- 1. Give the Θ for the following and justify your answer:
 - (a) $5x^2 + 4x + 3$
 - (b) $2^n + n!$
 - (c) $n^2 + 2^n$
 - (d) $\log(n) + n$
 - (e) log(n!)

2. Give a closed form for the following, then give the $\boldsymbol{\Theta}$

(a)
$$a_0 = 5$$

 $a_n = 3a_{n-1}$

(b)
$$a_4 = 2$$

 $a_n = a_{n-1} + \log_2(n)$

(c)
$$a_1 = 1$$

 $a_n = 2a_{n-2} + 1$

(d)
$$T(1) = 1$$

 $T(k) = 3T(k/2) + 1$

(e)
$$T(1) = 4$$

 $T(k) = T(k/3) + 4$

3. (Extra Credit):

$$T(0) = 1$$

 $T(k) = 3T(k-2) + 4(k-2) + 2$

- 4. Prove **theorem 2:** $x^k \in \mathbf{O}(x^{k+c})$
- 5. Prove **theorem 3:** $x^k + c \cdot x^{k-r} \in \mathbf{O}(x^k)$
- 6. Prove **theorem 5:** if $f(n) \in \mathbf{O}(g(n))$ and $g(n) \in \mathbf{O}(h(n))$, then $f(n) \in \mathbf{O}(h(n))$

7. Give the $\boldsymbol{\Theta}$ running time for the following selection sort algorithm

```
def selSort(1):
    for i in range(len(1)):
        min = 1[i]
        minI = i
        for j in range(i,len(1)):
            if 1[j] < min:
                  minI = j
                  min = 1[j]
                  #end if
            # end for
            (1[i], min) = (min, 1[i])
            # end for</pre>
```

8. Give the recurrence relation for badSort. Remember 1[a:b] copies the elements from 1[a] to 1[b], so even though it's an expression 1[a:b] runs in n-2 time.

```
def badSort(1): n = len(1)

if n == 1:
    return l

first = badSort(1[0:n-2])
  middle = badSort(1[1:n-1])
  end = badSort(1[2:n])

return [first[0]] + middle + [end[n-1]]
```

9. The following algorithm is the merge sort we way in class

```
def merge(low, high):
                                         def mergeSort(lst):
    i = 0
                                              n = len(lst)
    j = 0
                                              n2 = int(n/2)
    merged = []
    while i < len(low) and j < len(high):
                                             # base case:
        if low[i] < high[j]:</pre>
                                             if n <= 1:
            merged += [low[i]]
                                                  return 1st
            i += 1
                                              # recursive case:
        else:
                                              low = mergeSort(lst[0:n2])
            merged += [high[j]]
            j += 1
                                              high = mergeSort(lst[n2:n])
                                              lst = merge(low,high)
    return merged + low[i:] + high[j:]
                                              return 1st
```

- (a) Give the Θ running time for merge. hint: what is the input size for merge?
- (b) Use (a) to give a recurrence relation for the running time of mergeSort.
- (c) Solve the recurrence to get a Θ running time for mergesort.