



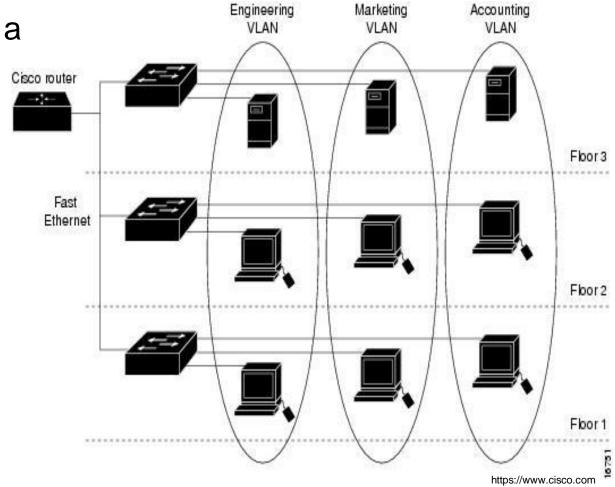
Teaser of Video at the end. Removed for Download.



VLAN

• set up multiple isolated virtual networks on a single physical infrastructure

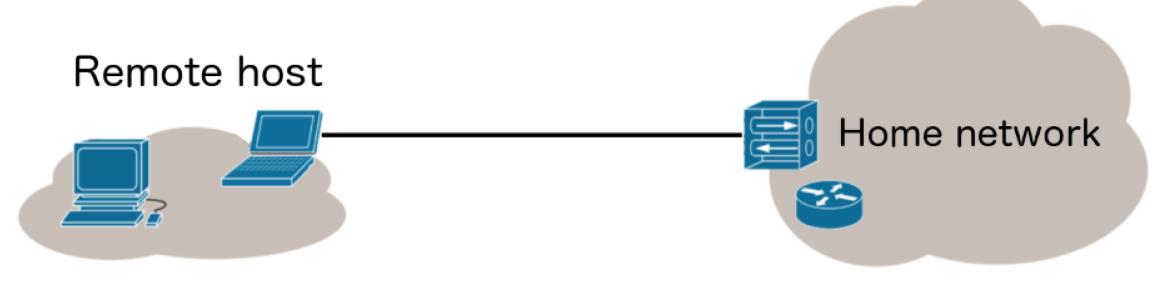
Resides on L2



1.1: Avoid opening ports

You are hosting some service at **home**, how can you access it without exposing it to the internet?

VPN Client to home network



1.2: Certificate Authority

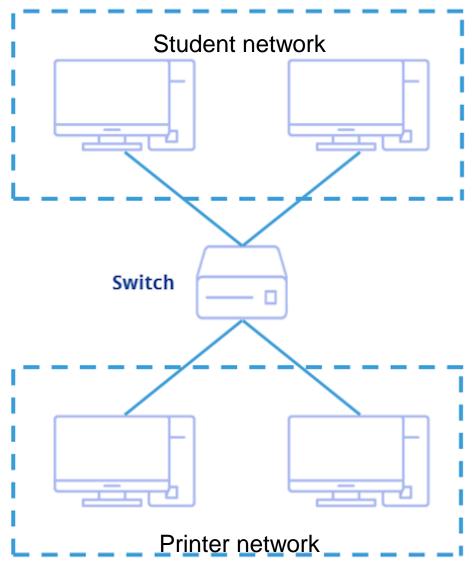
Accessing a **CA admin** panel from work

VPN provides client authentication

VPN as access control



1.3: University



Need to run **separate networks** for student access, staff and printers

- Classic case for VLAN configuration
- Guarantee isolation
- Can share infrastructure, no need to move printers close together

1.4: Print from home

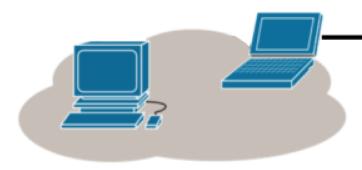
You are a student, want to print from home

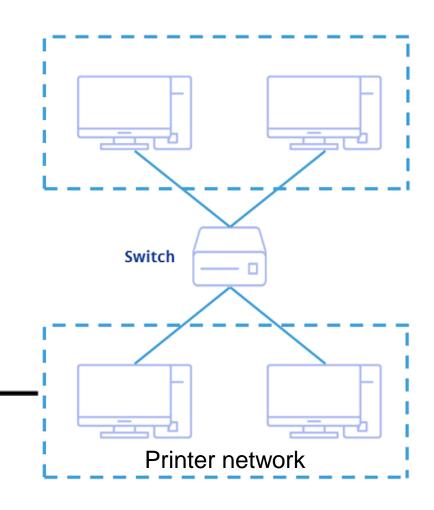
VPN

with endpoint in

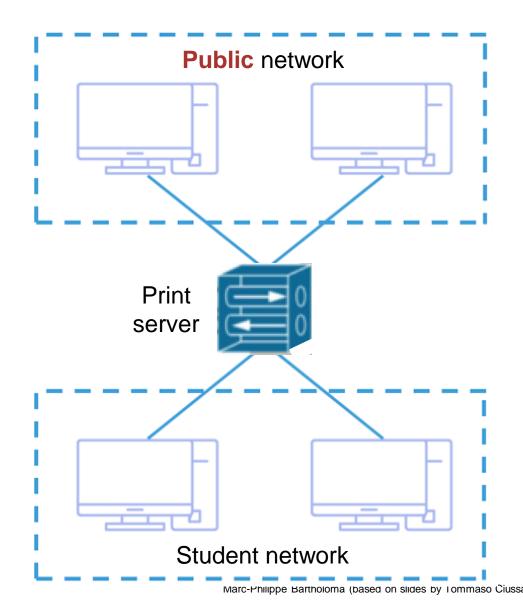
printer network

Remote host





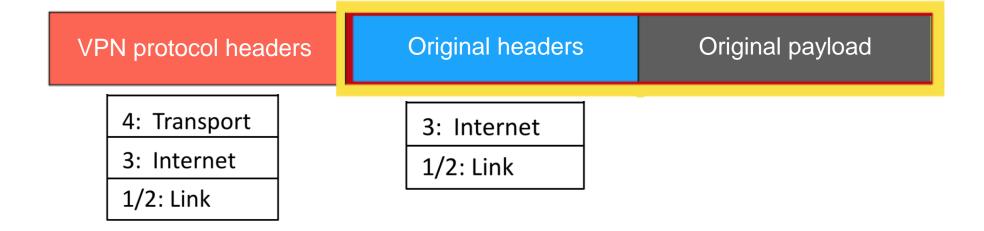
1.5: Print via email



How to enable print via email functionality?

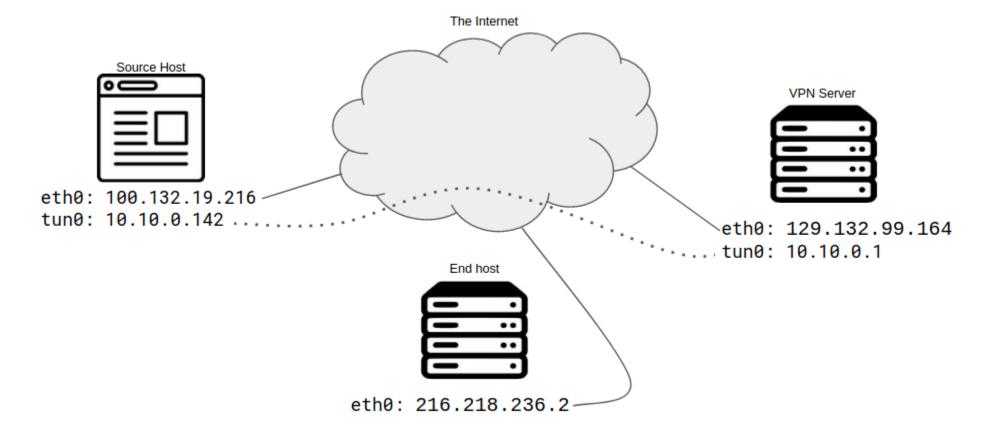
- Print + email server connected to both
 VLANs
- Receives emails from public network
- Sends them to the printer network
- Isolation ensured via software

2: Encapsulation of VPN packets



3: VPN Routing

Fix routing table

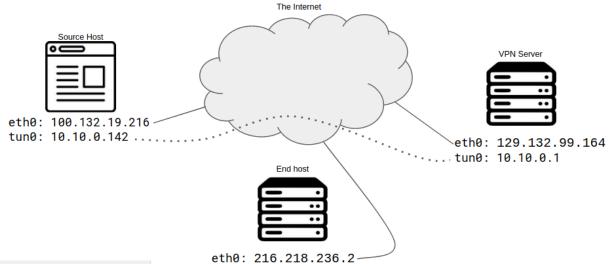


Kernel IP	routing	table
-----------	---------	-------

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
0.0.0.0	100.132.19.216	0.0.0.0	UG	0	0	0	eth0
100.132.19.0	0.0.0.0	255.255.255.0	U	0	0	0	eth0
10.10.0.0	0.0.0.0	255.255.0.0	U	0	0	0	tun0

3: VPN Routing

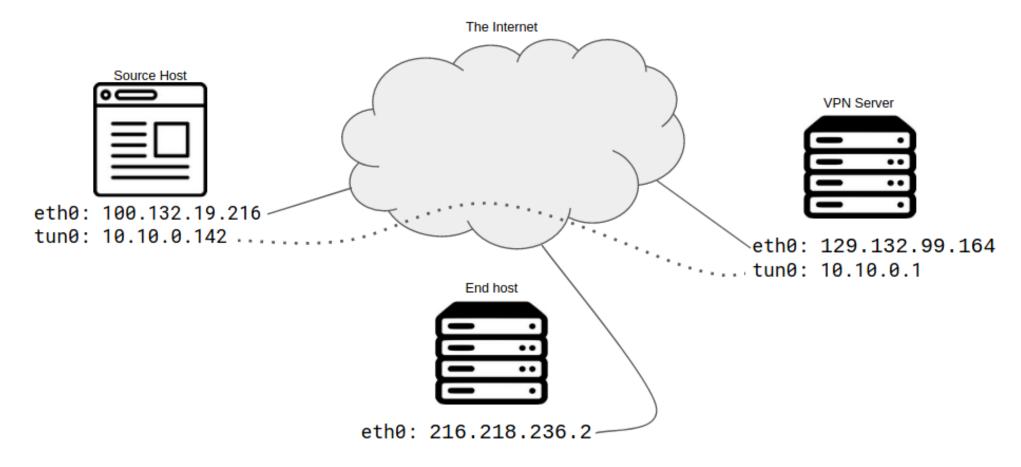
Track source and destination



packet	source IP	destination IP
original packet at source host (client)	not set	216.218.236.2
packet after routing decision		
headers of the outer protocol		
packet after decapsulation at VPN server		
packet after NATing at VPN server		

3: VPN Routing

Return route



4.1 Client certificate authentication

IPSec

- Supposed to be a private network
- VPN access must be restricted by definition
- Identifying both clients and servers is important
- Client certificates are exchanged

Usually no client certificate

TLS

- Websites are publicly accessible
- Client identity not important*
- Server identity critical
- MitM risk otherwise!

^{*}or achieved with different mechanisms

4.2: Client Cert leaks identity

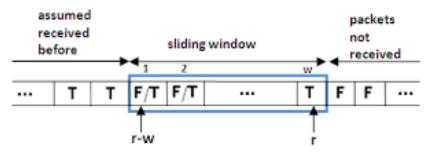
IPSec

- Encrypt using Anonymous Diffie Hellman!
- Passive MitM can't see certs
- Active MitM can still hijack the session (Remember: ANON DH)

TLS 1.3

- Also protects client identity
- (TLS 1.2 didn't)
- Error in exercise sheet: just replaced TLS 1.2 with TLS 1.3

4.3: Sequence numbers in IPSec (vs TLS)



T: true, the message has been received before

F: false, the message has not yet been received

Figure 1: The anti-replay window

- Sequence numbers used to prevent replay attacks
- Every party has a sliding window to avoid repeated packets

- IP is best effort, packets could be lost and order disrupted
- Need to transmit the numbers in IPSec header

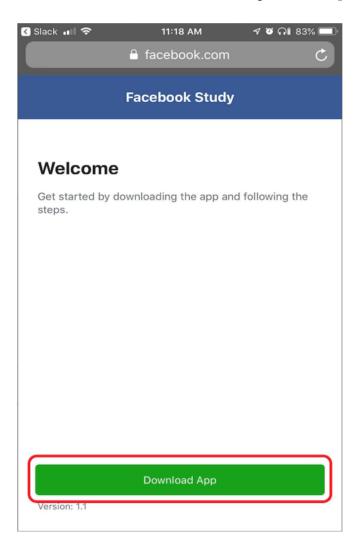
Facebook Research app

- Dear Facebook user, why don't you participate in a study?
- Join Facebook research program and get
 20\$/month for your collaboration!



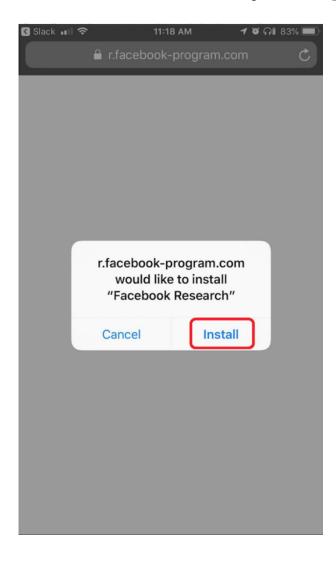
What's the catch?

What is actually happening



Phase 1:
Unsuspecting users
download the app

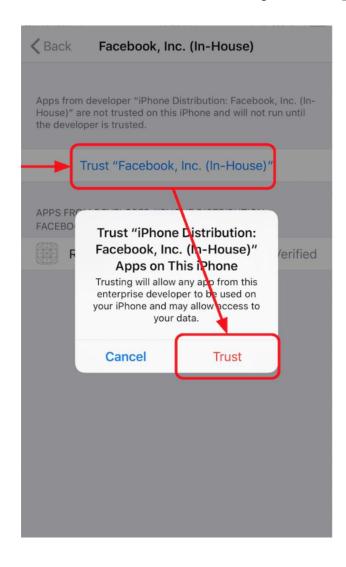
What is actually happening



Phase 2:

The app installs and activates a VPN pointing to Facebook

What is actually happening



Phase 3:

The app installs and activates a root certificate

Most users have no idea they just sold their privacy for 20\$

5: Facebook in the Middle

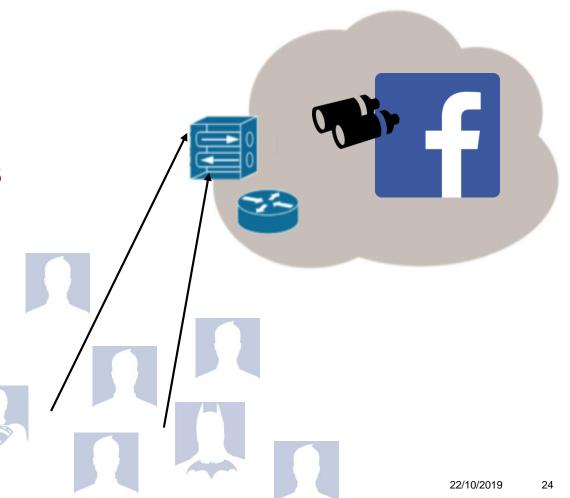
Your traffic is routed through Facebook

No effort to sniff your traffic

Facebook can create bogus certificates
 that your phone will trust!

 Can read your TLS sessions in plain text!





Wireguard basics

- Lightweight codebase, just 4K lines
- Limited set of cipher suites
- Every party has a static key pair
- Public key is not supposed to be available to everyone



Cryptokey routing

- Every interface has private key
 and a list of peers (public keys)
- List associates public keys with allowed IPs
- IPs used for routing
- IPs used for access control

```
[Interface]
PrivateKey = yAnz5TF+lXXJte14tji3zlMNq+hd2rYUIgJBgB3fBmk=
ListenPort = 51820
[Peer]
PublicKey = xTIBA5rboUvnH4htodjb6e697QjLERt1NAB4mZqp8Dg=
AllowedIPs = 10.192.122.3/32, 10.192.124.1/24
[Peer]
PublicKey = TrMvSoP4jYQlY6RIzBgbssQqY3vxI2Pi+y71l0WWXX0=
AllowedIPs = 10.192.122.4/32, 192.168.0.0/16
[Peer]
PublicKey = gN65BkIKy1eCE9pP1wdc8ROUtkHLF2PfAqYdyYBz6EA=
AllowedIPs = 10.10.10.230/32
```

6.1: PFS in WireGuard

- Tunnel_ephemeral keys derived from pubkeys and ephemeral keys
- Via DH and MAC computation

- PFS:
 - REKEY_AFTER_MESSAGES
 - REKEY_AFTER_TIME

PFS or weaker form?

Why is this different from Signal?

6.2: DoS attack mitigation

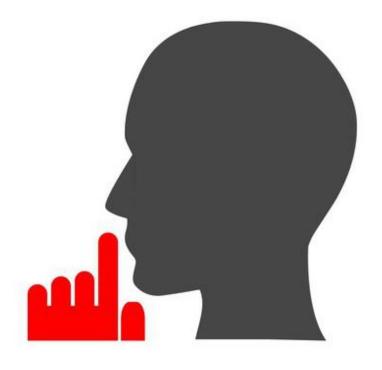
1. Silence is virtue

No reply in these situations:

- Sender doesn't know receiver public key
- Unauthorized originating IP

Attackers not able to flood, no state kept

Think about SYN flood...

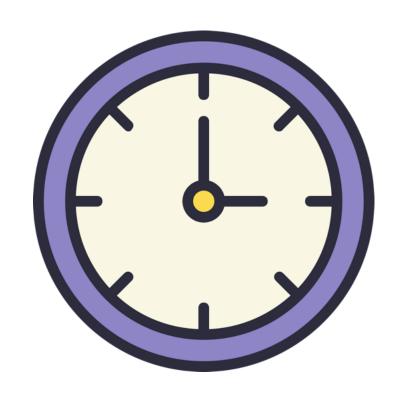


6.2: DoS attack mitigation

2. Timestamp

Sent with every packet:

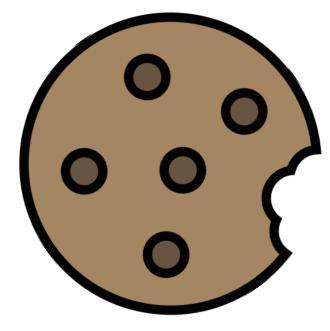
- Prevents replay attacks
- Same function as seqnum in IPSec



3. Cookie: server already under load, needs to avoid crypto

1. S replies with cookie:

- S has periodically changing random value R
- Cookie encrypted with client public key as symmetric key
- Content: MAC_R(IP of client C)

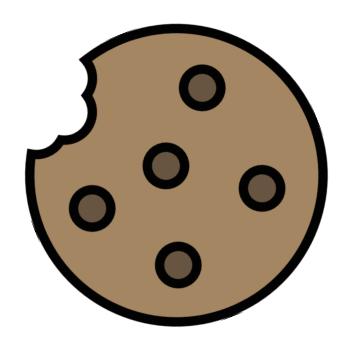


https://www.flaticon.com/authors/freepik

3. Cookie: server already under load, needs to avoid crypto

- 2. Client resends message, msg.mac2=MAC_{cookie}(msg)
 - Server checks a symm encryption and matches the IP
 - Proof of IP ownership

3. Server can apply rate limiting to IPs



Difference with IKEv2

Very similar cookie mechanism, but:

- IPSec maintains state
- Max num of open connections
- No msg.mac1 verification -> no silence
- Cookie transmitted in plaintext



5.4 VPNs and logging



Tom Scott, https://www.youtube.com/watch?v=WVDQEoe6ZWY

Your Questions

