#### hw5

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#### 1.

# 1.a. worse case time complexity: $\Theta(n^2)$

when all element place in the same bucket, and the bucket is sorted by selection sort, then the time complexity is  $\Theta(n^2)$ 

### 1.b. simple change make worst case time complexity to $O(n \log(n))$

use quick sort to sort the bucket,

then the worst-case time complexity is  $O(n \log(n))$ 

### 2.

## 2.a. algorithm

get length from 1 to n,

get max price in each length by split each length into two parts and get the max price in each part, then get the max price in length n.

```
maxPriceInLength[n]
for i in 1 to len(rod):
    for j in 1 to i:
        maxPriceInLength[i] = max(maxPriceInLength[i], price[j] + maxPriceInLength[i-j])
return maxPriceInLength[n]
```

### 2.b. time complexity

 $O(n^2)$ 

#### **3.**

input is length n split the input into A[0,i] and A[i+1,n] then for each char in A[0,i] find the mirror char in  $A[i+1,n] \rightarrow O(n^2)$  i can at from 0 to  $n-1 \rightarrow O(n)$ 

# 3.a. time complexity

$$O(n) * O(n^2) = O(n^3)$$

#### 4.

### 4.a. algorithm

from last server  $S_n$  to first server  $S_1$ , for all Server  $S_x$  if have copy then cost is 0 else calculate the cost of place a copy and the cost of get a copy from  $S_j$  where have copy and closest, then choose the minimum cost.

it can keep last server position to boost speed

# 4.b. time complexity

O(n)