimensional arrays to store sparse matrices has two very important drawbacks. First, since most of the entries in the sparse matrix are zero, this storage scheme wastes a lot of memory. Second, computations involving sparse matrices often need to operate only on the non-zero entries of the matrix. Use of dense storage format makes it harder to locate these non-zero entries. For these reasons sparse matrices are stored using different data structures.

The compressed storage format (CSR) is a widely used scheme for storing sparse matrices. In the CSR format, an $n \times n$ sparse matrix A that has k non-zero entries is stored using three

arrays: two integer arrays rowptr and colind, and one array of real values values. The array rowptr is of size n+1, and the other two arrays are each of size k. The array colind stores the column-indices of the non-zero entries in A, and the array values stores the corresponding non-zero entries. In particular, the array colind stores the column-indices of the first row followed by the column-indices of the second row followed by the column-indices of the third row, and so on. Similarly, the array values stores the corresponding non-zero entries of the first row followed by the



CSR format of a sample matrix

corresponding non-zero entries of the second row, and so on. The array rowptr is used to determine where the storage of a row starts and ends in the arrays colind and values. In particular, the column-indices of row i are stored starting at colind[rowptr[i]] and ending at (but not including) colind[rowptr[i+1]]. Similarly, the values of the non-zero entries of row i are stored starting at values[rowptr[i]] and ending at (but not including) values[rowptr[i+1]]. Also note that the number of non-zero entries for row i is simply rowptr[i+1]-rowptr[i].

Figure 0.11: The compressed storage format for sparse matrices.

My naming of the CSR/CSC fields is different from the common way of doing things.

I use "row" as the prefix for the row-wise representation (rowptr, rowind, rowval).

I use "col" as the prefix for the column-wise representation (colptr, colind, colval).

```
/*! Data structures for use within this module */
                /*----*/
                typedef struct {
                  int minfreg; /* the minimum frequency of a pattern */
                  int maxfreq; /* the maximum frequency of a pattern */
                  int minlen; /* the minimum length of the requested pattern */
                  int maxlen; /* the maximum length of the requested pattern */
                  int tnitems; /* the initial range of the item space */
                  /* the call-back function */
                  void (*callback)(void *stateptr, int nitems, int *itemids, int ntrans, int *transids);
The memory for
                  void *stateptr; /* the user-supplied pointer to pass to the callback */
those is allocated
once and reused.
                  /* workspace variables */
                  int *rmarker;
                  gk_ikv_t *cand;
                } isparams_t;
```

void gk_find_frequent_itemsets(int ntrans, ssize_t *tranptr, int *tranind, int minfreq, int maxfreq, int minlen, int maxlen, void (*process_itemset)(void *stateptr, int nitems, int *itemids, int ntrans, int *transids), void *stateptr) ssize_t i; gk_csr_t *mat, *pmat; isparams_t params; int *pattern; /* Create the matrix */ mat = gk_csr_Create(); mat->nrows = ntrans; mat->ncols = tranind[qk_iarqmax(tranptr[ntrans], tranind, 1)]+1; mat->rowptr = gk_zcopy(ntrans+1, tranptr, gk_zmalloc(ntrans+1, "gk_find_frequent_itemsets: mat.rowptr")); mat->rowind = gk_icopy(tranptr[ntrans], tranind, gk_imalloc(tranptr[ntrans], "gk_find_frequent_itemsets: mat.rowind")); mat->colids = gk_iincset(mat->ncols, 0, gk_imalloc(mat->ncols, "gk_find_frequent_itemsets: mat.colids")); /* Setup the parameters */ params.minfreg = minfreg; params.maxfreq = (maxfreq == -1 ? mat->nrows : maxfreq); params.minlen = minlen; params.maxlen = (maxlen == -1 ? mat->ncols : maxlen);params.tnitems = mat->ncols; params.callback = process_itemset; params.stateptr = stateptr; params.rmarker = gk_ismalloc(mat->nrows, 0, "gk_find_frequent_itemsets: rmarker"); = gk_ikvmalloc(mat->ncols, "gk_find_frequent_itemsets: cand"); /* Perform the initial projection */ gk_csr_CreateIndex(mat, GK_CSR_COL); pmat = itemsets_project_matrix(¶ms, mat, -1); gk_csr_Free(&mat); pattern = gk_imalloc(pmat->ncols, "gk_find_frequent_itemsets: pattern"); itemsets_find_frequent_itemsets(¶ms, pmat, 0, pattern);

gk_free((void **)&pattern, ¶ms.rmarker, ¶ms.cand, LTERM);

/ / total contration to total contration to the contration of the c

/*! The entry point of the frequent itemset discovery code */

This is to keep track of the original item-IDs, since their numbering will change due to reordering.

gk_csr_Free(&pmat);

```
/*! The recursive routine for DFS-based frequent pattern discovery */
void itemsets_find_frequent_itemsets(isparams_t *params, gk_csr_t *mat,
       int preflen, int *prefix)
 ssize_t i;
 gk_csr_t *cmat;
 /* Project each frequent column */
 for (i=0; i<mat->ncols; i++) {
   prefix[preflen] = mat->colids[i];
   if (preflen+1 >= params->minlen)
    (*params->callback)(params->stateptr, preflen+1, prefix,
        mat->colptr[i+1]-mat->colptr[i], mat->colind+mat->colptr[i]);
   if (preflen+1 < params->maxlen) {
    cmat = itemsets_project_matrix(params, mat, i);
    itemsets_find_frequent_itemsets(params, cmat, preflen+1, prefix);
    gk_csr_Free(&cmat);
   }
 }
```

```
gk_csr_t *itemsets_project_matrix(isparams_t *params, gk_csr_t *mat, int cid)
  ssize_t i, j, k, ii, pnnz;
  int nrows, ncols, pnrows, pncols;
  ssize_t *colptr, *pcolptr;
  int *colind, *colids, *pcolind, *pcolids, *rmarker;
  gk_csr_t *pmat;
  gk_ikv_t *cand;
  nrows = mat->nrows;
  ncols = mat->ncols;
  colptr = mat->colptr;
  colind = mat->colind;
  colids = mat->colids;
  rmarker = params->rmarker;
         = params->cand;
  cand
 /* Allocate space for the projected matrix based on what you know thus far */
  pmat = gk_csr_Create();
  pmat->nrows = pnrows = (cid == -1 ? nrows : colptr[cid+1]-colptr[cid]);
  /* Mark the rows that will be kept and determine the prowids */
 if (cid == -1) { /* Initial projection */
    gk_iset(nrows, 1, rmarker);
  else { /* The other projections */
   for (i=colptr[cid]; i<colptr[cid+1]; i++)</pre>
      rmarker[colind[i]] = 1;
 }
  /* Determine the length of each column that will be left in the projected matrix */
  for (pncols=0, pnnz=0, i=cid+1; i<ncols; i++) {
    for (k=0, j=colptr[i]; j<colptr[i+1]; j++) {</pre>
      k += rmarker[colind[j]];
                                                                      This is either 0 or 1.
   if (k >= params->minfreq && k <= params->maxfreq) {
      cand[pncols].val = i;
     cand[pncols++].key = k;
      pnnz += k;
   }
  }
 /* Sort the columns in increasing order */
  gk_ikvsorti(pncols, cand);
```

pncols: is the number of columns in the projected database.

pnnz: is the total number of items in the transactions of the projected database.

```
/* Sort the columns in increasing order */
gk_ikvsorti(pncols, cand);
/* Allocate space for the remaining fields of the projected matrix */
pmat->ncols = pncols;
pmat->colids = pcolids = gk_imalloc(pncols, "itemsets_project_matrix: pcolids");
pmat->colptr = pcolptr = qk zmalloc(pncols+1, "itemsets project matrix: pcolptr");
pmat->colind = pcolind = gk_imalloc(pnnz, "itemsets_project_matrix: pcolind");
/* Populate the projected matrix */
pcolptr[0] = 0;
for (pnnz=0, ii=0; ii<pncols; ii++) {</pre>
  i = cand[ii].val;
  for (j=colptr[i]; j<colptr[i+1]; j++) {</pre>
    if (rmarker[colind[j]])
      pcolind[pnnz++] = colind[j];
  }
  pcolids[ii] = colids[i];
  pcolptr[ii+1] = pnnz;
}
/* Reset the rmarker array */
if (cid == -1) { /* Initial projection */
  gk_iset(nrows, 0, rmarker);
}
else { /* The other projections */
                                                                             This optimization is
  for (i=colptr[cid]; i<colptr[cid+1]; i++)</pre>
                                                                             essential for sparse
    rmarker[colind[i]] = 0;
                                                                             data.
}
return pmat;
```