

# Approximate Incentive Compatibility

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# Incentive compatibility (IC)

Fundamental concept in mechanism design

Buyers maximize their utilities by bidding truthfully



# Many real-world mechanisms are not incentive compatible

# Discriminatory auctions

Multi-unit variant of first-price auction

#### Not incentive compatible

Used to sell treasury bills since 1929

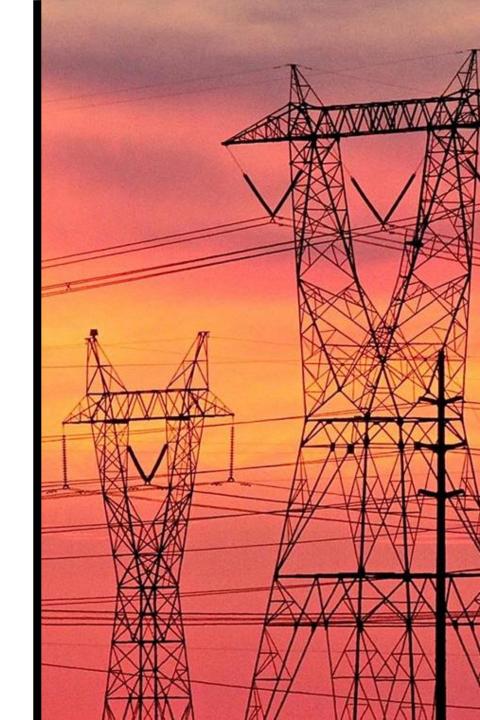


# Discriminatory auctions

Multi-unit variant of first-price auction

#### Not incentive compatible

Used to sell treasury bills since 1929 and electricity in the UK



#### GSP auction

Used for sponsored search

Not incentive compatible

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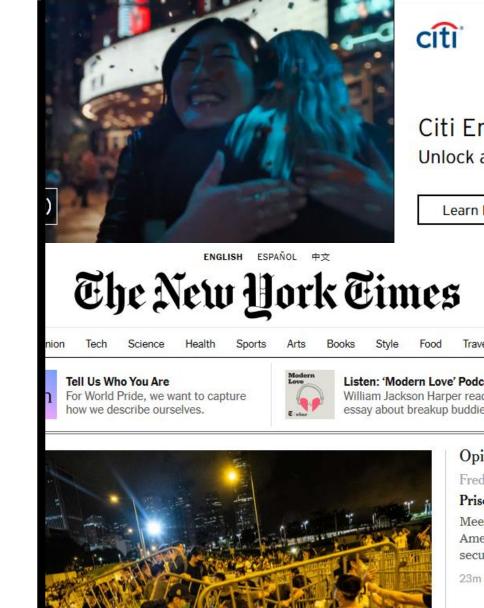


San Diego

# First-price auction

Ad exchanges transitioning to FP auction

Not incentive compatible



#### Combinatorial auctions

Nearly all fielded combinatorial auctions

(such as sourcing auctions)

aren't incentive compatible



# Why aren't real-world mechanisms incentive compatible?

# Why not IC?

**Expensive** to compute true values

Rules are easier to explain

Bids used to tune **future** parameters

Might leak **private** values

Agents not **risk** neutral



# Approximate incentive compatibility

Mechanism is  $\gamma$ -IC when for each bidder i:

If everyone except bidder i is truthful,

she can only increase utility by  $\gamma$  if she bids strategically

[Kothari, Parkes, and Suri, EC'03; Archer, Papadimitriou, Talwar, and Tardos, Internet Mathematics '04; Conitzer and Sandholm, IJCAI'07; Dekel, Fischer, and Procaccia, JCSS'10; Lubin and Parkes, Current Science '12; Mennle and Seuken, EC'14; Dütting, Fischer, Jirapinyo, Lai, Lubin, and Parkes TEAC'15; Azevedo and Budish, Review of Economic Studies '18; Feng, Narasimhan, and Parkes, AAMAS'18; Golowich, Narasimhan, and Parkes, IJCAI'18; Dütting, Feng, Narasimhan, Parkes, and Ravindranath, ICML'19]

# Approximate incentive compatibility

Mechanism is  $\gamma$ -IC when for each bidder i:

If everyone except bidder i is truthful,

she can only increase utility by  $\gamma$  if she bids strategically

...in expectation over **others'** values EX-INTERIM

(assume bidders independent)

...in expectation over **all** values

EX-ANTE

(no independence assumptions)

# Approximate incentive compatibility

Literature on  $\gamma$ -IC assumes distribution is known in advance



### Where does this knowledge come from?

#### We relax this assumption:

Assume only samples from distribution over agents' types

[Likhodedov and Sandholm, AAAI'04, AAAI'05; Balcan, Blum, Hartline, and Mansour, FOCS'05, JCSS'08; Elkind, SODA'07; Cole and Roughgarden, STOC'14; Mohri and Medina, ICML'14; Huang Mansour, and Roughgarden, EC'15; Sandholm and Likhodedov, OR'15; Morgenstern and Roughgarden, NeurIPS'15, COLT'16; Roughgarden and Schrijvers, EC'16; Devanur, Huang, and Psomas, STOC'16; Balcan, Sandholm, and Vitercik, NeurIPS'16, EC'18; Alon, Babaioff, Gonczarowski, Mansour, Moran, and Yehudayoff, NeurIPS'17; Gonczarowski and Nisan, STOC'17; Cai and Daskalakis, FOCS'17; Syrgkanis, NeurIPS'17, Medina and Vassilvitskii, NeurIPS'17, ...]

Estimate IC approximation factor ( $\gamma$ ) using samples

#### Our estimate (first try):

Maximum utility agent i can gain by misreporting her type,

on average over samples 
$$\left\{oldsymbol{t}_{-i}^{(1)},...,oldsymbol{t}_{-i}^{(N)}
ight\}$$
:

on average over samples 
$$\left\{ \boldsymbol{t}_{-i}^{(1)}, \dots, \boldsymbol{t}_{-i}^{(N)} \right\}$$
:
$$\max_{t_i, t_i' \in \mathbb{R}^D} \left\{ \frac{1}{N} \sum_{j=1}^N u\left(t_i, t_i', \boldsymbol{t}_{-i}^{(j)}\right) - u\left(t_i, t_i, \boldsymbol{t}_{-i}^{(j)}\right) \right\}$$

Utility from strategic bid

Utility from truthful bid

Estimate IC approximation factor ( $\gamma$ ) using samples

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Maximum utility agent i can gain by misreporting her type,

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Might not be finite-time procedure

Estimate IC approximation factor ( $\gamma$ ) using samples

#### Our estimate $\hat{\gamma}$ :

Maximum utility agent i can gain by misreporting her type,

on average over samples 
$$\{\boldsymbol{t}_{-i}^{(1)},...,\boldsymbol{t}_{-i}^{(N)}\}$$
,

if true & reported types from **finite subset** F of type space

$$\hat{\gamma} = \max_{t_i, t_i' \in F} \left\{ \frac{1}{N} \sum_{j=1}^{N} u\left(t_i, t_i', t_{-i}^{(j)}\right) - u\left(t_i, t_i, t_{-i}^{(j)}\right) \right\}$$



Estimate IC approximation factor ( $\gamma$ ) using samples

#### Our estimate $\hat{\gamma}$ :

Maximum utility agent i can gain by misreporting her type, on average over samples  $\{\boldsymbol{t}_{-i}^{(1)},...,\boldsymbol{t}_{-i}^{(N)}\}$ , if true & reported types from **finite subset** F of type space

#### Estimate used in mechanism design via deep learning:

Add constraint requiring this estimate be small

[Feng, Narasimhan, and Parkes, AAMAS'18; Golowich, Narasimhan, and Parkes, IJCAI'18; Dütting, Feng, Narasimhan, Parkes, and Ravindranath, ICML'19]



Estimate IC approximation factor ( $\gamma$ ) using samples

#### Our estimate $\hat{\gamma}$ :

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#### **Challenge:**

Might miss pairs of true & reported types with large utility gains

Estimate IC approximation factor ( $\gamma$ ) using samples

#### Our estimate $\hat{\gamma}$ :

Maximum utility agent i can gain by misreporting her type, on average over samples  $\{oldsymbol{t}_{-i}^{(1)}, ..., oldsymbol{t}_{-i}^{(N)}\}$ , if true & reported types from **finite subset** F of type space

- 1. Which finite subset? 2.  $|\hat{\gamma} \gamma| \leq ?$

#### Which finite subset?

#### 1. Uniform grid



Easy to construct



Works if distribution is "nice"

2. Learning theoretic cover (standard ML theory techniques)



Can be hard to construct



Always works

# Uniform grid

#### **Challenge:**

Utility functions are volatile

First-price auction

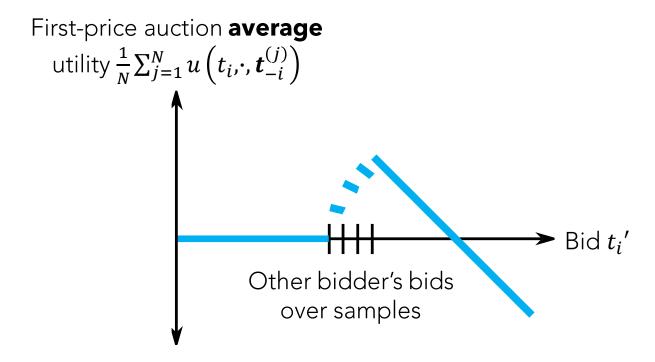
utility  $u(t_i, \cdot, t_{-i})$   $\longrightarrow$  Bid  $t_i'$ Other

bidder's bid

# Uniform grid

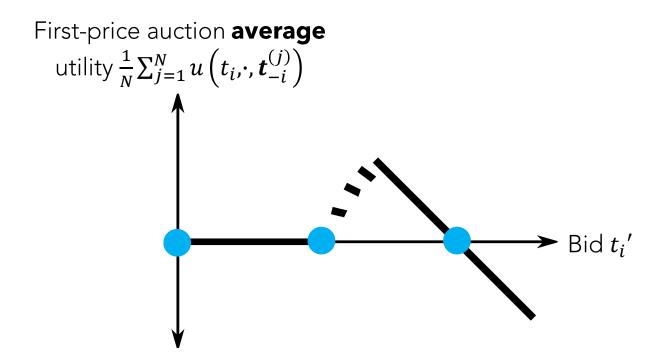
#### **Challenge:**

Utility functions are volatile



# Uniform grid

#### Coarse discretization can lead to poor utility estimation



# When is the distribution "nice" enough to use a grid?

# Dispersion

Functions  $u_1, ..., u_N$  are (w, k)-dispersed if:

Every w-ball contains discontinuities of  $\leq k$  functions

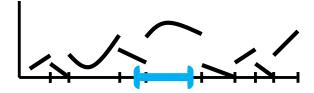
[Balcan, Dick, and Vitercik, FOCS'18]

Plot 
$$\frac{1}{N}\sum u_i$$
:

#### **Not** dispersed

Many discontinuities in interval

#### **Dispersed**

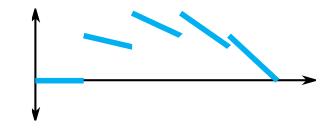


Few discontinuities in interval

#### Our estimate $\hat{\gamma}$ :

Maximum utility agent can gain by misreporting her type, on average over samples, if true & reported types from **finite subset** of type space

**Theorem** (informal): If utility functions induced by N samples are: (w,k)-dispersed and piecewise L-Lipschitz



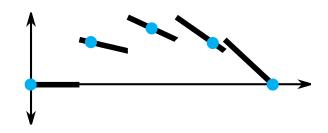
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 $\Rightarrow$  Can use w-grid as finite subset



**Theorem** (informal): If utility functions induced by *N* samples are:

(w,k)-dispersed and piecewise L-Lipschitz

 $\Rightarrow$  Can use w-grid as finite subset

Estimation error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(Lw + \frac{k}{N} + \sqrt{\frac{d}{N}}\right)$$

d = standard ML measure of utility functions' intrinsic complexity

**Theorem** (informal): If utility functions induced by N samples are: (w,k)-dispersed and piecewise L-Lipschitz

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Estimation error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(Lw + \frac{k}{N} + \sqrt{\frac{d}{N}}\right)$$

#### Proof idea:

• If snap types to grid, average utility only changes by  $\leq Lw + \frac{k}{N}$ 

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Estimation error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(Lw + \frac{k}{N} + \sqrt{\frac{d}{N}}\right)$$

#### Proof idea:

- If snap types to grid, average utility only changes by  $\leq Lw + \frac{k}{N}$
- $\sqrt{\frac{d}{N}}$  additional error incurred from sampling

**Theorem** (informal): If utility functions induced by *N* samples are:

(w,k)-dispersed and piecewise L-Lipschitz

 $\Rightarrow$  Can use w-grid as finite subset

**Estimation error:** 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(Lw + \frac{k}{N} + \sqrt{\frac{d}{N}}\right)$$

When 
$$w = O\left(\frac{1}{\sqrt{N}}\right)$$
,  $k = O(\sqrt{N})$ :

We prove these (w, k) values hold when distribution is **nice** 

# **Applications**

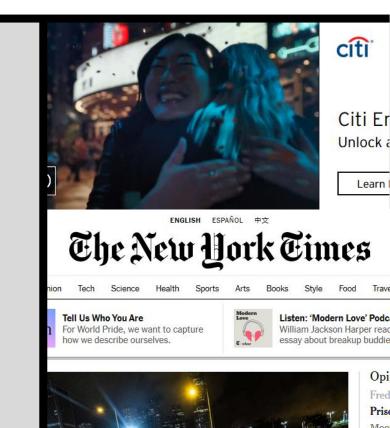
When does dispersion hold?

 $[0, \kappa]$  = range of density functions defining agents' type distributions

#### **First-price auction**

Error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(\frac{\text{(\#bidders)} + \kappa^{-1}}{\sqrt{\text{(\#samples)}}}\right)$$

Also analyze combinatorial first-price auctions



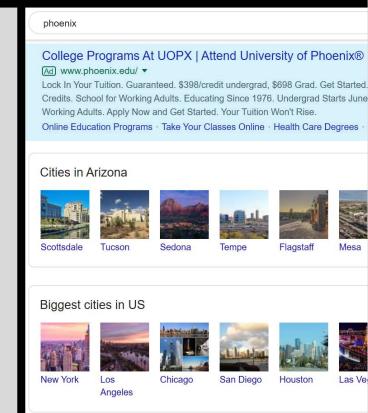
# Applications

When does dispersion hold?

 $[0, \kappa]$  = range of density functions defining agents' type distributions

#### **Generalized second-price auction**

Error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(\frac{(\text{\#bidders})^{3/2} + \kappa^{-1}}{\sqrt{(\text{\#samples})}}\right)$$



# Applications

When does dispersion hold?

 $[0, \kappa]$  = range of density functions defining agents' type distributions

#### Discriminatory and uniform price auctions

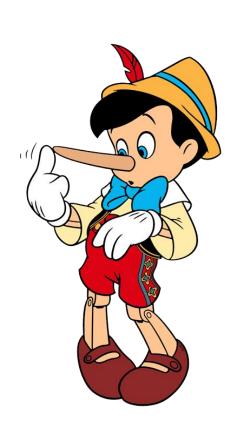
Generalization of first-price auction to multi-unit settings

Error: 
$$|\hat{\gamma} - \gamma| = \tilde{O}\left(\frac{(\text{\#bidders})(\text{\#units})^2 + \kappa^{-1}}{\sqrt{(\text{\#samples})}}\right)$$



#### Conclusion

- Provide techniques for estimating how far mechanism is from IC
- Introduce empirical variant of approximate IC
- Bound estimate's error using dispersion
- Guarantees for:
  - First-price (combinatorial) auction
  - Generalized second-price auction
  - Discriminatory auction
  - Uniform price auction
  - Second-price auction under spiteful agents





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