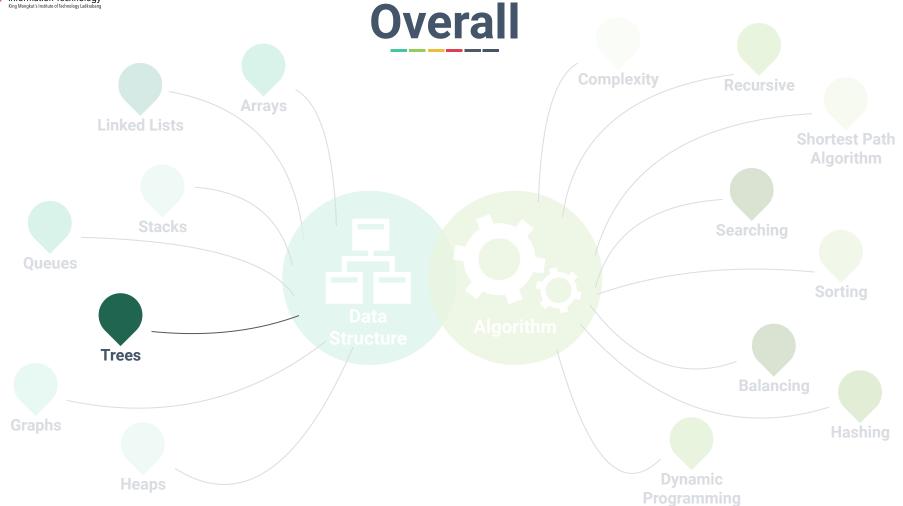


Chapter 6: Trees

Dr. Sirasit Lochanachit







Abstract Data Type

Linear:

- Arrays, Stacks, Queues, and Linked Lists
- Operations: push, pop, enqueue, dequeue
- Algorithms: Searching and sorting, insertion, deletion

Non-Linear:

- Trees and Graphs
- Algorithms: Traversal, Insertion and Deletion, Balancing, and etc.



Today's Outline

General Trees:

- Definition and examples
- Elements of Tree Structure

Binary Trees:

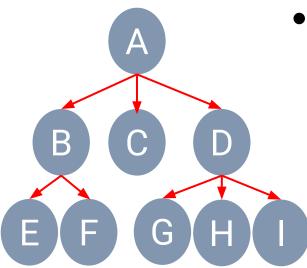
Definition, examples, properties, and types

Tree traversal algorithms:

- Depth-first Traversal (Preorder, postorder, and inorder)
- Breadth-first Traversal



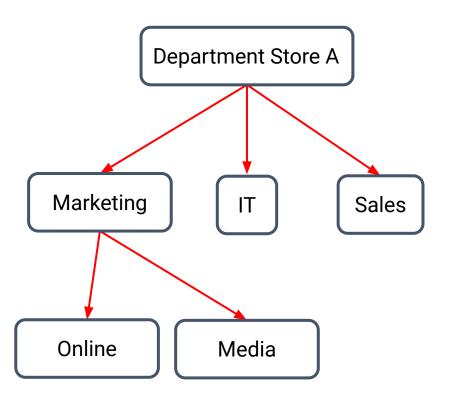
What is a Tree?



A tree stores elements in a hierarchical structure.



Tree Example

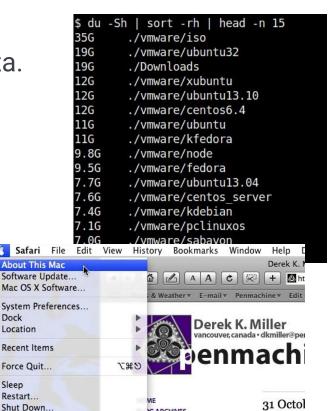






Tree Applications

- Tree represents natural organisation for data.
- Tree structure has been used widely in
 - File systems (Directory)
 - Graphical User Interfaces
 - Databases (Sub-categories, etc.)
 - Websites

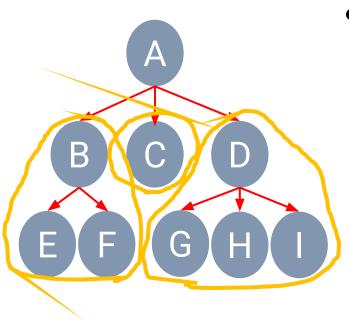


Log Out Derek K. Miller.

Retrieved from https://live.staticflickr.com/7315/13167827344_f2a0ab2015_o_d.png CC BY 2.0 https://live.staticflickr.com/2261/1817203755_c0b7db3f64_o_d.jpg CC BY-NC 2.0



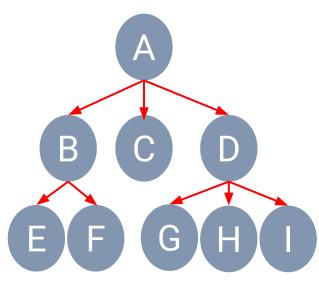
Formal Definition



• A **tree** is a collection of nodes that store elements with a **parent-child** relationship.



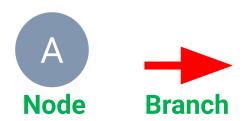
Formal Definition

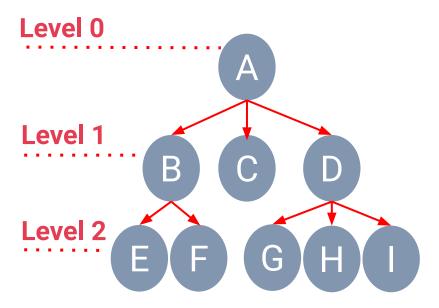


- Siblings
- External or <u>leaf</u> node
- Internal
- Ancestor
 - **Descendant**



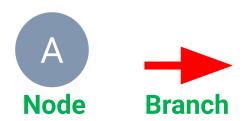
Basic Elements of Tree Structure

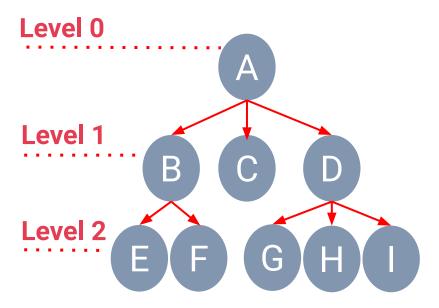






Basic Elements of Tree Structure





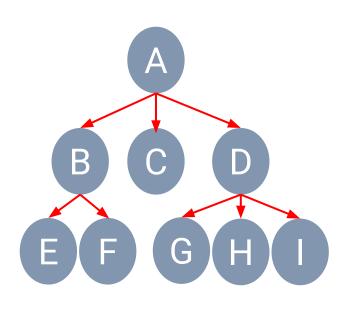


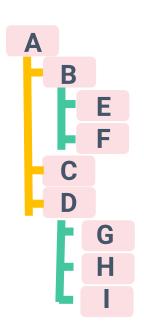
Implementation of Trees

Tree



Other Implementation of Trees





A{ B{E, F}, C, D{G, H, I} }

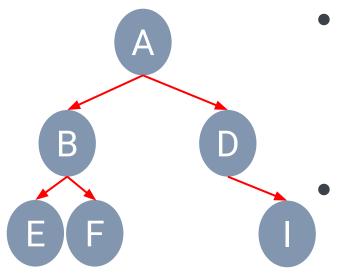
(a) General Form

(b) Tab Form

(c) Set Form



Binary Trees

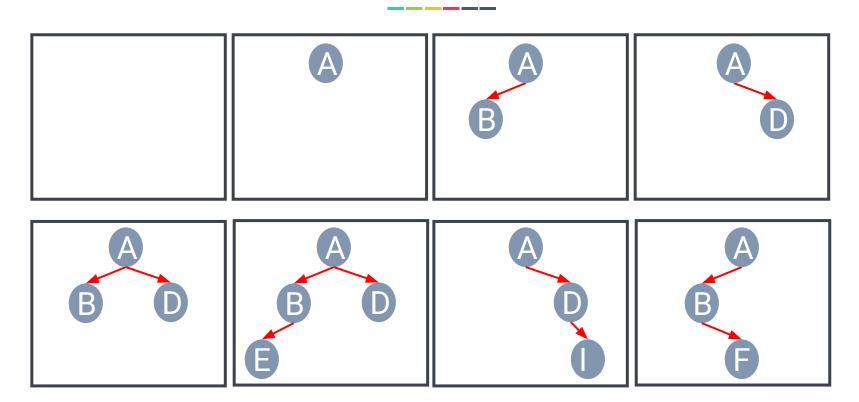


 A binary tree is a tree in which any node can have only two children at maximum.

Each child node is designated as being either a **left child** or **right child**.



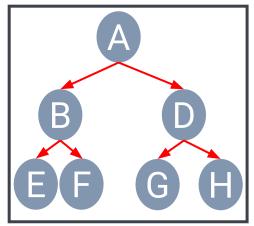
Binary Trees

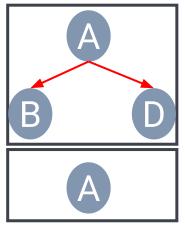


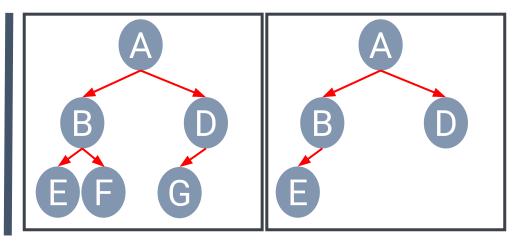


Types of Binary Trees

 A binary tree is nearly complete when every level, except the last, is completely filled and all leaf nodes are <u>as far left as possible</u>.







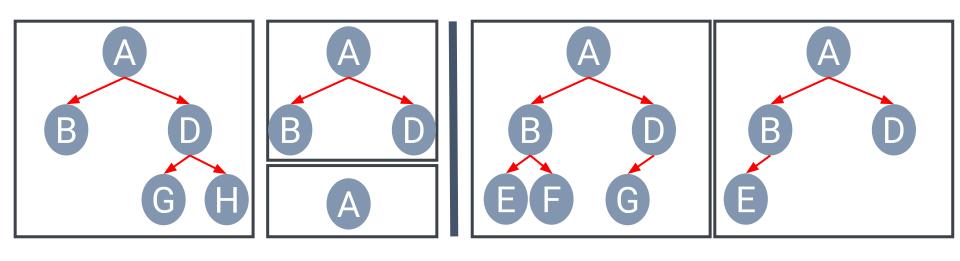
(a) Complete/Perfect Binary Tree

(b) Nearly Complete Binary Tree at level 2



Types of Binary Trees

• A binary tree is proper or full when every node has zero or two children.

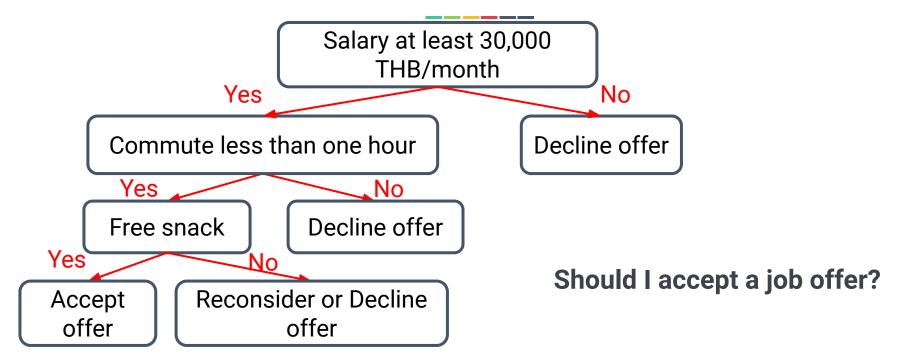


(a) **Proper/Full** Binary Tree
Each node has either 0 or 2 children

(b) Improper Binary Tree

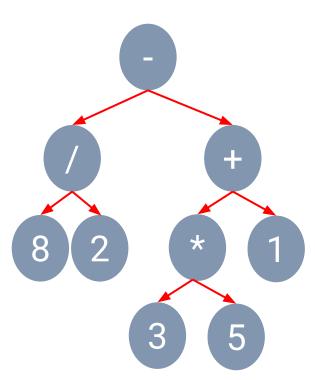


Examples of Binary Trees





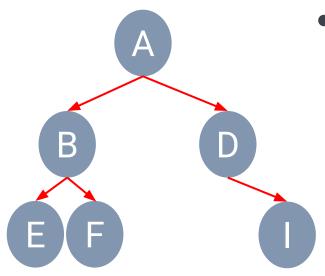
Examples of Binary Trees



- Binary trees can be used to represent an arithmetic expression.
- Leaves are associated with variables or constants.
- Root and Internal nodes are associated with operators.
- Subtree is a sub-expression.



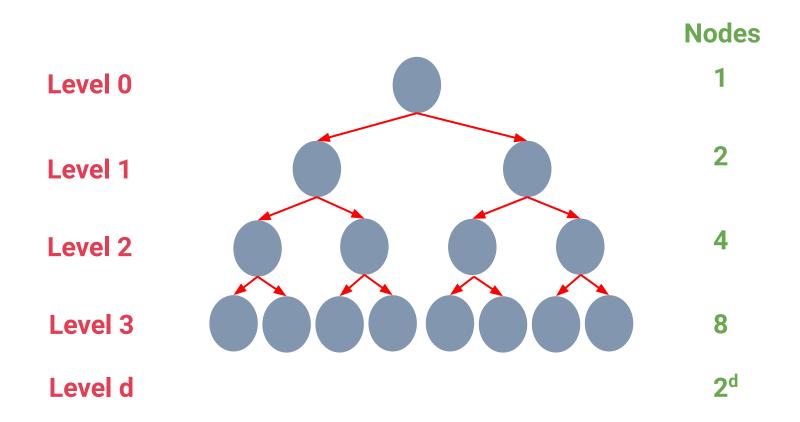
Recursive Binary Trees



- A binary tree is either empty or consists of:
 - A root node that stores an element
 - A binary left subtree (possibly empty)
 - A binary right subtree (possibly empty)

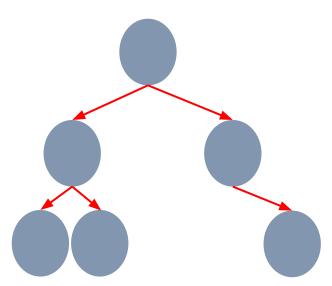


Binary Trees' Properties





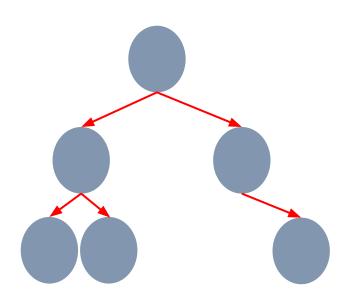
Binary Trees' Properties



- *n* denotes the number of nodes
- h denotes the height of tree



Binary Trees Implementation



Arrays (List of Lists)

or

Doubly Linked Lists

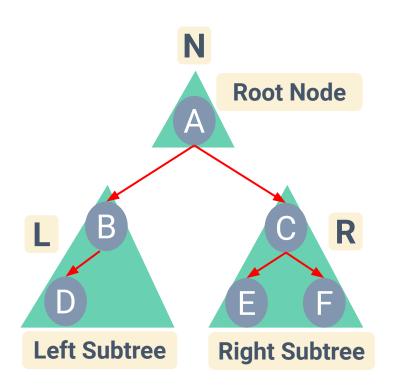


Binary Tree Traversal

- A traversal of a tree is a method to access all the tree nodes.
- Two main approaches: Depth-first and Breadth-first
 - Depth-first
 - Preorder traversal
 - Postorder traversal
 - Inorder traversal
 - \circ Breadth-first visit all the nodes at depth d before moving to depth d+1.



Depth-first Traversal



Preorder Traversal



Inorder Traversal



Postorder Traversal

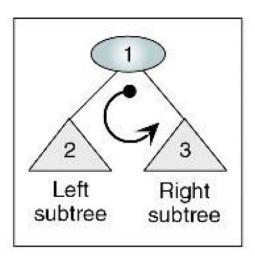




Preorder Traversal

Root -> Left subtree -> Right subtree

```
Algorithm preOrder(root):
    if (root is not null):
        process(root)
        preOrder(leftSubtree)
        preOrder(rightSubtree)
    end if
end preOrder
```





Preorder Traversal

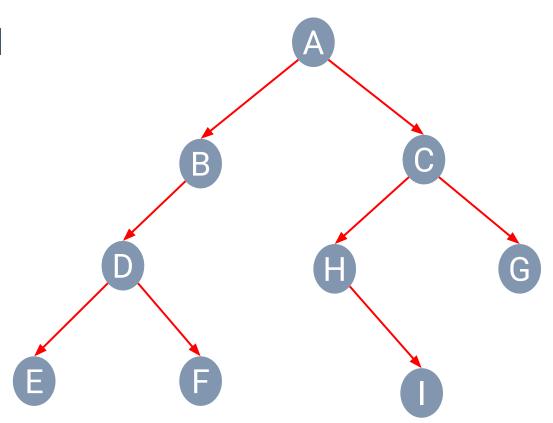
Preorder Traversal







Visited nodes





Preorder Traversal

Preorder Traversal

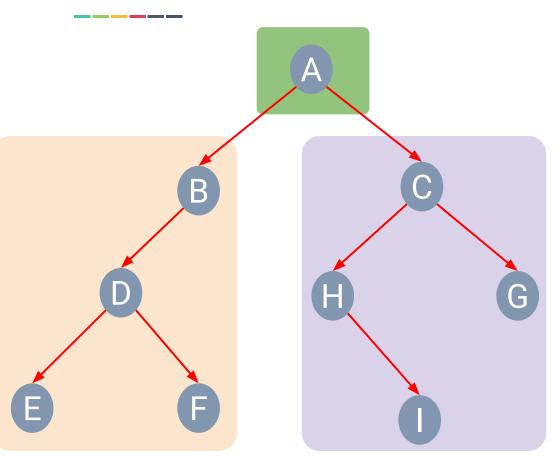






Visited nodes

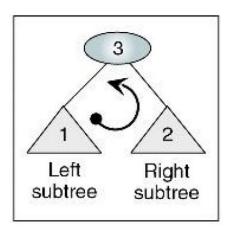






Left subtree -> Right subtree -> root

```
Algorithm postOrder(root):
    if (root is not null):
        postOrder(leftSubTree)
        postOrder(rightSubtree)
        process(root)
    end if
end postOrder
```

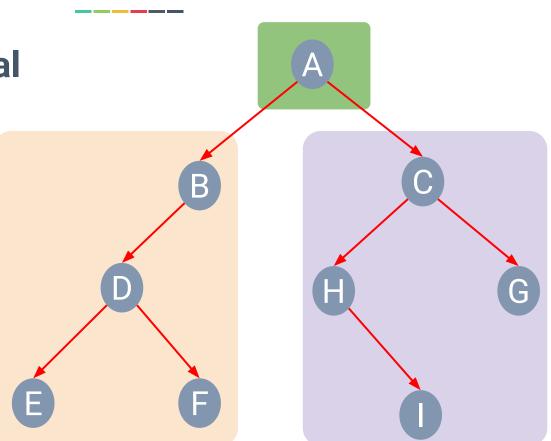




Postorder Traversal

L R N

Visited nodes

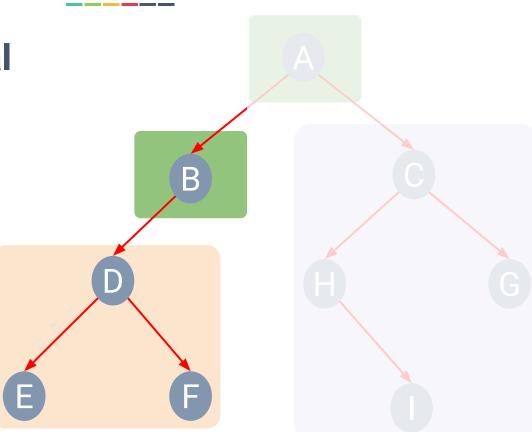




Postorder Traversal



Visited nodes





Postorder Traversal



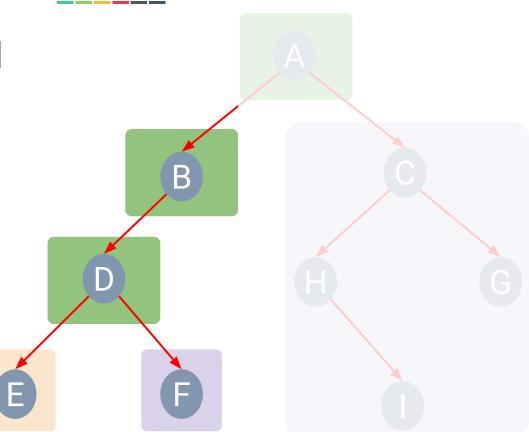










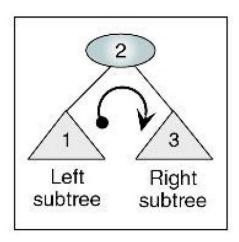




Inorder Traversal

Left subtree -> root -> Right subtree

```
Algorithm inOrder(root):
    if (root is not null):
        inOrder(leftSubTree)
        process(root)
        inOrder(rightSubtree)
    end if
end inOrder
```



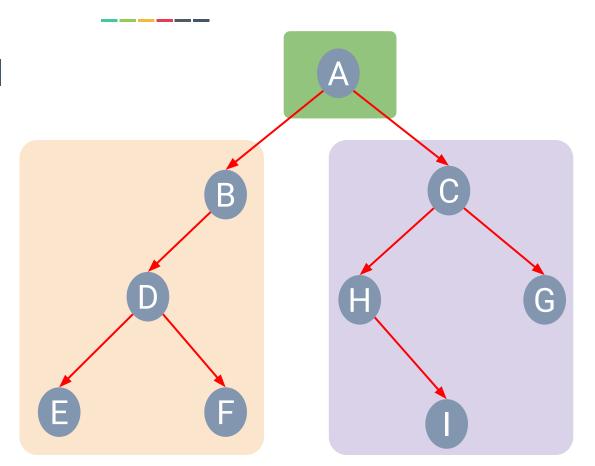


Inorder Traversal

Inorder Traversal



Visited nodes





Inorder Traversal

Inorder Traversal





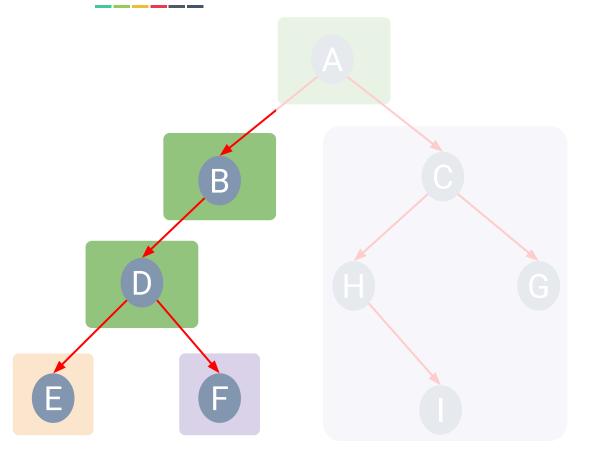


Visited nodes



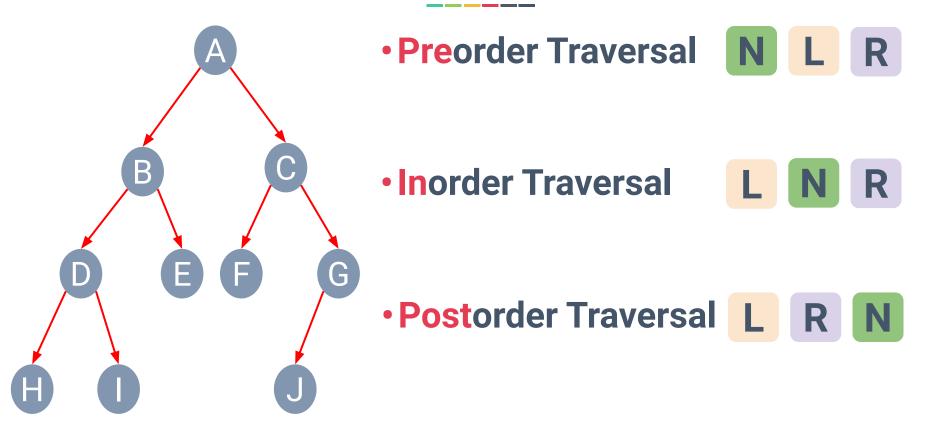






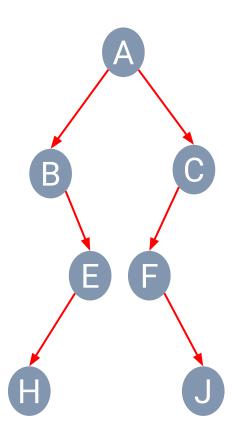


Depth-first Traversal Exercise 1





Depth-first Traversal Exercise 2



Preorder Traversal



. R

Inorder Traversal





R

Postorder Traversal

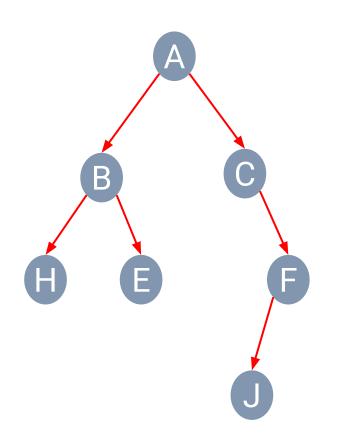








Depth-first Traversal Exercise 3



Preorder Traversal



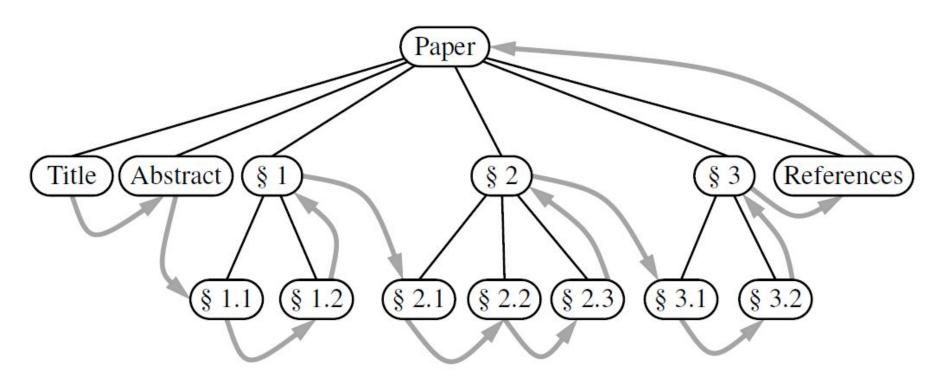
Inorder Traversal



Postorder Traversal L R



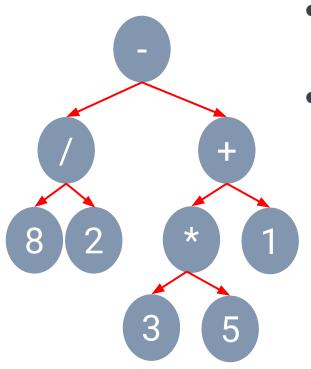
Depth-first Traversal Application



Preorder traversal of an ordered tree - Table of contents.



Depth-first Traversal Application



- Binary trees can be used to represent an arithmetic expression.
- The **inorder** traversal visits node in a consistent order with the expression.

Expression: (8/2) - ((3*5)+1)

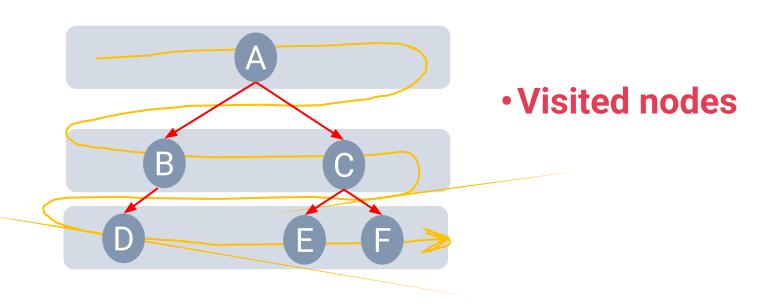




R



Breadth-first Traversal



• Visit all the nodes at depth *d* before moving to depth *d*+1.