

CSE 214

Computer Organization

Lecture 7

Programming

The Basic Computer

Mostafa I. Soliman
Professor of Computer Engineering
CSE Department

mostafa.soliman@ejust.edu.eg
mossol@yahoo.com



PROGRAMMING THE BASIC COMPUTER

- Introduction
- Machine Language
- Assembly Language
- Assembler
- Program Loops
- Programming Arithmetic Operations
- Programming Logic Operations
- Subroutines
- Input-Output Programming
- Advanced Topic:
Computer Parallelism

Self Reading

Making the simple complicated
is commonplace;
making the complicated
simple, awesomely simple,
that's creativity.
Charles Mingus

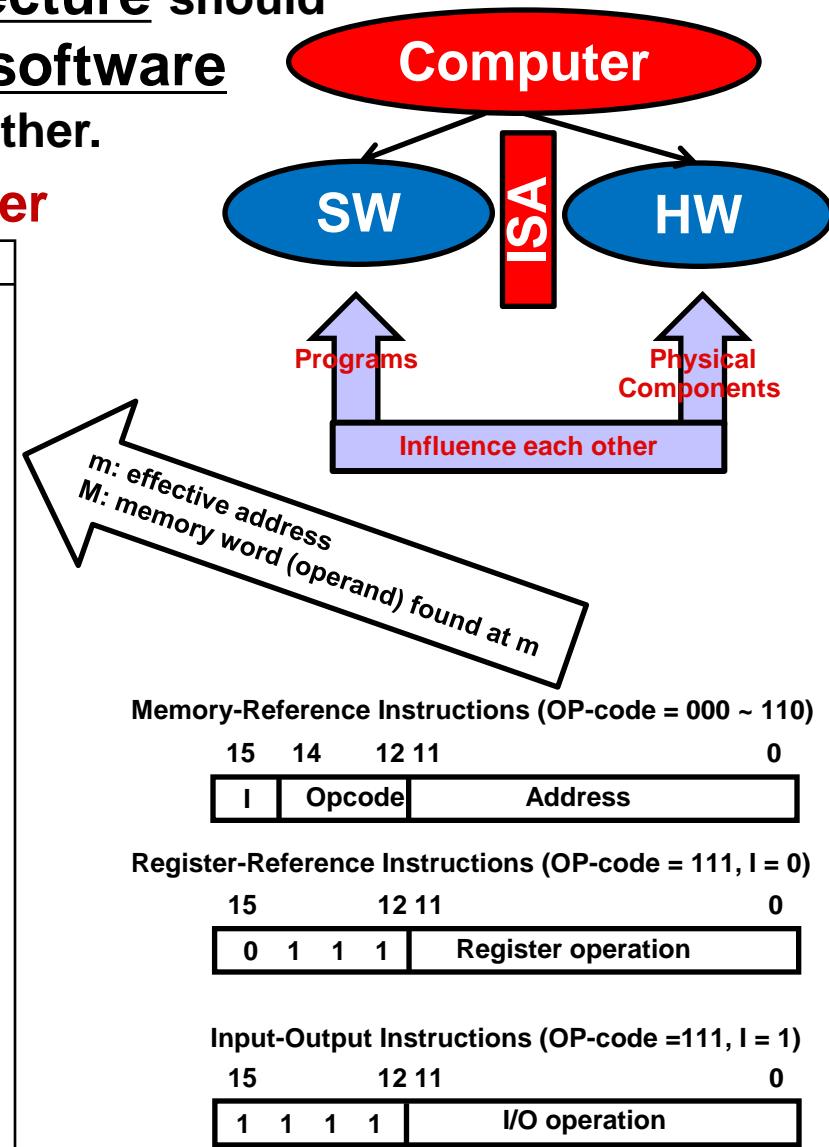
INTRODUCTION

Those concerned with computer architecture should have a knowledge of both hardware and software because the two branches influence each other.

Instruction Set (IS) of the Basic Computer

Symbol	Hexa code	Description
AND	0 or 8	AND M to AC
ADD	1 or 9	Add M to AC, carry to E
LDA	2 or A	Load AC from M
STA	3 or B	Store AC in M
BUN	4 or C	Branch unconditionally to m
BSA	5 or D	Save return address in m and branch to m+1
ISZ	6 or E	Increment M and skip if zero
CLA	7800	Clear AC
CLE	7400	Clear E
CMA	7200	Complement AC
CME	7100	Complement E
CIR	7080	Circulate right E and AC
CIL	7040	Circulate left E and AC
INC	7020	Increment AC, carry to E
SPA	7010	Skip if AC is positive
SNA	7008	Skip if AC is negative
SZA	7004	Skip if AC is zero
SZE	7002	Skip if E is zero
HLT	7001	Halt computer
INP	F800	Input information and clear flag
OUT	F400	Output information and clear flag
SKI	F200	Skip if input flag is on
SKO	F100	Skip if output flag is on
ION	F080	Turn interrupt on
IOF	F040	Turn interrupt off

IS = 25 Instructions



MACHINE LANGUAGE

- **Program (Sequence of machine instructions)**

A list of instructions or statements for directing the computer to perform a required data processing task

- **Various types of programming languages**

- Hierarchy of programming languages

- Machine-language

- Binary code (Programmer)

- Octal or hexadecimal code

- Assembly-language

- (Assembler)

- Symbolic code

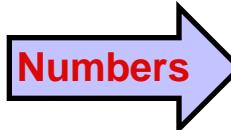
- High-level language

- (Compiler)

COMPARISON OF PROGRAMMING LANGUAGES

- Binary Program to Add Two Numbers

Numbers



Location	Instruction Code
000	0010 0000 0000 0100
001	0001 0000 0000 0101
010	0011 0000 0000 0110
011	0111 0000 0000 0001
100	0000 0000 0101 0011
101	1111 1111 1110 1001
110	0000 0000 0000 0000

- Hexa program

Location	Instruction
000	2004
001	1005
002	3006
003	7001
004	0053
005	FFE9
006	0000

- Program with Symbolic OP-Code

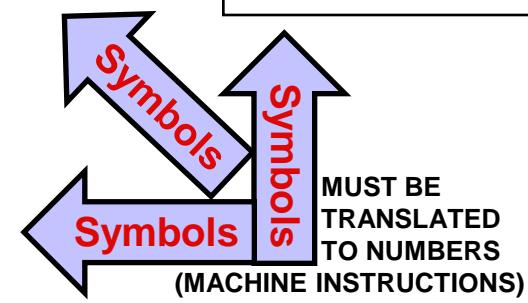
Location	Instruction	Comments
000	LDA 004	Load 1st operand into AC
001	ADD 005	Add 2nd operand to AC
002	STA 006	Store sum in location 006
003	HLT	Halt computer
004	0053	1st operand
005	FFE9	2nd operand (negative)
006	0000	Store sum here

- Assembly-Language Program

ORG	0	/Origin of program is location 0
LDA	A	/Load operand from location A
ADD	B	/Add operand from location B
STA	C	/Store sum in location C
HLT		/Halt computer
A,	DEC 83	/Decimal operand
B,	DEC -23	/Decimal operand
C,	DEC 0	/Sum stored in location C
	END	/End of symbolic program

- C Program

```
main()
{
    int A, B, C ;
    A = 83 ;
    B = -23 ;
    C = A + B ;
}
```



ASSEMBLY LANGUAGE

Syntax of the Basic Computer assembly language

Each line is arranged in three columns called fields

Label field

- May be empty or may specify a symbolic address consists of up to 3 characters
- Terminated by a comma

Instruction field

- Specifies a machine or a pseudo instruction
- May specify one of
 - * Memory reference instruction (MRI)
 - MRI consists of two or three symbols separated by spaces.

ADD OPR (direct address MRI)

ADD PTR I (indirect address MRI)

- * Register reference or input-output instruction

Non-MRI does not have an address part

- * Pseudo instruction with or without an operand

Inform the assembler that

Comment field “ Begin with / ”

- May be empty or may include a comment / for comment

ORG 0	/Origin of program is location 0
LDA A	/Load operand from location A
ADD B	/Add operand from location B
STA C	/Store sum in location C
HLT	/Halt computer
A,	/Decimal operand
B,	/Decimal operand
C,	/Sum stored in location C
END	/End of symbolic program

PSEUDO-INSTRUCTIONS

ORG N

Hexadecimal number N is the memory location
for the instruction or operand listed in the following line

END

Denotes the end of symbolic program

DEC N

Signed decimal number N to be converted to the binary

HEX N

Hexadecimal number N to be converted to the binary

Example: Assembly language program to subtract two numbers

ORG 100	/ Origin of program is location 0x100	
LDA SUB	/ Load subtrahend to AC	
CMA	/ Complement AC	
INC	/ Increment AC	
ADD MIN	/ Add minuend to AC	
STA DIF	/ Store difference	
HLT	/ Halt computer	
MIN,	DEC 83	/ Minuend
SUB,	DEC -23	/ Subtrahend
DIF,	HEX 0	/ Difference stored here
	END	/ End of symbolic program

TRANSLATION TO BINARY

Hexadecimal Code		Symbolic Program
Location	Content	
100	2107	ORG 100
101	7200	LDA SUB
102	7020	CMA
103	1106	INC
104	3108	ADD MIN
105	7001	STA DIF
106	0053	HLT
107	FFE9	MIN, SUB, DIF,
108	0000	DEC 83 DEC -23 HEX 0 END

Symbol	Hexa code	Description
AND	0 or 8	AND M to AC
ADD	1 or 9	Add M to AC, carry to E
LDA	2 or A	Load AC from M
STA	3 or B	Store AC in M
BUN	4 or C	Branch unconditionally to m
BSA	5 or D	Save return address in m and branch to m+1
ISZ	6 or E	Increment M and skip if zero
CLA	7800	Clear AC
CLE	7400	Clear E
CMA	7200	Complement AC
CME	7100	Complement E
CIR	7080	Circulate right E and AC
CIL	7040	Circulate left E and AC
INC	7020	Increment AC, carry to E
SPA	7010	Skip if AC is positive
SNA	7008	Skip if AC is negative
SZA	7004	Skip if AC is zero
SZE	7002	Skip if E is zero
HLT	7001	Halt computer
INP	F800	Input information and clear flag
OUT	F400	Output information and clear flag
SKI	F200	Skip if input flag is on
SKO	F100	Skip if output flag is on
ION	F080	Turn interrupt on
IOF	F040	Turn interrupt off

ASSEMBLER (FIRST PASS)

Assembler

Source Program - Symbolic Assembly Language Program
 Object Program - Binary Machine Language Program

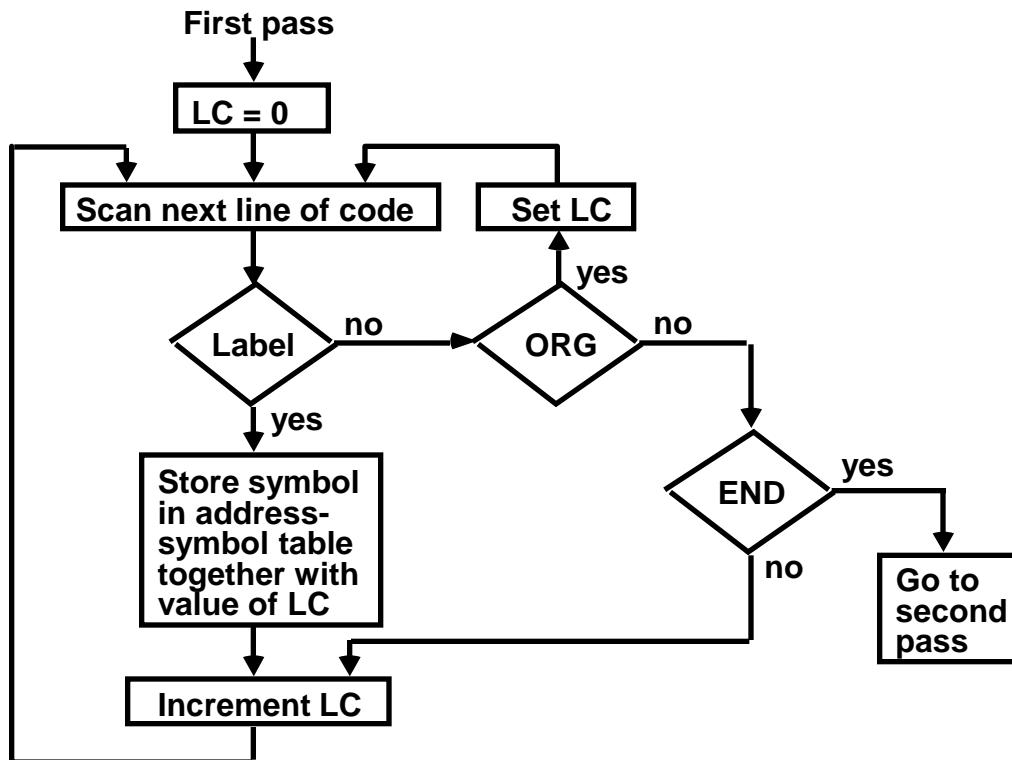
Two pass assembler

1st pass: generates a table that correlates all user defined (address) symbols with their binary equivalent value

2nd pass: binary translation

First pass

Generate
Symbol
Table



LC: Location Counter
 Keep track of the memory location of instructions

Example

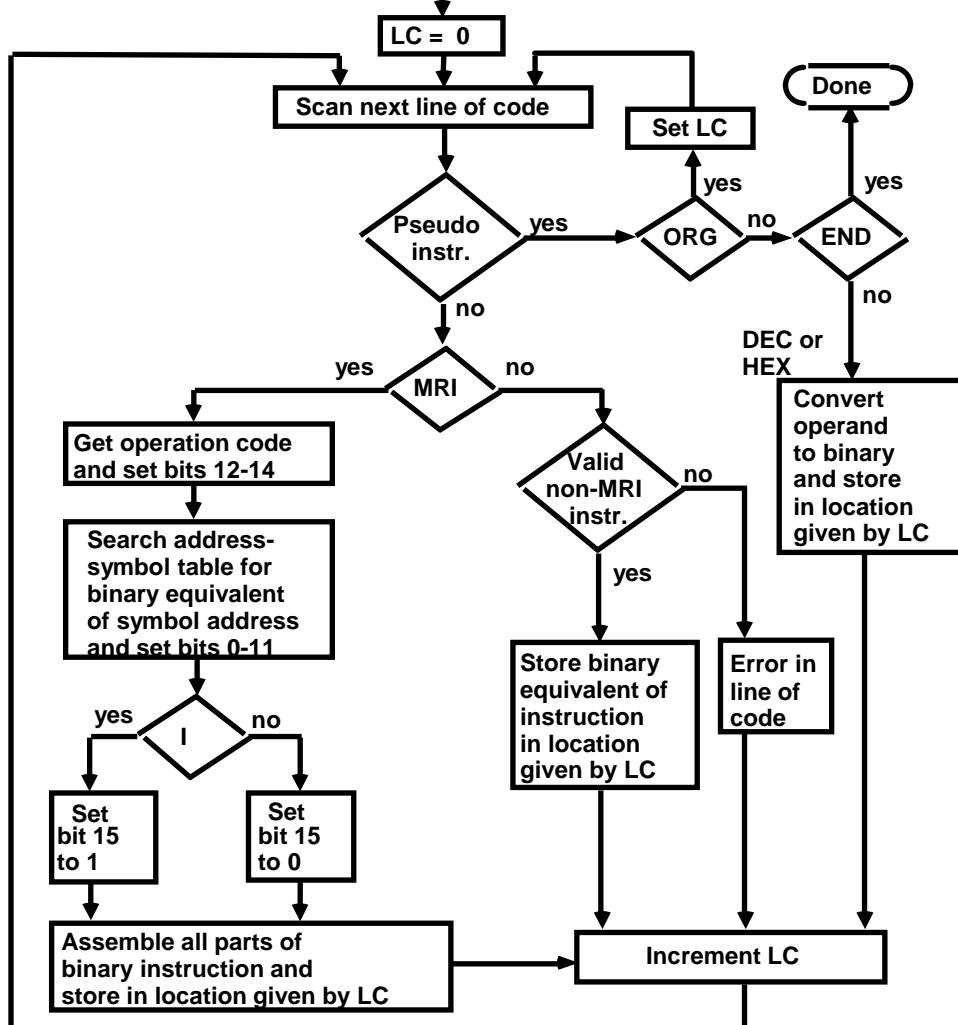
Symbol	LC (Address)
MIN	0x106
SUB	0x107
DIF	0x108

- LC = 0 ORG 100 /Origin of program is location 100
- LC = 100 LDA SUB /Load subtrahend to AC
- LC = 101 CMA /Complement AC
- LC = 102 INC /Increment AC
- LC = 103 ADD MIN /Add minuend to AC
- LC = 104 STA DIF /Store difference
- LC = 105 HLT /Halt computer
- LC = 106 MIN, DEC 83 /Minuend
- LC = 107 SUB, DEC -23 /Subtrahend
- LC = 108 DIF, HEX 0 /Difference stored here
- LC = 109 END /End of symbolic program
- LC = 10A

Memory word	Symbol or (LC)*	Hexadecimal code	Binary representation
1	M I	4D 49	0100 1101 0100 1001
2	N ,	4E 2C	0100 1110 0010 1100
3	(LC)	01 06	0000 0001 0000 0110
4	S U	53 55	0101 0011 0101 0101
5	B ,	42 2C	0100 0010 0010 1100
6	(LC)	01 07	0000 0001 0000 0111
7	D I	44 49	0100 0100 0100 1001
8	F ,	46 2C	0100 0110 0010 1100
9	(LC)	01 08	0000 0001 0000 1000

ASSEMBLER (SECOND PASS)

Machine instructions are translated by means of table-lookup procedures;
Second pass



1. Pseudo-Instruction Table,
2. MRI Table,
3. Non-MRI Table
4. Address Symbol Table

Memory-Reference Instructions (OP-code = 000 ~ 110)

15	14	12 11	0
I	Opcode	Address	

Register-Reference Instructions (OP-code = 111, I = 0)

15	12 11	0
0	1 1	Register operation

Input-Output Instructions (OP-code = 111, I = 1)

15	12 11	0
1 1	1 1	I/O operation

Example

Hexadecimal code

Location	Content	Symbolic program
100	2107	ORG 100
101	7200	LDA SUB
102	7020	CMA
103	1106	INC
104	3108	ADD MIN
105	7001	STA DIF
106	0053	HLT
107	FFE9	MIN, SUB,
108	0000	DEC 83 DEC -23 DIF, HEX 0 END

Pseudo-instruction

ORG	N
DEC	N
HEX	N
END	--

Symbol	LC
MIN	0x106
SUB	0x107
DIF	0x108

Symbol	Hexa code	Description
AND	0 or 8	AND M to AC
ADD	1 or 9	Add M to AC, carry to E
LDA	2 or A	Load AC from M
STA	3 or B	Store AC in M
BUN	4 or C	Branch unconditionally to m
BSA	5 or D	Save return address in m and branch to m+1
ISZ	6 or E	Increment M and skip if zero
CLA	7800	Clear AC
CLE	7400	Clear E
CMA	7200	Complement AC
CME	7100	Complement E
CIR	7080	Circulate right E and AC
CIL	7040	Circulate left E and AC
INC	7020	Increment AC, carry to E
SPA	7010	Skip if AC is positive
SNA	7008	Skip if AC is negative
SZA	7004	Skip if AC is zero
SZE	7002	Skip if E is zero
HLT	7001	Halt computer
INP	F800	Input information and clear flag
OUT	F400	Output information and clear flag
SKI	F200	Skip if input flag is on
SKO	F100	Skip if output flag is on
ION	F080	Turn interrupt on
IOF	F040	Turn interrupt off

MRI

Non-MRI

PROGRAM LOOPS

Loop: A sequence of instructions that are executed many times, each with a different set of data

Program to accumulate 100 numbers

Assembly-language program to accumulate 100 numbers:

	ORG 100	/ Origin of program is HEX 100
	LDA ADS	/ Load first address of operand
	STA PTR	/ Store in pointer
	LDA NBR	/ Load -100
	STA CTR	/ Store in counter
	CLA	/ Clear AC
LOP,	ADD PTR I	/ Add an operand to AC
	ISZ PTR	/ Increment pointer
	ISZ CTR	/ Increment counter
	BUN LOP	/ Repeat loop again
	STA SUM	/ Store sum
	HLT	/ Halt
ADS,	HEX 150	/ First address of operands
PTR,	HEX 0	/ Reserved for a pointer
NBR,	DEC -100	/ Initial value for a counter
CTR,	HEX 0	/ Reserved for a counter
SUM,	HEX 0	/ Sum is stored here
	ORG 150	/ Origin of operands is HEX 150
	DEC 75	/ First operand
	:	
	DEC 23	/ Last operand
	END	/ End of symbolic program

PROGRAMMING ARITHMETIC AND LOGIC OPERATIONS

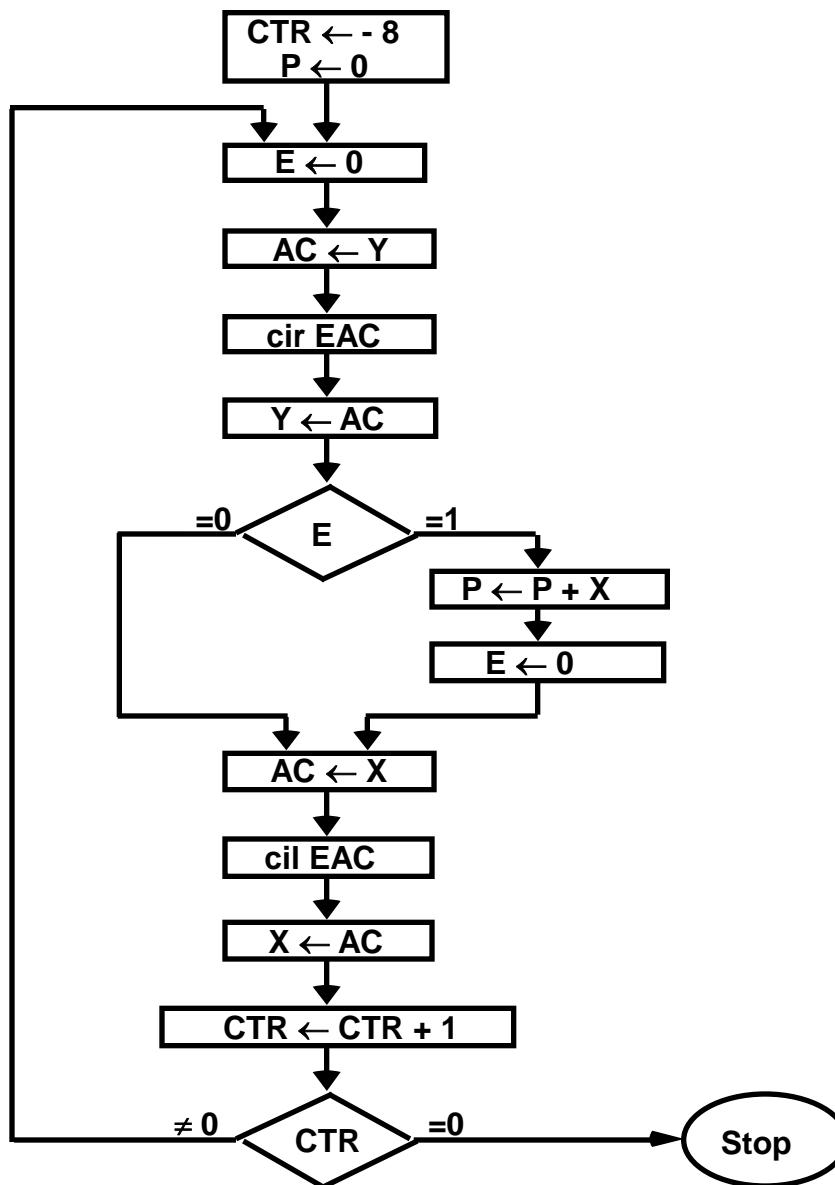
Implementation of Arithmetic and Logic Operations

- Software Implementation
 - Implementation of an operation with a program using machine instruction set
 - Usually when the operation is not included in the instruction set
- Hardware Implementation
 - Implementation of an operation in a computer with one machine instruction

Software Implementation example:

- * Multiplication
 - For simplicity, unsigned positive numbers
 - 8-bit numbers -> 16-bit product

FLOWCHART OF A PROGRAM - Multiplication -



X holds the multiplicand
Y holds the multiplier
P holds the product

Example with four significant digits

$$\begin{array}{r}
 \begin{array}{r} X = 0000\ 1111 \\ Y = 0000\ 1011 \end{array} \quad \begin{array}{r} P \\ \hline 0000\ 0000 \end{array} \\
 \begin{array}{r} 0000\ 1111 \\ 0001\ 1110 \\ 0000\ 0000 \\ 0111\ 1000 \end{array} \quad \begin{array}{r} 0000\ 1111 \\ 0010\ 1101 \\ 0010\ 1101 \\ 1010\ 0101 \end{array} \\
 \hline
 \begin{array}{r} 1010\ 0101 \end{array}
 \end{array}$$

ASSEMBLY LANGUAGE PROGRAM - Multiplication -

```
ORG 100
LOP,    CLE          / Clear E
        LDA Y        / Load multiplier
        CIR          / Transfer multiplier bit to E
        STA Y        / Store shifted multiplier
        SZE          / Check if bit is zero
        BUN ONE      / Bit is one; goto ONE
        BUN ZRO      / Bit is zero; goto ZRO
ONE,    LDA X        / Load multiplicand
        ADD P        / Add to partial product
        STA P        / Store partial product
        CLE          / Clear E
ZRO,    LDA X        / Load multiplicand
        CIL          / Shift left
        STA X        / Store shifted multiplicand
        ISZ CTR      / Increment counter
        BUN LOP      / Counter not zero; repeat loop
        HLT          / Counter is zero; halt
CTR,    DEC -8       / This location serves as a counter
X,      HEX 000F     / Multiplicand stored here
Y,      HEX 000B     / Multiplier stored here
P,      HEX 0         / Product formed here
END
```

ASSEMBLY LANGUAGE PROGRAM

- Double Precision Addition -

E

A = AH AL

+ + +

B = BH BL

C = CH CL

```
LDA AL      / Load A low
ADD BL      / Add B low, carry in E
STA CL      / Store in C low
CLA          / Clear AC
CIL          / Circulate to bring carry into AC(16)
ADD AH      / Add A high and carry
ADD BH      / Add B high
STA CH      / Store in C high
HLT          
```

Self Reading

ASSEMBLY LANGUAGE PROGRAM

- Logic and Shift Operations -

- Logic operations

- BC instructions : AND, CMA, CLA
- Program for OR operation

$$\overline{\overline{A}} \text{ OR } \overline{\overline{B}} = \overline{\overline{A}} \text{ AND } \overline{\overline{B}}$$

LDA A	/ Load 1st operand
CMA	/ Complement to get A'
STA TMP	/ Store in a temporary location
LDA B	/ Load 2nd operand B
CMA	/ Complement to get B'
AND TMP	/ AND with A' to get A' AND B'
CMA	/ Complement again to get A OR B

- Shift operations - BC has *Circular Shift* only

- Logical shift-right operation
- Logical shift-left operation

CLE
CIR

CLE
CIL

- Arithmetic right-shift operation

CLE	/ Clear E to 0
SPA	/ Skip if AC is positive
CME	/ AC is negative
CIR	/ Circulate E and AC

SUBROUTINES

Subroutine

- A set of common instructions that can be used in a program many times.
- Subroutine *linkage* : a procedure for branching to a subroutine and returning to the main program

Example

<i>Loc.</i>			
100		ORG 100	/ Main program
101		LDA X	/ Load X
102		BSA SH4	/ Branch to subroutine
103		STA X	/ Store shifted number
104		LDA Y	/ Load Y
105		BSA SH4	/ Branch to subroutine again
106		STA Y	/ Store shifted number
107	X,	HLT	
108	Y,	HEX 1234	
			/ Subroutine to shift left 4 times
109	SH4,	HEX 0	/ Store return address here
10A		CIL	/ Circulate left once
10B		CIL	
10C		CIL	
10D		CIL	/ Circulate left fourth time
10E		AND MSK	/ Set AC(13-16) to zero
10F		BUN SH4 I	/ Return to main program
110	MSK,	HEX FFF0	/ Mask operand
		END	

SUBROUTINE PARAMETERS AND DATA LINKAGE

Linkage of Parameters and Data between the Main Program and a Subroutine

- via Registers
- via Memory locations
-

Example: Subroutine performing *LOGICAL OR operation*; Need two parameters

<i>Loc.</i>		
200		ORG 200
201		LDA X / Load 1st operand into AC
202		BSA OR / Branch to subroutine OR
203		HEX 3AF6 / 2nd operand stored here
204		STA Y / Subroutine returns here
205	X,	HLT
206	Y,	HEX 7B95 / 1st operand stored here
207	OR,	HEX 0 / Result stored here
208		CMA / Subroutine OR
209		STA TMP / Complement 1st operand
20A		LDA OR I / Store in temporary location
20B		CMA / Load 2nd operand
20C		AND TMP / Complement 2nd operand
20D		CMA / AND complemented 1st operand
20E		ISZ OR / Complement again to get OR
20F		BUN OR I / Increment return address
210	TMP,	HEX 0 / Return to main program
		END / Temporary storage

SUBROUTINE - Moving a Block of Data -

	/ Main program
BSA MVE	/ Branch to subroutine
HEX 100	/ 1st address of source data
HEX 200	/ 1st address of destination data
DEC -16	/ Number of items to move
HLT	
MVE,	HEX 0 / Subroutine MVE
	LDA MVE I / Bring address of source
	STA PT1 / Store in 1st pointer
	ISZ MVE / Increment return address
	LDA MVE I / Bring address of destination
	STA PT2 / Store in 2nd pointer
	ISZ MVE / Increment return address
	LDA MVE I / Bring number of items
	STA CTR / Store in counter
	ISZ MVE / Increment return address
LOP,	LDA PT1 I / Load source item
	STA PT2 I / Store in destination
	ISZ PT1 / Increment source pointer
	ISZ PT2 / Increment destination pointer
	ISZ CTR / Increment counter
	BUN LOP / Repeat 16 times
	BUN MVE I / Return to main program
PT1,	--
PT2,	--
CTR,	--

- Fortran subroutine

```

SUBROUTINE MVE (SOURCE, DEST, N)
DIMENSION SOURCE(N), DEST(N)
DO 20 I = 1, N
20 DEST(I) = SOURCE(I)
      RETURN
      END
    
```

INPUT OUTPUT PROGRAM

Program to Input one Character(Byte)

CIF,	SKI	/ Check input flag
	BUN CIF	/ Flag=0, branch to check again
	INP	/ Flag=1, input character
	OUT	/ Display to ensure correctness
	STA CHR	/ Store character
	HLT	
CHR,	--	/ Store character here

Program to Output a Character

COF,	LDA CHR	/ Load character into AC
	SKO	/ Check output flag
	BUN COF	/ Flag=0, branch to check again
	OUT	/ Flag=1, output character
	HLT	
CHR,	HEX 0057	/ Character is "W"

CHARACTER MANIPULATION

Subroutine to Input 2 Characters and pack into a word

```
IN2, --          / Subroutine entry
FST, SKI
BUN FST
INP           / Input 1st character
OUT
BSA SH4       / Logical Shift left 4 bits
BSA SH4       / 4 more bits
SCD, SKI
BUN SCD
INP           / Input 2nd character
OUT
BUN IN2 I    / Return
```

PROGRAM INTERRUPT

Tasks of Interrupt Service Routine

- Save the Status of CPU
 Contents of processor registers and Flags
- Identify the source of Interrupt
 Check which flag is set
- Service the device whose flag is set
 (Input Output Subroutine)
- Restore contents of processor registers and flags
- Turn the interrupt facility on
- Return to the running program
 Load PC of the interrupted program

INTERRUPT SERVICE ROUTINE

<i>Loc.</i>			
0	ZRO,	-	/ Return address stored here
1		BUN SRV	/ Branch to service routine
100		CLA	/ Portion of running program
101		ION	/ Turn on interrupt facility
102		LDA X	
103		ADD Y	/ Interrupt occurs here
104		STA Z	/ Program returns here after interrupt
200	SRV,	STA SAC	/ Interrupt service routine
		CIR	/ Store content of AC
		STA SE	/ Move E into AC(1)
		SKI	/ Store content of E
		BUN NXT	/ Check input flag
		INP	/ Flag is off, check next flag
		OUT	/ Flag is on, input character
		STA PT1 I	/ Print character
	NXT,	ISZ PT1	/ Store it in input buffer
		SKO	/ Increment input pointer
		BUN EXT	/ Check output flag
		LDA PT2 I	/ Flag is off, exit
		OUT	/ Load character from output buffer
	EXT,	ISZ PT2	/ Output character
		LDA SE	/ Increment output pointer
		CIL	/ Restore value of AC(1)
		LDA SAC	/ Shift it to E
		ION	/ Restore content of AC
		BUN ZRO I	/ Turn interrupt on
	SAC,	-	/ Return to running program
	SE,	-	/ AC is stored here
	PT1,	-	/ E is stored here
	PT2,	-	/ Pointer of input buffer
			/ Pointer of output buffer

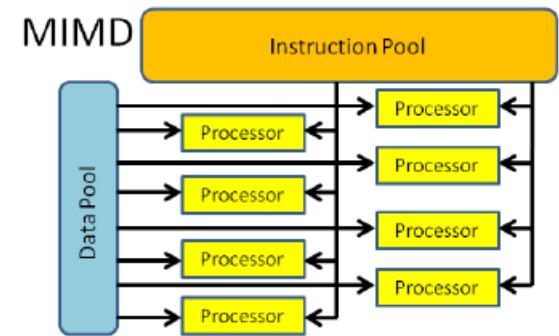
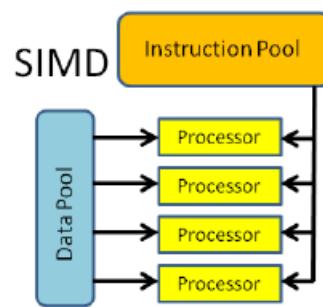
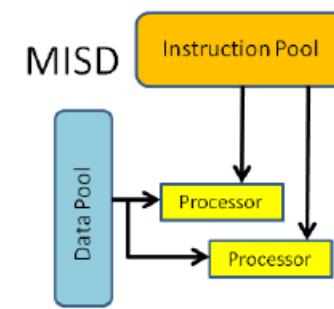
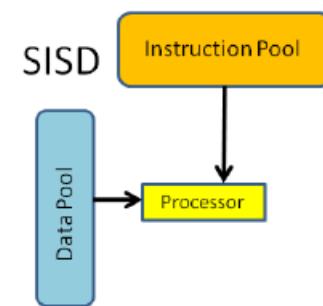
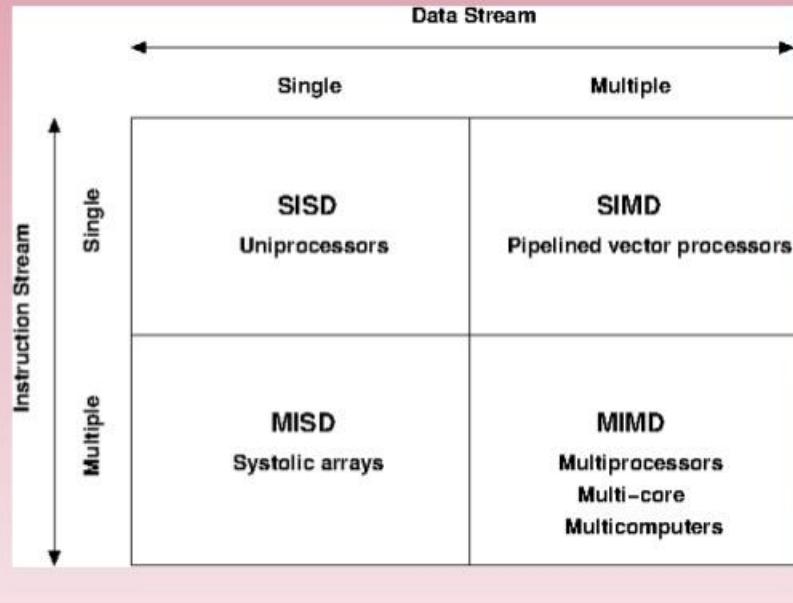
Advanced Topic

Computer Parallelism

Flynn's Taxonomy

- First proposed by Michael J. Flynn in 1966, Flynn's taxonomy is a specific classification of parallel computer architectures that are based on the number of concurrent instruction (single or multiple) and data streams (single or multiple) available in the architecture. The four categories in Flynn's taxonomy are the following:
- (SISD) single instruction, single data
- (MISD) multiple instruction, single data
- (SIMD) single instruction, multiple data
- (MIMD) multiple instruction, multiple data

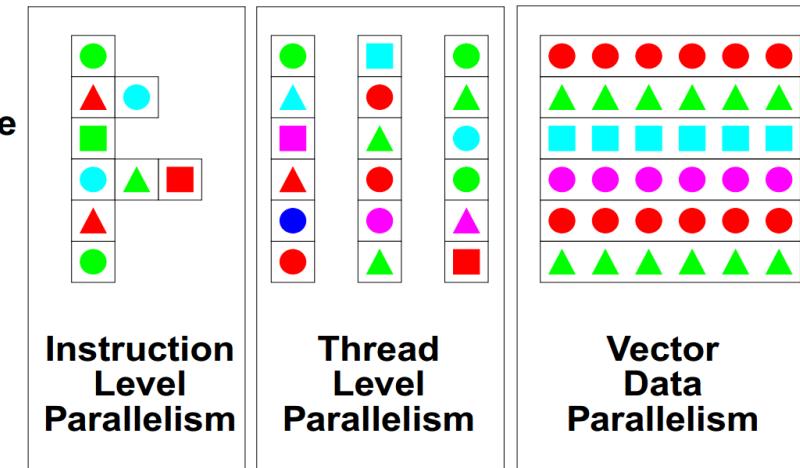
Flynn's Classification



Computer Parallelism

- Taking advantage of parallelism is one of the most important methods for improving performance.
- Classes of architectural parallelism:

- Instruction-Level Parallelism (ILP)
- Thread-Level Parallelism (TLP)
- Data-Level Parallelism (DLP)



- Compiler and hardware conspire to exploit ILP implicitly without the programmer's attention
- DLP and TLP are explicitly parallel
 - requiring the restructuring of the application so that it can exploit explicit parallelism.

Few Powerful or Many Simple Cores?

*If you were plowing a field,
which would you rather use:
two strong oxen
or
1024 chickens?*

Seymour Cray, Father of the Supercomputer
(arguing for two powerful vector processors versus many simple
processors)

References

- M. Mano, “Computer System Architecture,” Pearson Publisher, 3rd Edition, 1992.
- D. Patterson and J. Hennessy, “Computer Organization and Design: the Hardware/Software Interface,” Morgan Kaufmann; 5th Edition, 2014.
- M. Mano, “Digital Design: With an Introduction to the Verilog HDL, VHDL, and SystemVerilog,” Pearson Publisher, 6th Edition, 2017.
- M. Flynn, “Some Computer Organizations and Their Effectiveness,” IEEE Trans. Computers, C-21, No.9, Sept. 1972, pp. 948-960.
- A. Legrand, Parallel Computing: from KiloFlops to Exascale. Evolution in the Last Decades and Recent Challenges, http://polaris.imag.fr/arnaud.legrand/teaching/2013/PC_01_parallel_architectures.pdf

Thank you