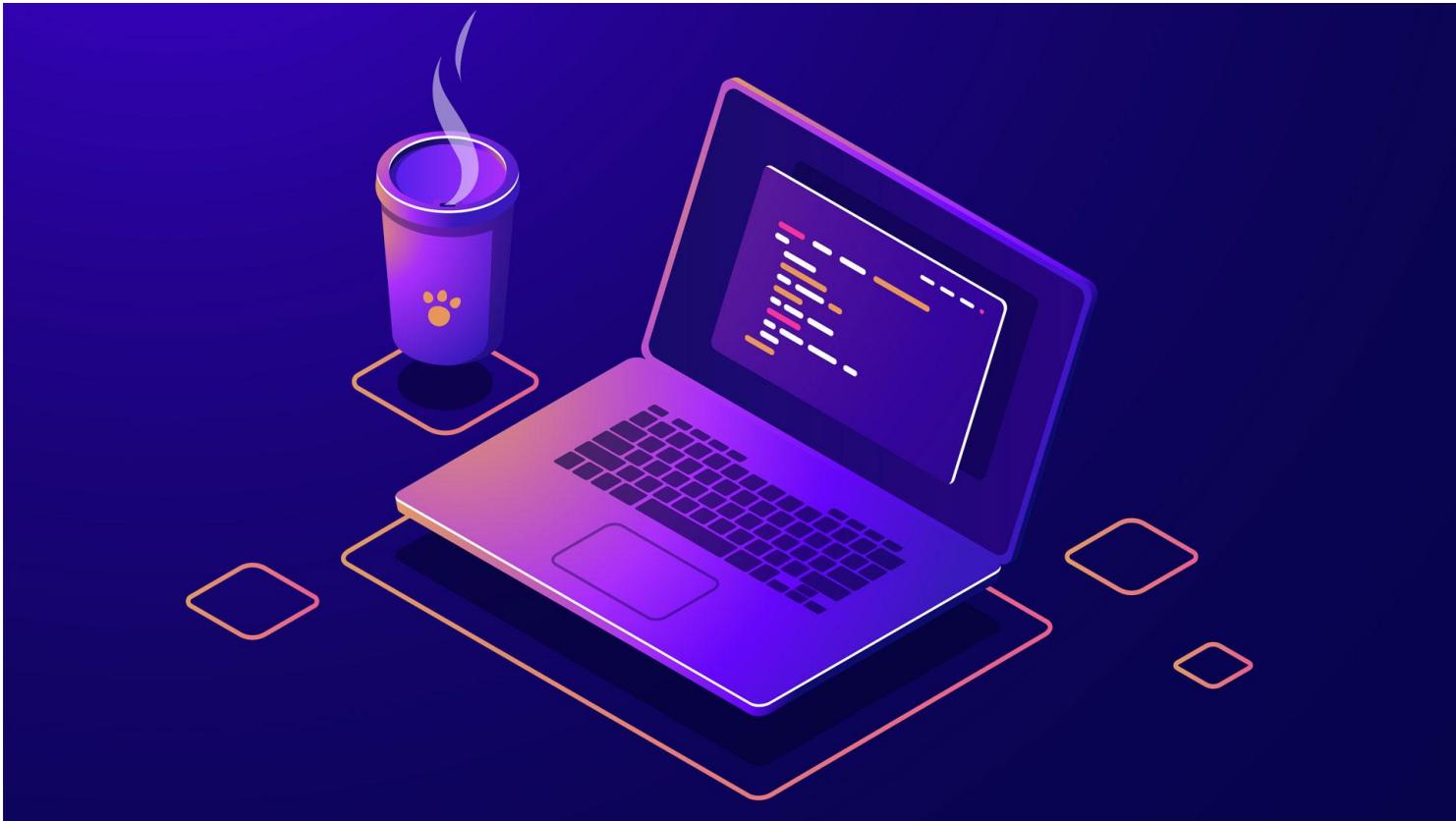


# Task Definition



# Task Definition



- ▶ Can the computation be divided in parts?<sup>1</sup>
  - ▶ Task decomposition: based on the processing to do (e.g. functions, loop iterations)
  - ▶ Data decomposition: based on the data to be processed (e.g. elements of a vector, rows of a matrix) (implies task decomposition)
  - ▶ There may be (data or control) dependencies between tasks
- ▶ Metrics to understand how our task/data decomposition can potentially behave
- ▶ Factors: granularity and overheads

# Task Definition



- ▶ TDG: directed acyclic graph to represent tasks and dependencies between them
- ▶ Metrics:
  - ▶  $T_1 = \sum_{i=1}^{nodes} (work\_node_i)$
  - ▶  $T_\infty = \sum_{i \in criticalpath} (work\_node_i)$ , assuming sufficient (infinite) resources
  - ▶  $Parallelism = T_1/T_\infty$
  - ▶  $P_{min}$  is the minimum number of processors necessary to achieve  $Parallelism$
- ▶ Task granularity vs. number of tasks



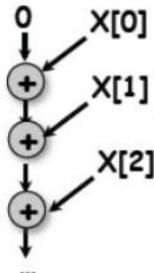
# Example 1: Vector Sum

Compute the sum of elements  $X[0] \dots X[n-1]$  of a vector  $X$

```
sum = 0; for ( i=0 ; i< n ; i++ ) sum += X[i];
```

**Task definition:** each iteration of the  $i$  loop is a task.

- ▶ **TDG** (with input data):



- ▶ **Metrics:**

$$T_1 \propto n$$

$$T_\infty \propto n$$

$$\text{Parallelism} = 1$$

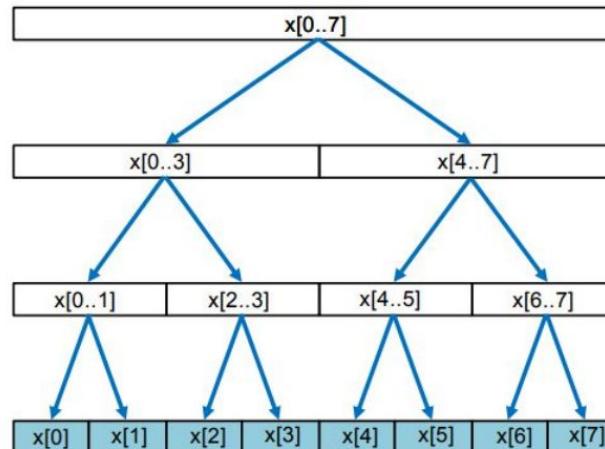
How can we design an algorithm which leads to a TDG with more parallelism?



# Example 1: Vector Sum

Writing a **recursive version** of the sequential program to compute the sum of elements  $X[0] \dots X[n-1]$  of a vector  $X$ , following a *divide-and-conquer* strategy:

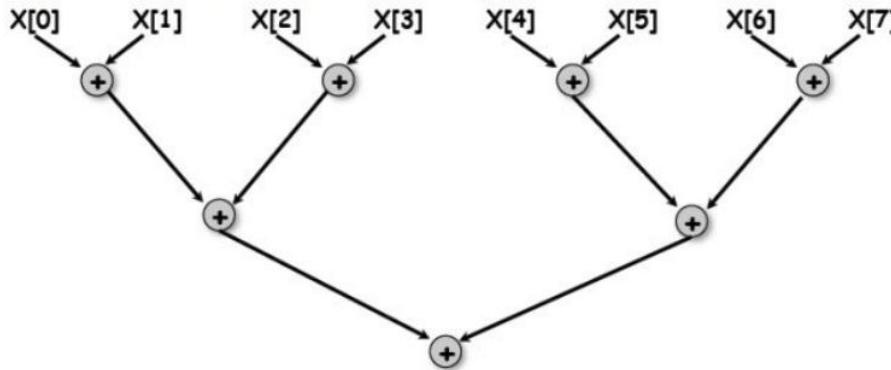
```
int recursive_sum(int *X, int n) {  
    int ndiv2 = n/2;  
    int sum=0;  
  
    if (n==1) return X[0];  
  
    sum1 = recursive_sum(X, ndiv2);  
    sum2 = recursive_sum(X+ndiv2, n-ndiv2);  
    return sum1+sum2;  
}  
  
void main() {  
    int sum, X[N];  
    ...  
    sum = recursive_sum(X,N);  
    ...  
}
```





# Example 1: Vector Sum

- ▶ **Task definition:** each invocation to recursive\_sum
- ▶ **TDG (with input data):**



- ▶ **Metrics:**  
 $T_1 \propto n$ ;  $T_\infty \propto \log_2(n)$ ; *Parallelism*  $\propto (n \div \log_2(n))$
- ▶ Same problem can be expressed with different algorithms/implementations leading to different metrics

# Instructor Social Media

**Youtube: Lucas Science**



**Instagram: lucaasbazilio**



**Twitter: lucasebazilio**

