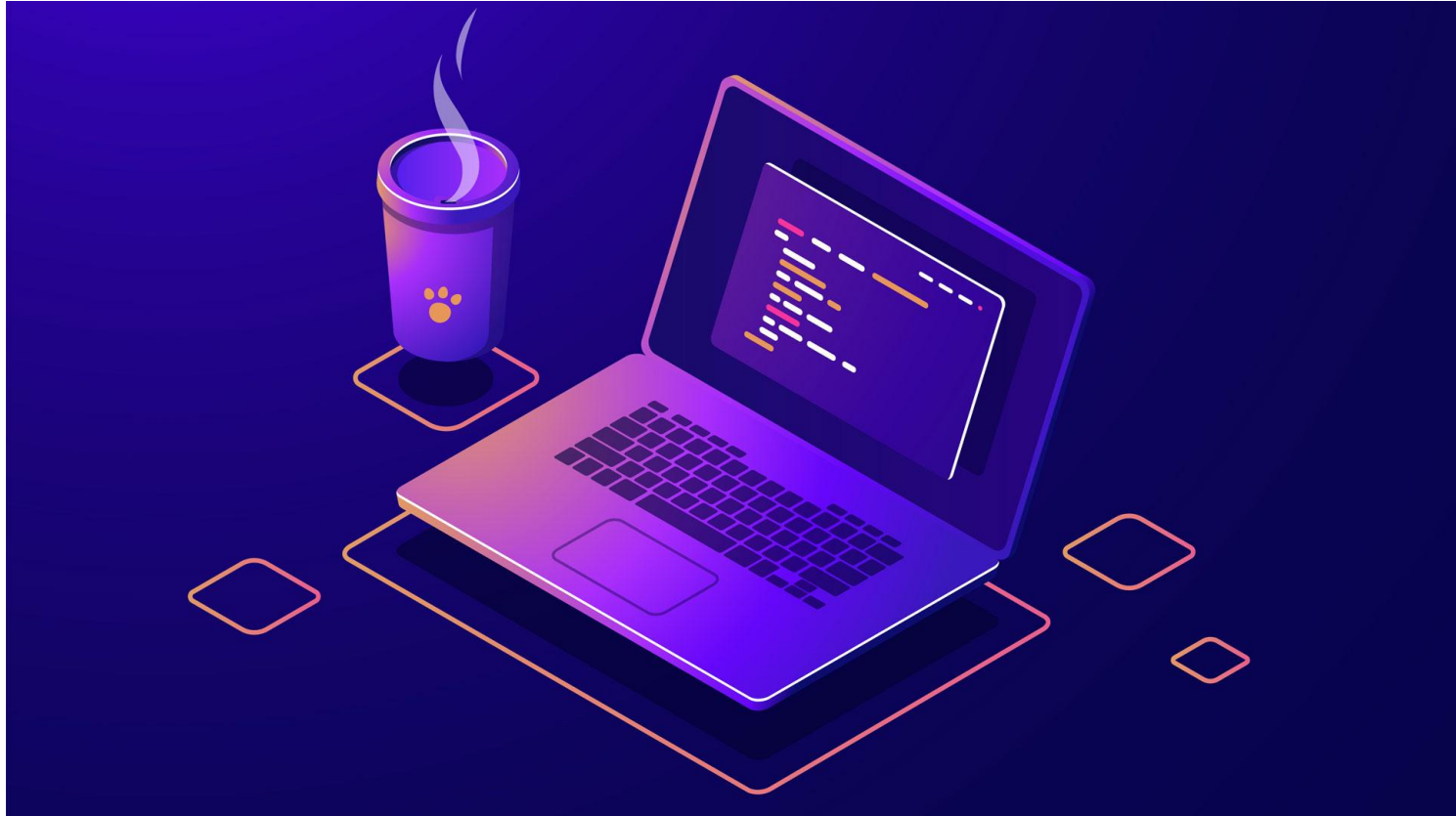


Task Definition



Task Definition



- ▶ Can the computation be divided in parts?¹
 - ▶ Task decomposition: based on the processing to do (e.g. functions, loop iterations)
 - ▶ Data decomposition: based on the data to be processed (e.g. elements of a vector, rows of a matrix) (implies task decomposition)
 - ▶ There may be (data or control) dependencies between tasks
- ▶ Metrics to understand how our task/data decomposition can potentially behave
- ▶ Factors: granularity and overheads

Task Definition



- ▶ TDG: directed acyclic graph to represent tasks and dependencies between them
- ▶ Metrics:
 - ▶ $T_1 = \sum_{i=1}^{nodes} (work_node_i)$
 - ▶ $T_\infty = \sum_{i \in criticalpath} (work_node_i)$, assuming sufficient (infinite) resources
 - ▶ $Parallelism = T_1 / T_\infty$
 - ▶ P_{min} is the minimum number of processors necessary to achieve $Parallelism$
- ▶ Task granularity vs. number of tasks

Example 1: Vector Sum

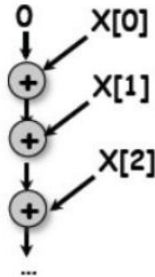


Compute the sum of elements $X[0] \dots X[n-1]$ of a vector X

```
sum = 0; for ( i=0 ; i< n ; i++ ) sum += X[i];
```

Task definition: each iteration of the i loop is a task.

► **TDG** (with input data):



► **Metrics:**

$$T_1 \propto n$$

$$T_\infty \propto n$$

$$Parallelism = 1$$

How can we design an algorithm which leads to a TDG with more parallelism?

Example 1: Vector Sum



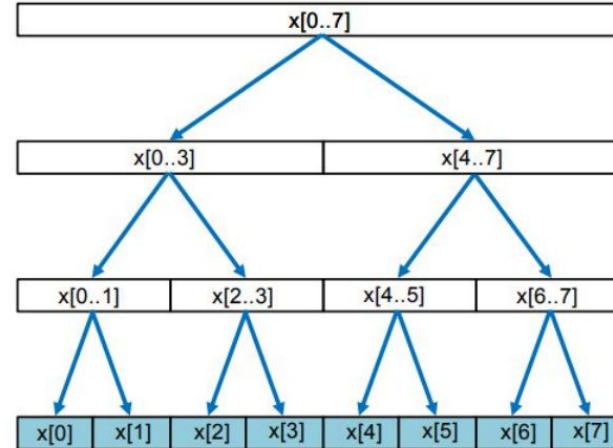
Writing a **recursive version** of the sequential program to compute the sum of elements $X[0] \dots X[n-1]$ of a vector X , following a *divide-and-conquer* strategy:

```
int recursive_sum(int *X, int n) {
    int ndiv2 = n/2;
    int sum=0;

    if (n==1) return X[0];

    sum1 = recursive_sum(X, ndiv2);
    sum2 = recursive_sum(X+ndiv2, n-ndiv2);
    return sum1+sum2;
}

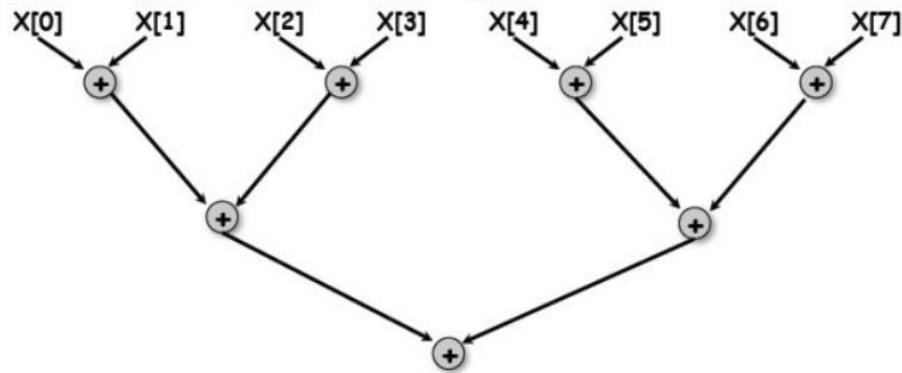
void main() {
    int sum, X[N];
    ...
    sum = recursive_sum(X,N);
    ...
}
```



Example 1: Vector Sum



- ▶ **Task definition:** each invocation to `recursive_sum`
- ▶ **TDG** (with input data):



- ▶ **Metrics:**
 $T_1 \propto n$; $T_\infty \propto \log_2(n)$; $Parallelism \propto (n \div \log_2(n))$
- ▶ Same problem can be expressed with different algorithms/implementations leading to different metrics

Instructor Social Media

Youtube: Lucas Science



Instagram: lucaasbazilio



Twitter: lucasebazilio

