

SYMBOLIC PROGRAM SLICING ON SMART CONTRACTS

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ABSTRACT

we propose a method to do program slicing on stack-based programming language. With slicing, we can analyse properties more efficiently. We take the EVM bytecode[1] as the example language, which is used to write smart contract[2] on Ethereum[1].

Keywords— BlockChain, Ethereum, Slicing, Verification

1. INTRODUCTION

2. RELATED WORK

3. METHOD

4. EXPERIMENT

5. CONCLUSION

6. REFERENCES

- [1] Gavin Wood, “Ethereum: A secure decentralised generalised transaction ledger eip-150 revision (759dccc - 2017-08-07),” 2017, Accessed: 2018-01-03.
- [2] Nick Szabo, “Advances in distributed security,” 2003, Accessed: 2016-04-31.

Algorithm 1: AnalysisEnvironment

Input: *StackDG*

/ Environment variables */*

```
1 evald =  $\emptyset$ ;
2 conAddrs = dictionary();
3 addrDepPCs = dictionary();
4 conOffsets = dictionary();
5 offsetDepPCs = dictionary();
6 conVals = dictionary();
7 valDepPCs = dictionary();
8 rwDependency = dictionary();
/* Return All dependant program counters */
9 function depPCs(inst)
10   return addrDepPCs[inst.pc]
11      $\cup$  offsetDepPCs[inst.pc]
12      $\cup$  valDepPCs[inst.pc]
13 end function
/* check insts if have same addr parameters */
14 function addrOverlap(insA, insB)
15   rangeA = product(conAddrs[insA.pc],
16                   conOffsets[insA.pc]);
17   rangeB = product(conAddrs[insB.pc],
18                   conOffsets[insB.pc]);
19   for (Aa, Ao), (Ba, Bo)  $\in$  product(rangeA, rangeB) do
20     if {Aa..Aa + Ao}  $\cap$  {Ba..Ba + Bo}  $\neq \emptyset$  then
21       return True;
22   end
23   return False;
24 end function
/* Initialize Environment */
25 function initialize()
26   global Environment variables;
27   readIns = StackDG.readIns ;
28   writeIns = StackDG.writeIns ;
29   /* eval instructions' parameters (addr) */
30   addrs = [ i.pc, eval(i.addrs) | i  $\in$  readIns  $\cup$  writeIns ];
31   conAddrs = { pc: con | (pc, (con, _))  $\in$  addrs };
32   addrDepPCs = { pc: dep | (pc, (_, dep))  $\in$  addrs };
33   /* eval instructions' parameters (offset) */
34   offsets = [ i.pc, eval(i.offsets) | i  $\in$  readIns  $\cup$  writeIns ];
35   conOffsets = { pc: con | (pc, (con, _))  $\in$  offsets };
36   offsetDepPCs = { pc: dep | (pc, (_, dep))  $\in$  offsets };
37   /* eval instructions' parameters (val) */
38   vals = [ i.pc, eval(i.vals) | i  $\in$  writeIns ];
39   conVals = { pc: con | (pc, (con, _))  $\in$  vals };
40   valDepPCs = { pc: dep | (pc, (_, dep))  $\in$  vals };
41   /* write insts that can't be re-evald */
42   evald = {pc  $\in$  addrDepPCs | addrDepPCs[pc] =  $\emptyset$ }
43      $\cap$  {pc  $\in$  offsetDepPCs | offsetDepPCs[pc] =  $\emptyset$ }
44      $\cap$  {pc  $\in$  valDepPCs | valDepPCs[pc] =  $\emptyset$ }
45 end function
```

Algorithm 2: buildAddressDependency

Input: *CFG, StackDG*

Output: *AddressDG*

```
1 import AnalysisEnvironment as env
2 function buildAddressDependency()
3   /* initialize the environment */
4   env.initialize();
5   /* declare and alias variables */
6   visit, alter =  $\emptyset$ ,  $\emptyset$  ;
7   sreads = StackDG.SLOADs ;
8   swrites = StackDG.SSTOREs ;
9   mreads = StackDG.MSTOREs ;
10  mwrites = StackDG.{MLOAD  $\cup$  SHA3
11     $\cup$  CREATE  $\cup$  CALL  $\cup$  RETURN}s
12  /* build addr dependency of */
13  /* storage and memroy */
14  while True do
15    evald = env.evald.copy();
16    buildDepend(swrites, sreads, visit);
17    buildDepend(mwrites, mreads, visit);
18    if exist write inst can be re-evald then
19      for inst  $\in$  re-evald do
20        update Environment variables in env
21        with eval(instruction parameters)
22      end
23    if env.evald  $\setminus$  evald =  $\emptyset$  then
24      break;
25  end
26  return env.rwDependency;
27 end function
/* helper function */
28 function buildDepend(writes, reads, visit)
29   concrete = {ins  $\in$  (writes  $\setminus$  visit)
30     | env.depPCs(ins) =  $\emptyset$ };
31   while concrete  $\neq \emptyset$  do
32     for inst  $\in$  concrete do
33       block = CFG.blockOf(inst);
34       dfsCFG(inst, block, writes, reads,  $\emptyset$ );
35     end
36     visit.update(concrete);
37     for inst  $\in$  (writes  $\setminus$  visit) do
38       if (env.depPCs(inst)  $\setminus$  env.evald) =  $\emptyset$  then
39         update Environment variables in env
40         with eval(instruction parameters)
41       end
42     concrete = {ins  $\in$  (writes  $\setminus$  visit)
43       | env.depPCs(ins) =  $\emptyset$ };
44   end
45 end function
```

Algorithm 3: dfsCFG

Input: *CFG, StackDG*
Output: *AddressDG*

```
1 import AnalysisEnvironment as env
  /* do CFG dfs for building dependency */
2 function dfsCFG(wInst, block, writes, reads, visit)
3   visit.add(block);
4   rwInsts = (block.insts  $\cap$  writes  $\cap$  reads);
5   if wInst  $\in$  block then
6     rwInsts = rwInsts \
7       { ins  $\in$  block.insts | ins.pc > wInst.pc };
8   for inst  $\in$  rwInsts do
9     /* if "exist the probability" to re-write the
       same address then return, "probability"
       means the "or" part */
10    if inst.name = wInst.name and
11      (env.addrOverlap(wInst, inst)
12       or env.addrDepPCs[inst.pc]  $\neq \emptyset$ ) then
13      visit.remove(block);
14      return;
15    if inst  $\in$  reads then
16      depPCs = env.addrDepPCs[inst.pc]  $\cup$ 
17        env.offsetDepPCs[inst.pc];
18      if depPCs  $\neq \emptyset$  and
19        depPCs \ env.ealed =  $\emptyset$  then
20        update Environment variables in env
21        with eval(instruction parameters)
22      if addrOverlap(wInst, inst) then
23        env.ealed.add(pc);
24        env.rwDependency[inst.pc].add(wInst);
25        env.conVals[inst.pc].update(
26          env.conVals[wInst.pc]);
27    end
28  for nextBlock  $\in$  block.outgoingBlock do
29    dfsCFG(wInst, nextBlock, writes, reads, visit);
30  end
31 end function
```

Algorithm 4: eval

Input: *StackDG, instruction, visit*
Output: *concrete values, dependant PCs*

```
1 function eval(inst, visit)
2   if inst  $\in$  visit then
3     return {}, {inst};
4   visit.add(inst);
5   concrete, dependant =  $\emptyset$ ,  $\emptyset$ ;
6   if inst.name.startswith('PUSH') then
7     return {int(op.operand)}, {};
8   cons, deps = {map(eval(_, visit), argList)
9     | argList  $\in$  inst.argLists}T;
10  for argList  $\in$  cons do
11    val = None;
12    if None  $\in$  argList then
13      continue;
14    else if inst.name = 'ADD' then
15      val = let x, y = argList in x + y;
16    else if inst.name = 'SUB' then
17      val = let x, y = argList in x - y;
18    else if inst.name = 'MUL' then
19      val = let x, y = argList in x * y;
20    else if inst.name = 'DIV' then
21      val = let x, y = argList in x / y;
22    else if inst.name = 'EXP' then
23      val = let x, y = argList in xy;
24    else if inst.name = 'ISZERO' then
25      val = let [x] = argList in  $\begin{cases} 0, & \text{if } x = 0 \\ 1, & \text{otherwise} \end{cases}$ 
26    else if inst.name = 'NOT' then
27      val = let [x] = argList in (1 << 256) - 1 - x;
28    else if inst.name = 'AND' then
29      val = let x, y = argList in x & y;
30    else if inst.name = 'OR' then
31      val = let x, y = argList in x | y;
32    else if inst.name = 'EQ' then
33      val = let x, y = argList in x == y;
34    else if inst.name  $\in$  {'MLOAD', 'SLOAD', 'SHA3'}
35      then
36      concrete.update(env.conVals[inst.pc]);
37    else
38      /* SHA3 not impl yet */
39      throw Exception("not handle the inst yet");
40    if val  $\neq$  None then
41      concrete.add(val);
42      dependant.update(concat(deps));
43  end
44  visit.remove(inst);
45  return concrete, dependant;
46 end function
```
