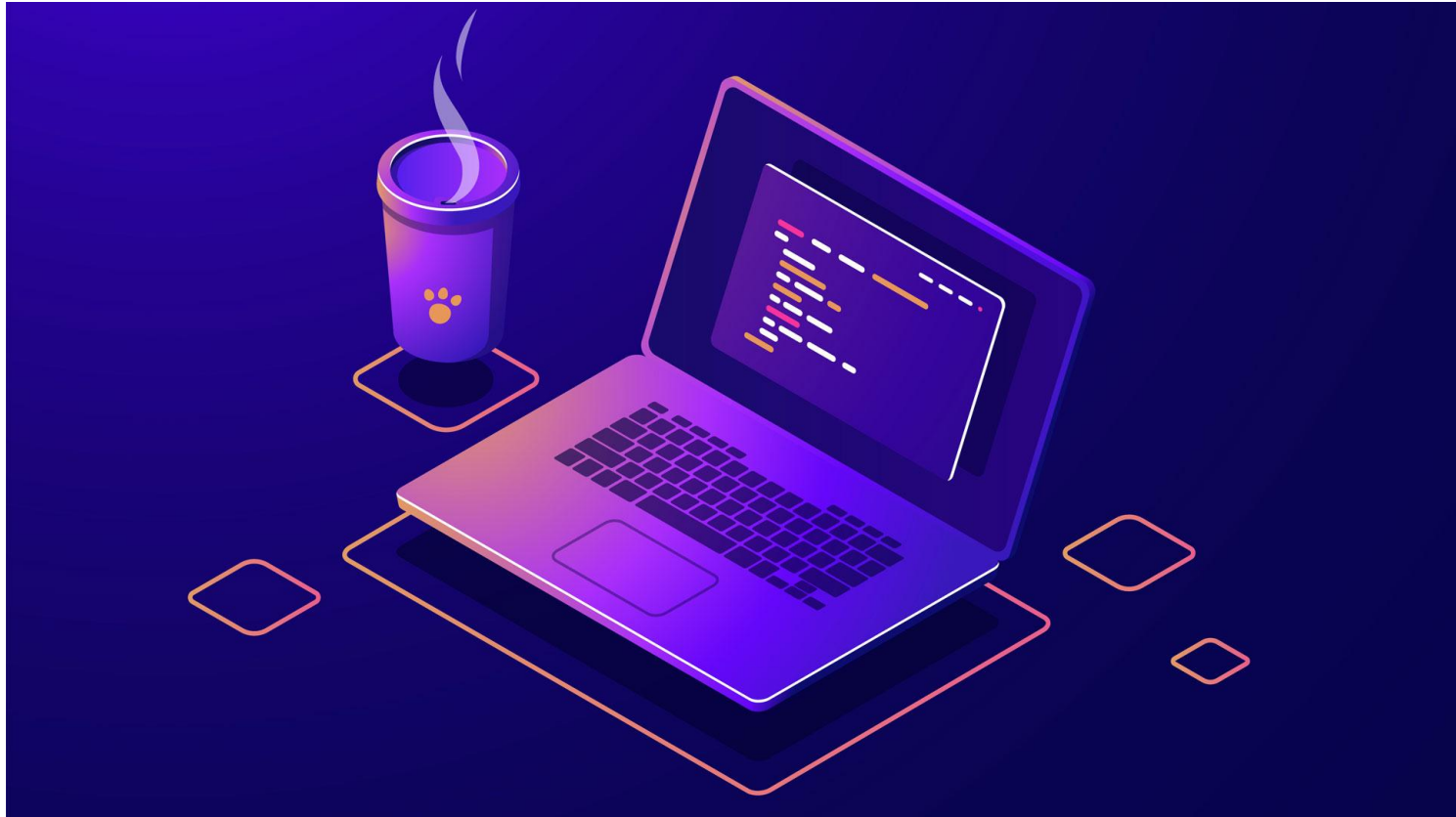
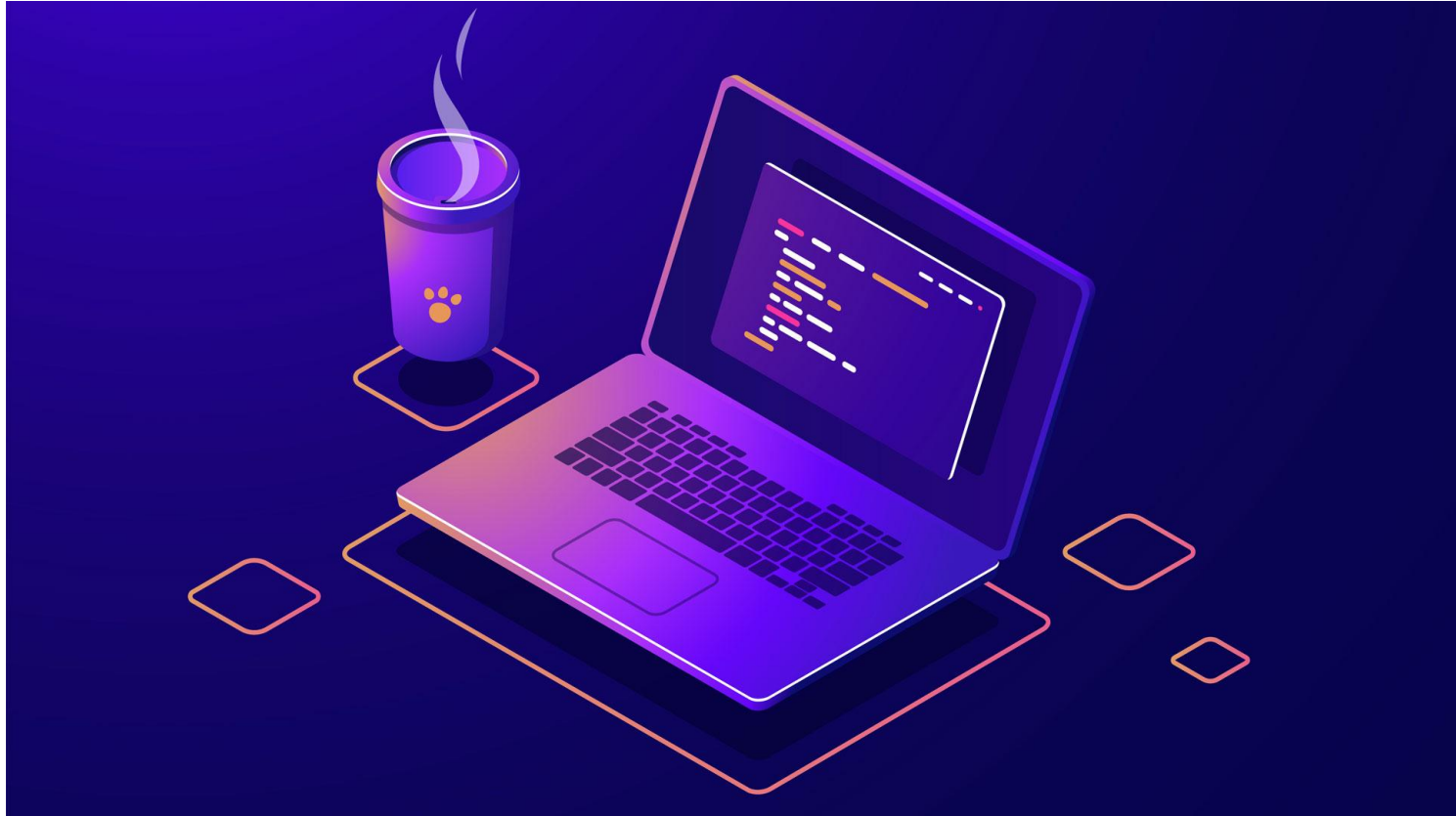


# OpenMP Problems



## Problem 2 - Point C



## Problem 2



The following sequential code in C finds all positions in vector DBin in which a set of keys (contained in vector keys) appear. Positions where keys appear are stored in a new vector DBout (the order in DBout of the positions found is irrelevant).

```
int main() {
    double keys[nkeys], DBin[DBsize], DBout[nkeys][DBsize];
    unsigned int i, k, counter[nkeys];

    getkeys(keys, nkeys);           // get keys
    init_DBin(DBin, DBsize);        // initialize elements in DBin
    clear_DBout(DBout, nkeys, DBsize); // initialize elements in DBout
    clear_counter(counter, nkeys);   // initialize counter to zero

    for (i = 0; i < DBsize; i++)
        for (k = 0; k < nkeys; k++)
            if (DBin[i] == keys[k]) DBout[k][counter[k]++] = i;
}
```

# Problem 2



```
#define DBsize 1048576
#define nkeys 16 // the number of processors can be larger than the number of keys

int main() {
    double keys[nkeys], DBin[DBsize], DBout[nkeys][DBsize];
    unsigned int i, k, counter[nkeys];

    getkeys(keys, nkeys);           // get keys
    init_DBin(DBin, DBsize);        // initialize elements in DBin
    clear_DBout(DBout, nkeys, DBsize); // initialize elements in DBout
    clear_counter(counter, nkeys);   // initialize counter to zero

    for (i = 0; i < DBsize; i++)
        for (k = 0; k < nkeys; k++)
            if (DBin[i] == keys[k]) DBout[k][counter[k]++] = i;
}
```

# Problem 2



- (a) Write a first *OpenMP* parallelisation that implements an **iterative task decomposition strategy** of the outermost loop `i`, making use of the `taskloop` directive, in which you minimise the serialisation introduced by the synchronisation that you may introduce. **Note:** you are not allowed to change the structure of the two loops.
- (b) Write a second *OpenMP* parallelisation that also implements an **iterative task decomposition strategy**, but this time applied to the innermost loop `k`, again making use of the `taskloop` directive, in which you maximise the parallelism that can be exploited. **Notes:** 1) `taskloop` has an implicit `taskgroup` synchronisation that you can omit with the `nogroup` clause; 2) observe that the number of keys is not large when compared to the possible number of processors to use; and 3) you are not allowed to change the structure of the two loops.
- (c) Finally, write a third *OpenMP* parallelisation that implements a **task-based recursive divide-and-conquer decomposition strategy**, with the following requirements: 1) the recursion splits the input vector `DBin` in two almost identical halves, with a base case that corresponds to checking a single element of `DBin`; 2) uses a **cut-off strategy based on the size of the input vector**, so that tasks are only generated while that size is larger than `CUT_SIZE`; 3) only uses *OpenMP* pragmas and clauses for the implementation of the cut-off strategy; and 4) you have to use the synchronisation mechanism, if needed, that maximises the parallelism in the program.

# Mutual Exclusion



Mutual exclusion: mechanism to ensure that only one task at a time executes the code within a region.



# Mutual Exclusion



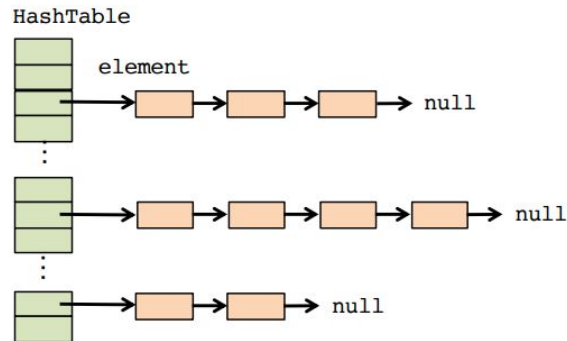
Associate a lock variable with each slot in the hash table, protecting the chain of elements in an slot

```
omp_lock_t hash_lock[SIZE_HASH];

#pragma omp parallel
#pragma omp single
{
    for (i = 0; i < SIZE_HASH; i++) omp_init_lock(&hash_lock[i]);

    #pragma omp taskloop
    for (i = 0; i < SIZE_TABLE; i++) {
        int index = hash_function (dataTable[i], SIZE_HASH);
        omp_set_lock (&hash_lock[index]);
        insert_element (dataTable[i], index, HashTable);
        omp_unset_lock (&hash_lock[index]);
    }

    for (i = 0; i < SIZE_HASH; i++) omp_destroy_lock(&hash_lock[i]);
}
```



Threads may be inserting elements into the hash table in parallel, as long as these elements hash to different slots

# Recursive Task Decomposition



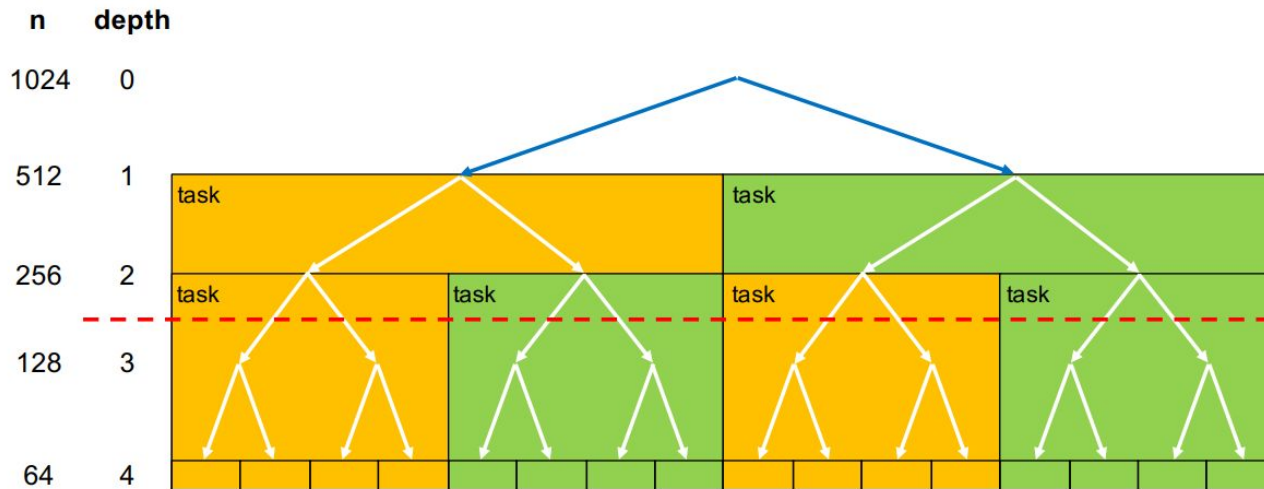
In recursive task decomposition strategies one can control task granularity by controlling recursion levels where tasks are generated (cut-off control).



# Recursive Task Decomposition

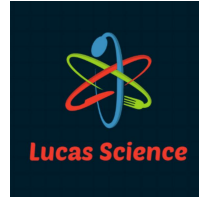


Tree strategy with **depth recursion control**



# Instructor Social Media

**Youtube: Lucas Science**



**Instagram: lucaasbazilio**



**Twitter: lucasebazilio**

