

# Conceptual Database Design

# 1. Conceptual Database Design



Database Design

# 1.2 Entity-Relationship Model

- The Entity-Relationship (ER) model is a high-level conceptual data model (Chen in 1966).
- ER is used mainly as a design tool.

## 1.2 Entity-Relationship Model<sub>(cont)</sub>

- *Entity type*: Group of object with the same properties
- *Entity*: member of an entity type - analogous to an object.
- *Attribute*: a property of object
- *Relationship*: among objects
  - ER can model “n-way” relationship,
  - ER models a relationship and its inverse by a single relationship.

## 1.2.1 Entity and Attributes

- *Entities* represent things in the real world.
- *Attributes* describe properties of entities.
- Attributes may be
  - simple(atomic) e.g. sex = 'Female', or
  - composite e.g. name consists of title (Dr), Initials (C.C.), family name (Chen).

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- Each entity has values for each attribute.
- Attributes may be
  - *single-valued* e.g. student number, name, or
  - *multivalued* e.g. keywords = neural networks, computer graphics, databases.

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- Each simple attribute has a *value set (domain)*: the set of possible values for that attribute.
- In a composite attribute  $A = (A_1, \dots, A_n)$ , suppose that  $V_1, \dots, V_n$  are the domains of  $A_1, \dots, A_n$ .
- The domain  $V$  of  $A$  is  $V_1 \times \dots \times V_n$ .
- Mathematically, an attribute  $A$  of an entity type  $E$  is a function

$$A : E \rightarrow \wp(V) .$$

- where  $V$  is the domain of  $A$ , and  $\wp(V)$  is the power set of  $V$
- For single-valued attributes,  $A(e)$  must be a singleton.

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- An attribute can have a null value if, for example:
  - there is no suitable value e.g. a student may have no interests: keywords = NULL
  - the true value is not known e.g. the marriage date of a person is not known: marriage date = NULL.
- A derived attribute is one whose value can be derived from other attributes and entities. e.g. number of students.



## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- An *entity* type is a set of entities with the same attributes.
- It is described by an *entity* schema: a name and a list of attributes.
- The set of individual entity *instances* at a particular moment in time is called an extension of the entity type.

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

| Schema<br>(Intension)    | RESEACHER<br>Name, Payroll_no, No_of_students,<br>Keywords                                | DEPARTMENT<br>Name                           |
|--------------------------|---|--|
| Instances<br>(Extension) | (Dr C.C. Chen, 230-0013, 3, Neural Networks)<br>(Dr R. Wilkinson, 231-0091, 1, Databases) | Computer Science<br>Psychology<br>Management |

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- An entity type usually has a *key*: a set of attributes that uniquely identifies an entity. For example:
  - {payroll number} is a key of RESEARCHER,
  - {name} is a key of DEPARTMENT.
- There may be more than one possible key.
- An important constraint is the key constraint: in any extension of the entity type, there cannot be two entities having the same values for their key attributes.

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- We can describe schemata with composite attributes using ()'s and with multi-valued attributes using {}'s. e.g.

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

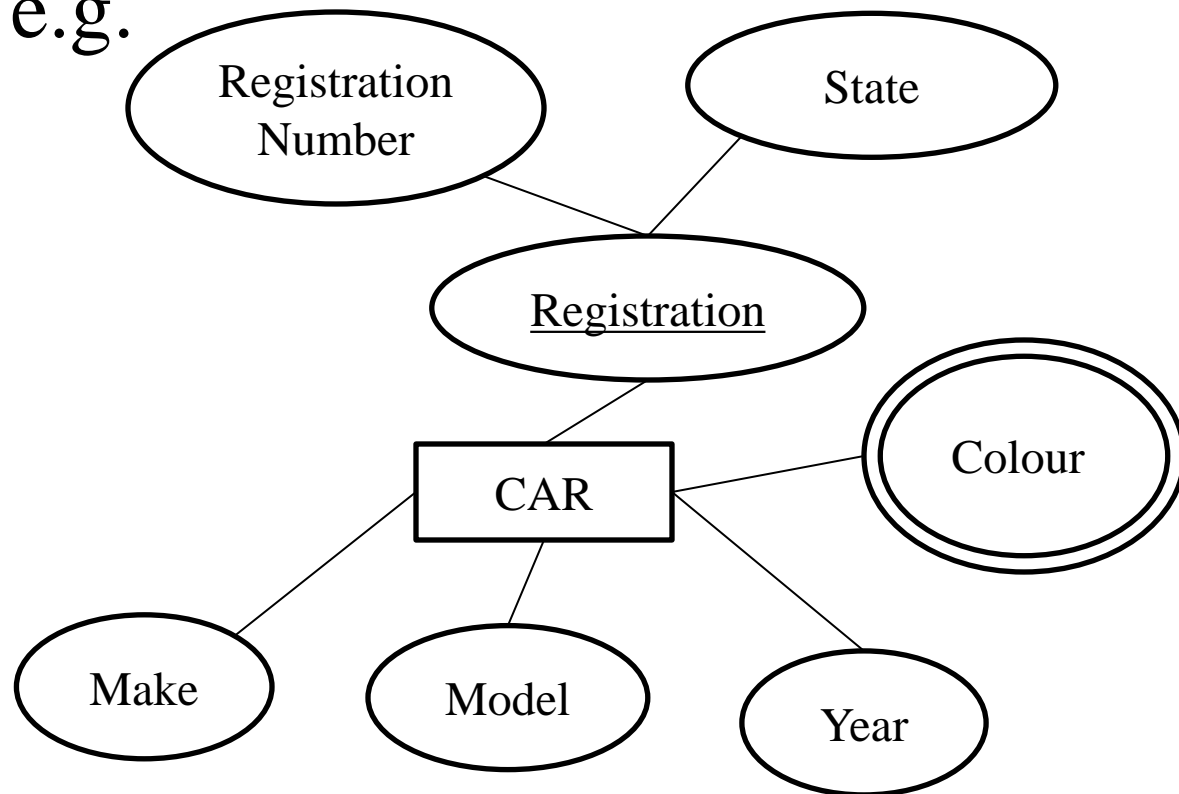
### CAR

**Registration(Registration No, State), Make, Model, Year, {Colour}**

((ARQ) 595, Vic), Datsun, 120Y, 1972, {green})  
((8HR) 696, WA), Mazda, 929, 1979, {grey, black})

## 1.2.1 Entity and Attributes<sub>(cont)</sub>

- Entities and their attributes can also be described with Entity-Relationship Diagrams (ERDs). e.g.



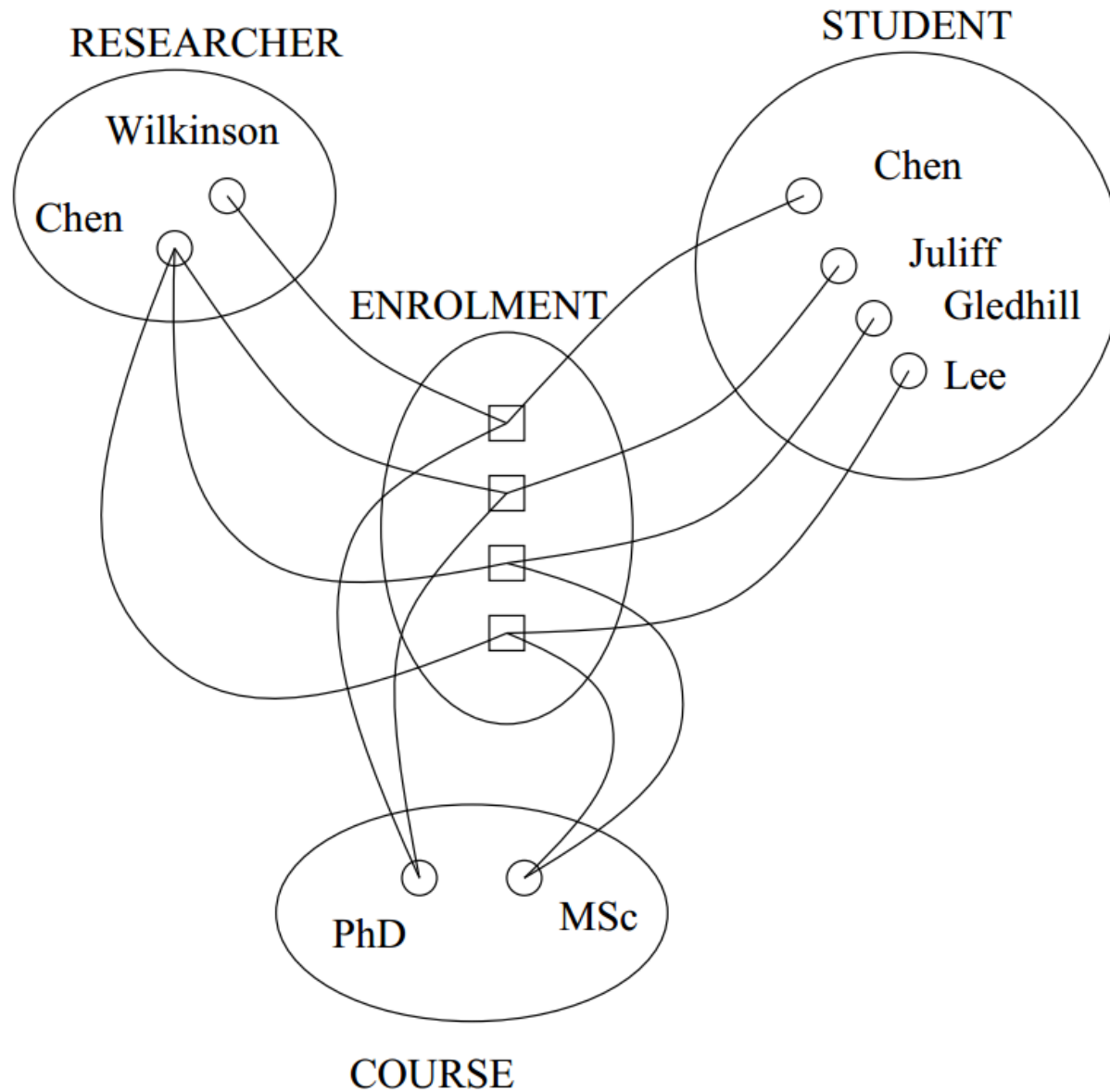
## 1.2.2 Relationships

- A *relationship* represents an association between things.
- A *relationship* type  $R$  among  $n$  entity types  $E_1, \dots, E_n$  is a set of associations among entities from these types.
- Mathematically, a relationship type  $R$  among entity types  $E_1, \dots, E_n$  is a subset of  $E_1 \times \dots \times E_n$ .
- Each instance  $r = (e_1, \dots, e_n)$  in  $R$  is a relationship.

## 1.2.2 Relationships<sub>(cont)</sub>

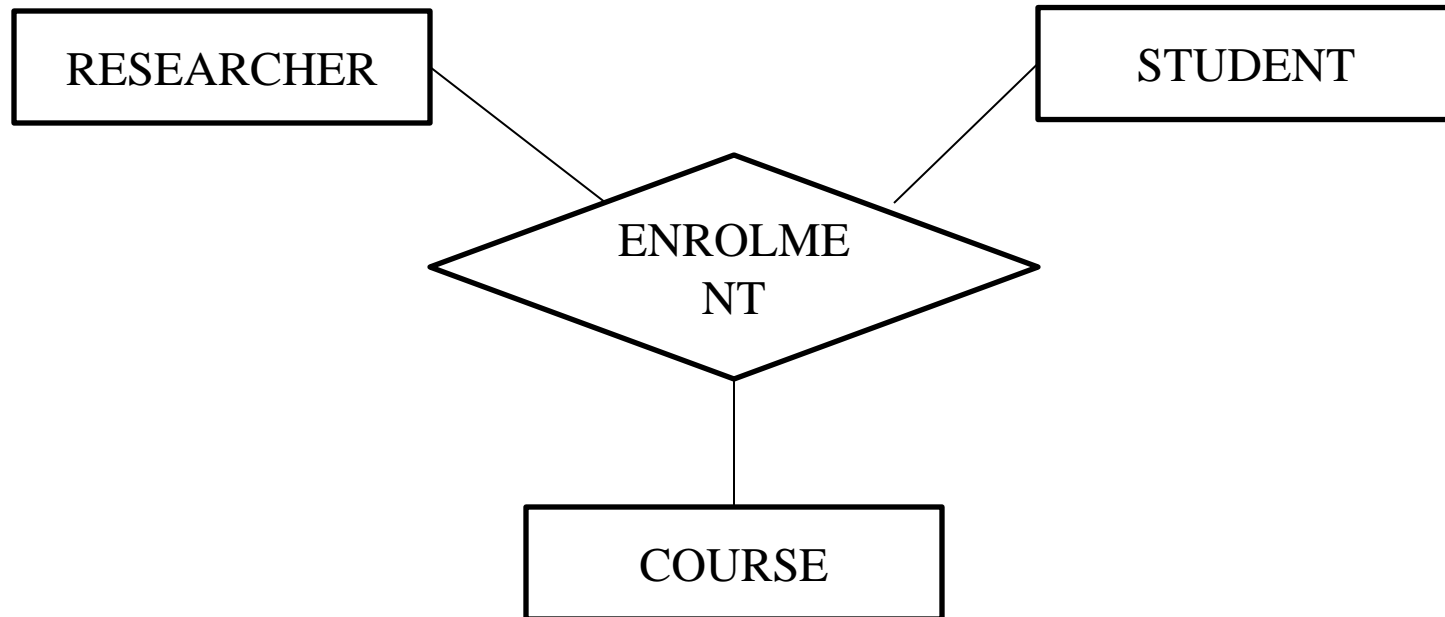
- We say that  $E_1, \dots, E_n$  participate in  $R$ .
- Similarly if  $r = (e_1, \dots, e_n)$  is an instance of  $R$ , we say that each  $e_i$  participates in  $r$ .
- The *degree* of  $R$  is the number of participating entity types. For example,
  - ENROLMENT could be a ternary (degree 3) relationship between RESEARCHER, STUDENT and COURSE.
- We can illustrate this using an occurrence diagram:





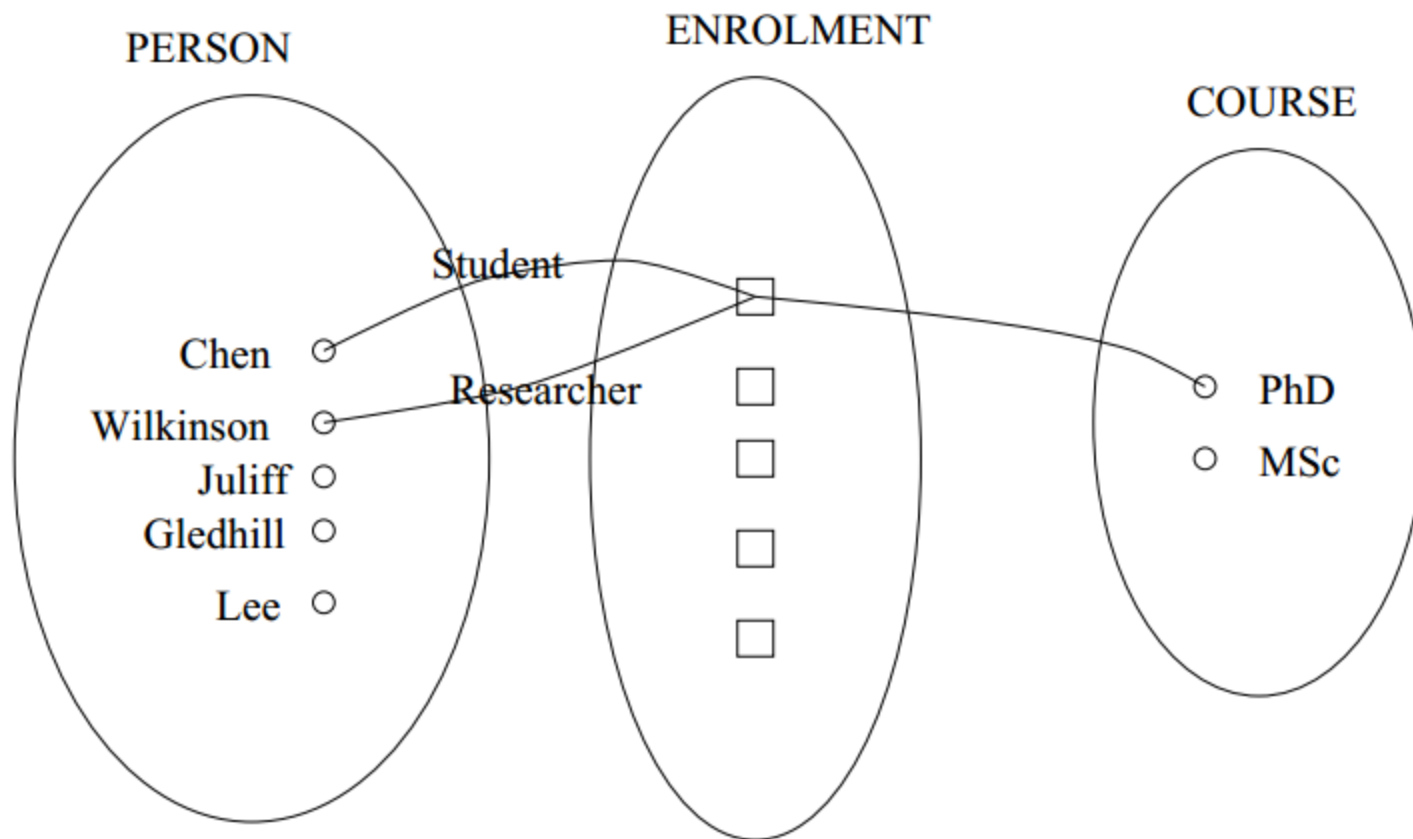
## 1.2.2 Relationships<sub>(cont)</sub>

- Entities and their relationships can also be represented using Entity-Relationship diagrams:



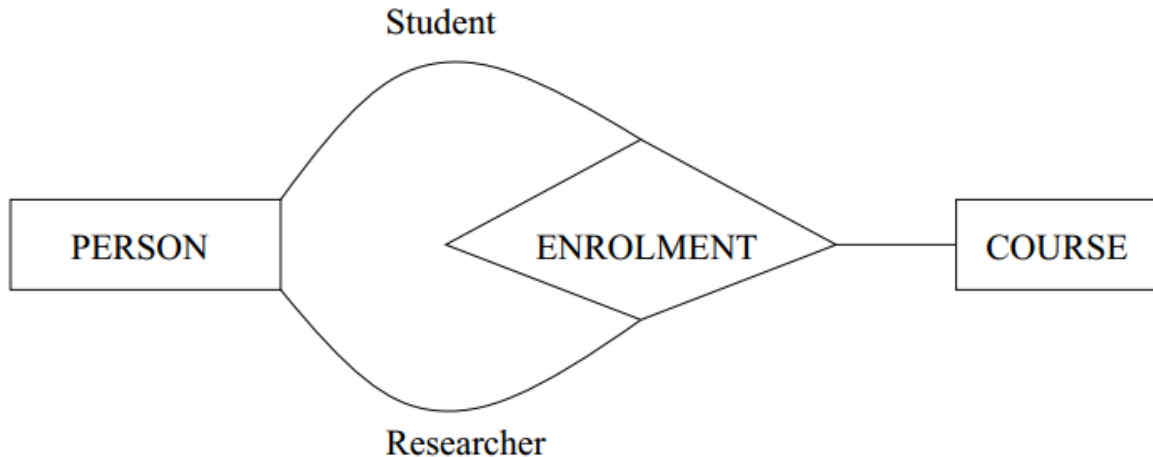
## 1.2.2 Relationships<sub>(cont)</sub>

- Each entity type that participates in a relationship plays a particular *role* in the relationship.
- An entity type can play
  - different roles in different relationships, or
  - more than one role in a relationship.
- A role name can be used to distinguish these.
- For example, ENROLMENT could be a relationship between PERSON(as researcher), PERSON(as student) and COURSE as in the diagram below:



## 1.2.2 Relationships<sub>(cont)</sub>

- Or, using an ERD:



- This is called a recursive relationship.

## 1.2.3 Weak entity types

- Some entity types do not have a key of their own.
- Such entity types are called weak entity types.
- Entities of a weak entity type can be identified by a partial key and by being related to another entity type - *owner*.
- The relationship type between a weak entity type to its owner is the *identifying relationship* of the weak entity type.

## 1.2.3 Weak entity types<sub>(cont)</sub>

- For example, a TAX PAYER entity may be related to several DEPENDENT, identified by their names.
- In this example, DEPENDENT is called a weak entity, {Name} is a partial key for it. The identifying relationship between DEPENDENT and TAX PAYER is IS DEPENDENT OF. TAX PAYER is said to *own* DEPENDENT.

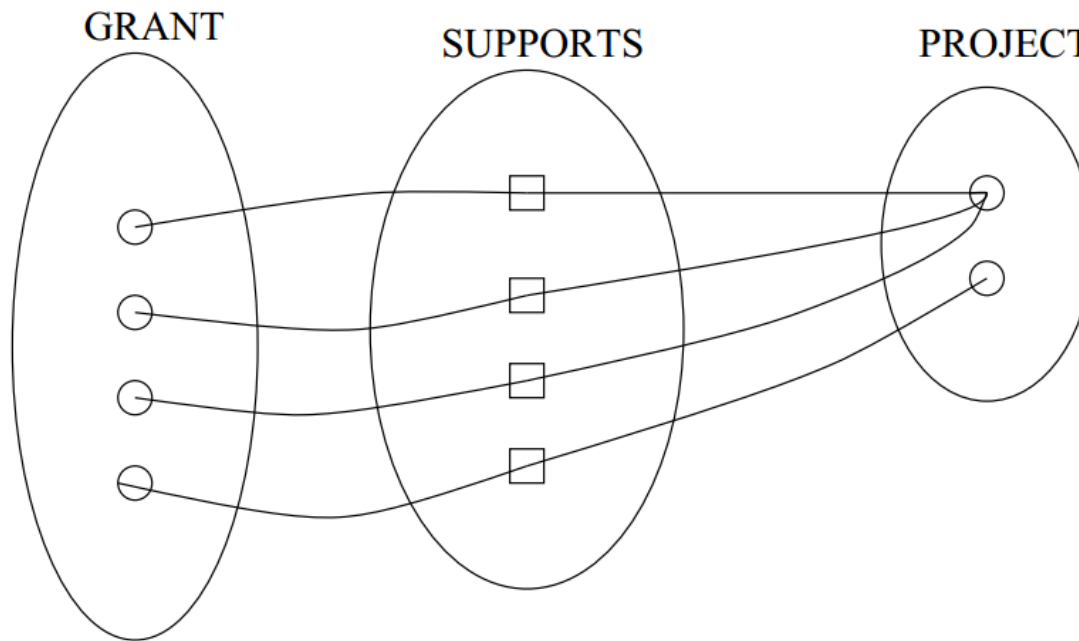
## 1.2.4 Constraints on relationship types

- Relationship types usually have certain constraints that limit the possible combinations of entities participating in relationship instances.
- They should reflect the correct factors
- *Cardinality ratio constraint*: specifies the number of relationship instances an entity can participate in.
- Example: A research grant supports only one research project, but a research project may be supported by many grants. PROJECT:GRANT is a 1 : N relationship.



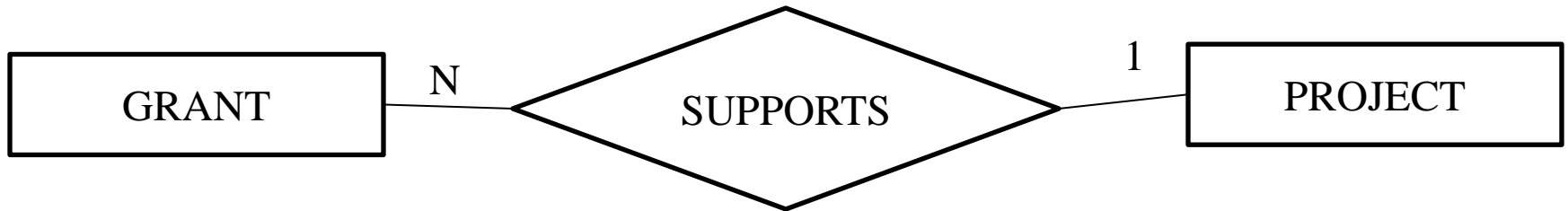
## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- This is illustrated in the occurrence diagram below:



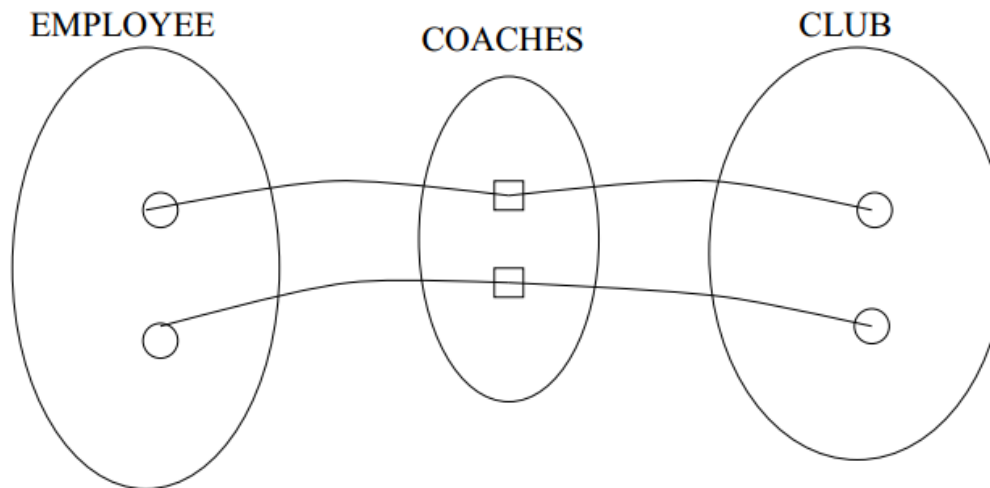
## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- We can also show this in an ERD:



## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- Example: Consider a database of AFL (here substitute your favourite team sport) statistics. The relationship of head coaches to clubs is an example of a 1 : 1 relationship.



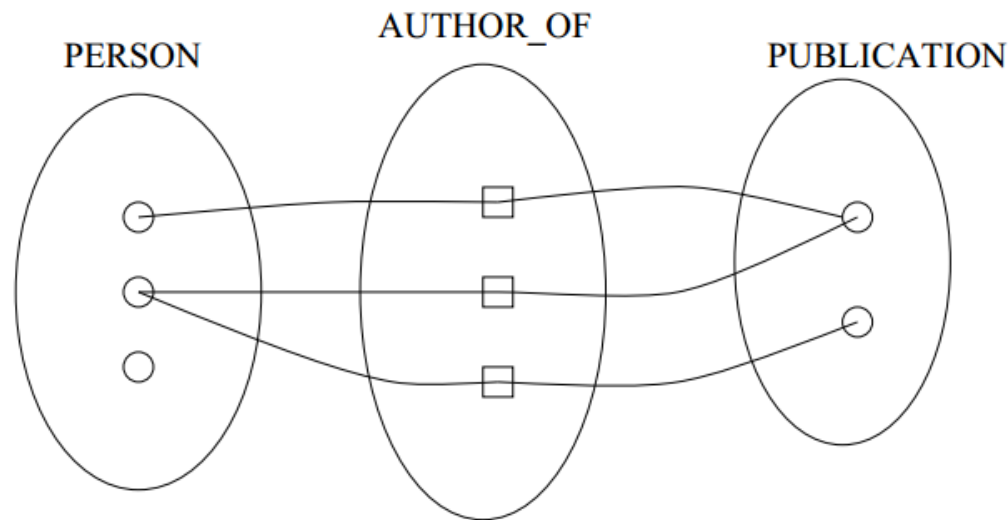
## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- With an ERD:



## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- Example: An example of an N : M relationship is authorship of publications:



## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- The equivalent ERD:



## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- Another kind of constraint that can be represented using the ER model is a
  - *Participation constraint*: participation of an entity in a relationship can be:
    - *total*: every entity must participate e.g. every publication has an author.
    - *partial*: not necessarily total. e.g. not every person has publications.

## 1.2.4 Constraints on relationship types<sub>(cont)</sub>

- This can be shown with an ERD like the one below:

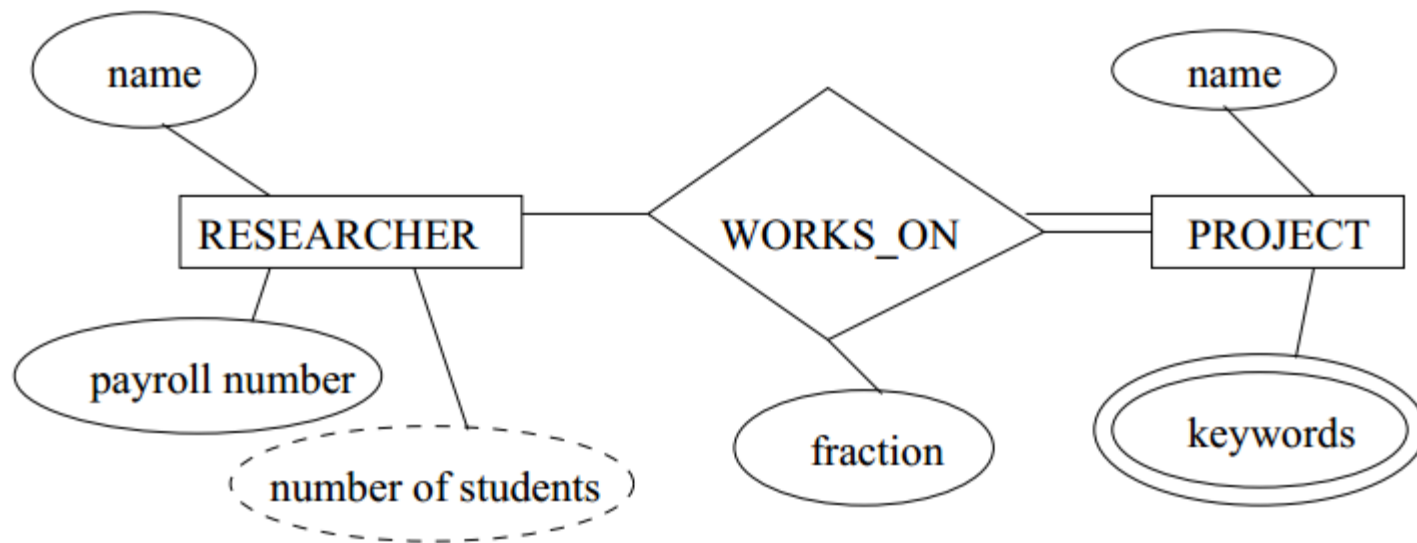




## 1.2.5 Attributes of relationship types

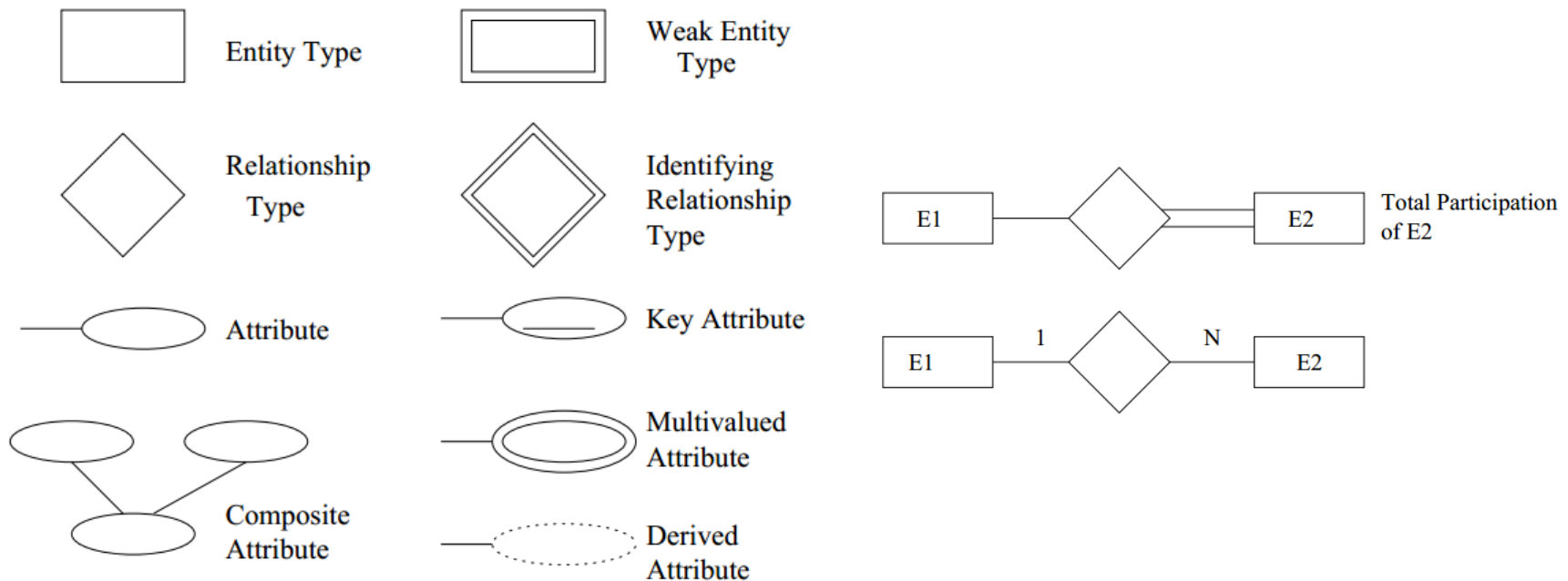
- Relationship types can have attributes – for example,
  - a researcher may work on several projects. The fraction of her time devoted to a particular project could be an attribute of the WORKS ON relationship type.
- This can be shown in an ERD as below:

## 1.2.5 Attributes of relationship types<sub>(cont)</sub>



## 1.2.5 Attributes of relationship types<sub>(cont)</sub>

- The notation used for ERDs is summarised in Elmasre/Navathe Figure 3.15.



## 1.2.6 Enhanced ER (EER) model

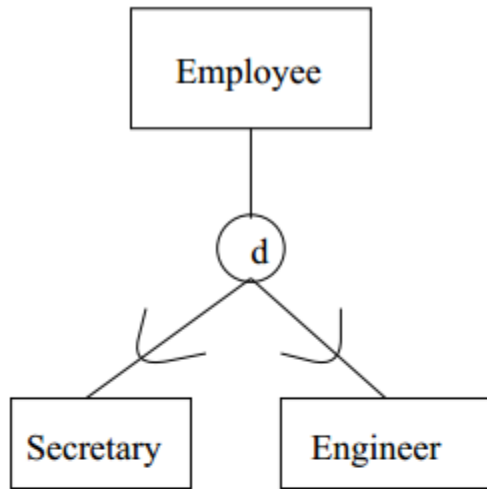
- Designers must use additionally modelling concepts to
  - represent the requirements from applications as accurately and explicitly as possible.

## 1.2.6 Enhanced ER (EER) model<sub>(cont)</sub>

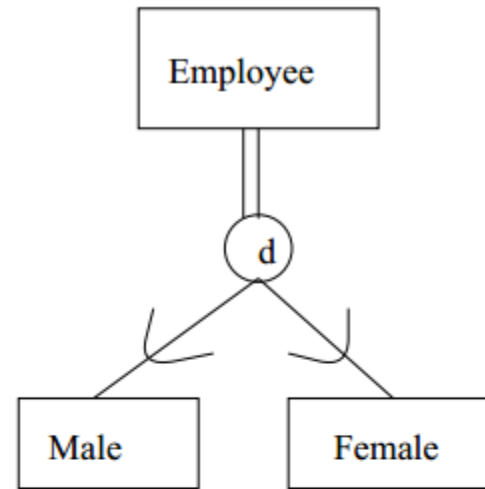
- There are many extensions to the ER model. We will look at one:
  - *Specialisation*: the process of defining a set of subclasses of an entity type; this entity type is called the superclass of the specialization.
  - *Generalisation*: a reverse process of specialisation.
- A subclass inherits all the attributes of the superclasses.

## 1.2.6 Enhanced ER (EER) model<sub>(cont)</sub>

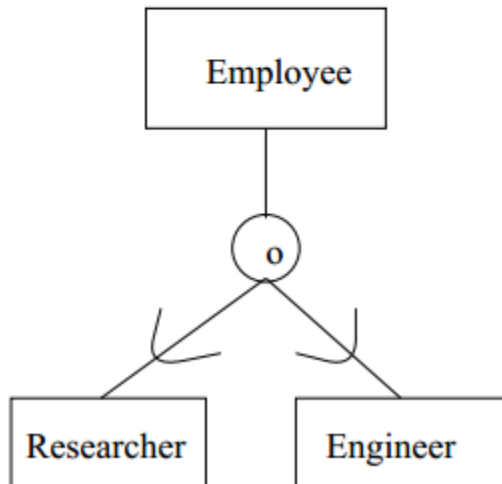
- A specialisation involves the following aspects:
  - Define a set of subclasses of an entity type.
  - Associate additional specific attributes with each subclass.
  - Establish additional specific relationship types between each subclass and other entity types, or other subclasses.
- A subclass may have multiple superclasses.
- A specialisation:
  - may be either total or partial; and
  - may be either disjoint or overlapping.



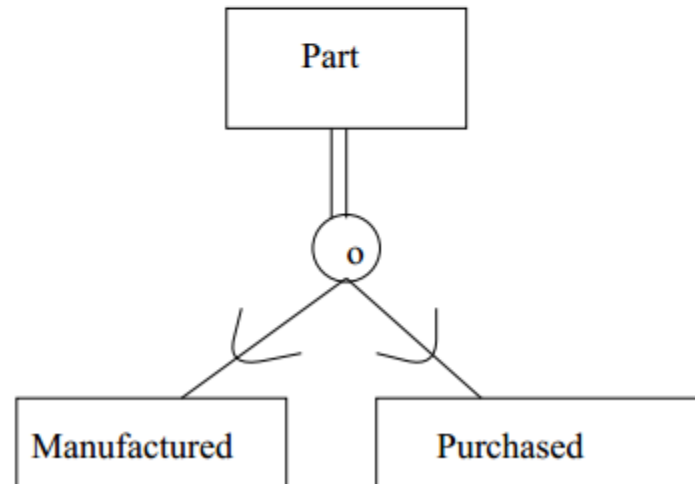
Partial disjoint



Total disjoint



Partial Overlapping



Total Overlapping

## 1.2.7 Design Principles

- Faithfulness: reflect reality.
- Avoid redundancy.
- Picking the right kind of element.