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May 26, 2014

Composable Data Consistency Policies

Abstract	
off strong consistency for availability and performance, s algorithms exhibiting benign data races. Developing such may vary across system components or for different syste underlying environment often provide weaker or strong code complex interaction patterns among the environm consistency properties for each component, without com components. While ongoing research has produced several method opers reason about, express and check consistency policie models, as employed in distributed systems, or around stature. Hence, we do not have a unified theory to describ scales, neither we have comprehensive semantics to descend in this research prospectus, we overview classic and recefor further research. We believe that studying consistent the way to creating elegant, generic frameworks and prefercy management patterns, in both distributed systems	me evident, with an increasing variety of systems trading uch as eventually consistent replicated key-value stores, or systems is a challenging, because consistency requirements em clients; even when uniform consistency is required, the er consistency than desired. Thus, developers have to enent and various system components to obtain the desired promising consistency, performance or availability of other dologies, tools, and programming models that assist develops, most work is polarized either around relaxed consistency rict models described in shared-memory concurrency litered diverse, commonly used consistency policies on different cribe how various consistency models interact or compose. Ent work on managing consistency, and motivate the need cry as a single, universal property of applications will pave ogramming languages to express a wide variety of consistency policies can lead to better safety and performance property policies can lead to better safety and performance property policies can lead to better safety and performance property policies can be described to be the safety and performance property policies can lead to better safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be described to be the safety and performance property policies can be a safety and performance property policies can
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Research Advisor's approval (Prof. Nate Nystrom):	Date:
PhD Director's approval (Prof. Stefan Wolf):	Date:

1 Introduction and Problem Domain

1.1 Replication and Partitioning

Data replication and partitioning appear in hardware and software systems at every scale. Consider, on a small scale, the replication of data among caches and main memory, or partitioning program data among objects in objects oriented systems, to be later composed and used as logically meaningful modules. Replicating data across web servers, or partitioning data across tables in a database are some larger scale examples. Both replication and partitioning introduce several challenges to system designers and programmers. A major challenge is the problem of data consistency: that an update to a replica or partition will be propagated correctly to clients of other replicas or partitions.

1.2 Correctness of Concurrent Operations

The notion of correctness means different things in different contexts, depending on different requirements. Traditionally (mostly before the internet era), correctness meant that systems behave as if all operations were performed serially at one replica or across all partitions, meaning that operations form a total order. examples of consistency models that possess this quality are one-copy serializability [Bernstein et al. 1987]. sequential consistency [Leslie some year] or linearizability [wing some year]. Enforcing a total order on operations occurring within a large distributed system, with many replicas can prove expensive, as well as unnecessary for many applications. In this era of massive, distributed online applications, Weaker models, such as eventual consistency [Sheth et al. 1991] or session consistency [Ref], are the norm not the exception. In fact, common issues in distributed systems, such as loss of data during communication or inavailability of processing nodes makes it challenging to enforce even those weaker models. Weak consistency models can be found on a much smaller scale, in the absence of replication as well. Class ConcurrentSkipListMap[ref], a part of Java standard library, contains execution paths that do not correspond to any sequential execution, because some carefully characterized operations that lost races are not retired, the consistency of instances is enforced by the semantics of the class.

1.3 Enforcing and Strengthening Consistency

Enforcing data consistency requires coordinating update requests, and determining when those updates become visible to clients. There are several mechanisms for achieving this, depending on the required consistency and performance properties such as hardware level CAS operations, explicit synchronization through locking, transactionsâĂę etc. I want to say here that all these mechanisms must determine what to do on failure of operation, specify the retry policy