Contract Use: Past, Present, and Future

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2019-09-18

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Who am I?

Software developer all century



- I have a purple house.
- First time presenting at CppCon
- First time presenting at a Conference

Who am I?

Software developer all century



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- First time presenting at CppCon
- First time presenting at a Conference (Be gently please)

Who am I?

- Bloomberg LP since 2017
- Joined BDE team in 2018
- Contract checking and deployment with BSLS_REVIEW
- WG21 participation to make contracts better
- SG21 participation with same goal

- Introduction
- 2 Basic Contracts
 - English Contracts
 - In Code Contracts
- Oping Stuff With Contracts
- 4 SG21
- Conclusion

Contracts are an agreement between two parties

Software contracts are an agreement between a library writer and client

Function contracts can be rendered in english

```
T* binsearch(T*begin, T*end, const T& val);

// Return a pointer to an element between the

// specified 'begin' and 'end' that is greater

// than or equal to the specified 'val', or 'end'

// if no such value exists. This function will

// perform no more than log(distance(begin,end))

// comparisons. The behavior is undefined

// unless '[begin,end)' is a contiguous sorted

// range.
```

Describe what a function will do

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Describe what behavior is not supported

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Undefined Behavior

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behavior for which this document imposes no requirements

N4830 - Working Draft, Standard for Programming Language C^{++}

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behavior for which a library contract provides no guarantees

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```

Preconditions

```
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Postconditions

```
T* binsearch(T*begin, T*end, const T& val);

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```

Essential Behavior

```
T* binsearch(T*begin, T*end, const T& val);

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```

Violating a contract is a bug

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Bugs are contract violations

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Possibly a contract no one wrote down

What can be checked?

• Parts of the english contract might be checkable with standard C++ expressions.

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- Parts might have readbale representations that cannot be implemented

What can be checked?

- Parts of the english contract might be checkable with standard C++ expressions.
- Parts might have readbale representations that cannot be implemented
- Parts might be statements beyond the scope of a single function execution

Some parts can be rendered with code

```
T* binsearch(T*begin, T*end, const T& val);

// Return a pointer to an element between the

// specified 'begin' and 'end' that is greater

// than or equal to the specified 'val', or 'end'

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```

Simple boolean predicates

```
begin != nullptr
end != nullptr
begin <= end
T* binsearch(T*begin, T*end, const T& val);
  // Return a pointer to an element between the
  // specified 'begin' and 'end' that is greater
 // than or equal to the specified 'val', or 'end'
 // if no such value exists. This function will
 // perform no more than log(distance(begin, end))
 // comparisons. The behavior is undefined
 // unless '[begin, end)' is a contiquous sorted
  // range.
```

Predicates about returned value

```
return val >= begin
return val <= end
return val == end *return val >= val
T* binsearch(T*begin, T*end, const T& val);
  // Return a pointer to an element between the
 // specified 'begin' and 'end' that is greater
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 // perform no more than log(distance(begin, end))
  // comparisons. The behavior is undefined
 // unless '[begin, end)' is a contiquous sorted
  // range.
```

Hard to check things

```
is_sorted(begin,end)
```

```
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// comparisons. The behavior is undefined

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```

Uncheckable things?

```
is_reachable_from(begin,end)
```

```
T* binsearch(T*begin, T*end, const T& val);

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```

Properties of repeated execution

?????????????????

```
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Prove software correctness

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- PROFIT

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- Compiler identifies all violated contracts
- Edge cases must be throught through
- All assumptions are captured in compiled code
- Mostly, if it compiles, it doesn't have bugs (contract violations)
- If there is a bug, contracts just need to be elaborated

No need for any checks

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- No need for any checks
- More knowledge for the compiler
 - __builtin_assume
 - Removing excess branches
 - Vectorization/SIMD instructions
- Smaller code size

• WARNING:

WARNING: MACROS INCOMING

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- Bloomberg has been doing this for 15 years

- WARNING: MACROS INCOMING
- How to leverage contracts without a language feature
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- See the BDE open source repostory for the real implementation
 - https://github.com/bloomberg/bde/blob/master/groups/bsl/bsls/bsls_assert.h
 - https://github.com/bloomberg/bde/blob/master/groups/bsl/bsls/bsls_review.h

What do you do if you can't prove a contract is being followed?

What do you do if you can't prove a contract is being followed?

Experiment

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Initial benefit of contracts in code

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#define ASSERT(X)

- Initial benefit of contracts in code
- Bloomberg specific naming

#define BSLS_ASSERT(X)

- Initial benefit of contracts in code
- Bloomberg specific naming
- Avoid code rot

```
#define BSLS_ASSERT(X) sizeof( (X)?true:false )
```

- Initial benefit of contracts in code
- Bloomberg specific naming
- Avoid code rot
- ... wish we had done that originally

```
#ifdef BSLS_ASSERT_VALIDATE_DISABLED_MACROS
#define BSLS_ASSERT(X) sizeof( (X)?true:false )
#else
#define BSLS_ASSERT(X)
#endif
```

- Initial benefit of contracts in code
- Bloomberg specific naming
- Avoid code rot
- ... wish we had done that originally
- ... or at least this to require a ;

```
#define BSLS_ASSERT(X) ((void)0)
```

- Initial benefit of contracts in code
- Bloomberg specific naming
- Avoid code rot
- ... wish we had done that originally
- ... or at least this to require a ;
- For simplicity

#define ASSERT(X)

• Identifying violations would be nice

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- The safest thing to do is stop immediately

```
#define ASSERT(X) if (!(X)) { std::abort(); }
```

- Identifying violations would be nice
- The safest thing to do is stop immediately
- Nice if ASSERT(X) needs a semicolon

```
#define ASSERT(X) do { if (!(X)) { std::abort(); } } while (false)
```

- Identifying violations would be nice
- The safest thing to do is stop immediately
- Nice if ASSERT(X) needs a semicolon
- For simplicity

```
#define ASSERT(X) if (!(X)) { std::abort(); }
```

• Checks of contracts are redundant if they're not broken

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- NDEBUG might be a way to control enablement

```
#ifdef NDEBUG
#define ASSERT(X)
#else
#define ASSERT(X) if (!(X)) { std::abort(); }
#endif
```

- Checks of contracts are redundant if they're not broken
- NDEBUG might be a way to control enablement
- This reminds me of something

```
#include <cassert>
#define ASSERT(X) assert(X)
```

- Checks of contracts are redundant if they're not broken
- NDEBUG might be a way to control enablement
- This reminds me of something
- Separating out controls from behavior helps

```
#define ASSERT_IMP(X) if (!(X)) { std::abort(); }
```

- Checks of contracts are redundant if they're not broken
- NDEBUG might be a way to control enablement
- This reminds me of something
- Separating out controls from behavior helps

```
\#define\ ASSERT\_IMP(X) \qquad \qquad if\ (!(X))\ \{\ std::abort();\ \} \ \#define\ ASSERT\_DISABLED\_IMP(X)
```

- Checks of contracts are redundant if they're not broken
- NDEBUG might be a way to control enablement
- This reminds me of something
- Separating out controls from behavior helps

```
#define ASSERT_IMP(X) if (!(X)) { std::abort(); }
#define ASSERT_DISABLED_IMP(X)

#ifdef ASSERT_LEVEL_ASSERT
#define ASSERT(X) ASSERT_IMP(X)
#else
#define ASSERT(X) ASSERT_DISABLED_IMP(X)
#endif
```

• Aborting with no information sucks

```
#define ASSERT_IMP(X) if (!(X)) {
    /*POOF*/;
    std::abort();
}
```

- Aborting with no information sucks
- Logging something helps

```
#define ASSERT_IMP(X) if (!(X)) {
    printf("ASSERTION FAILED!\n");
    std::abort();
}
```

- Aborting with no information sucks
- Logging something helps
- The preprocessor can give us more help

What about this guy?

A ERROR

IF YOU'RE SEEING THIS, THE CODE IS IN WHAT I THOUGHT WAS AN UNREACHABLE STATE.

I COULD GIVE YOU ADVICE FOR WHAT TO DO. BUT HONESTLY, WHY SHOULD YOU TRUST ME? I CLEARLY SCREWED THIS UP. I'M WRITING A MESSAGE THAT SHOULD NEVER APPEAR, YET I KNOW IT WILL PROBABLY APPEAR SOMEDAY.

ON A DEEP LEVEL, I KNOW I'M NOT UP TO THIS TASK. I'M SO SORRY.



NEVER WRITE ERROR MESSAGES TIRED.

- Aborting with no information sucks
- Logging something helps
- The preprocessor can give us more help
- Delegating to a pluggable function helps that

```
#define ASSERT_IMP(X) if (!(X)) {
    bb::Assert::invoke_violation_handler(__FILE__, __LINE__, #X);
    std::abort();
}
```

- Aborting with no information sucks
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```
#define ASSERT_IMP(X) if (!(X)) {
    bb::assert_violation violation(__FILE__, __LINE__, #X); \
    bb:Assert::invoke_violation_handler(violation); \
    std::abort();
}
```

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- Leave all behavior up to the violation handler

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#define ASSERT_IMP(X) if (!(X)) {
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```
void xkcd::violationHandler(const bb::assert_violation &violation) {
  printf("Error\n");
  printf("If you're seeing this, the code is in what\n");
  printf("I thought was an unreachable state.");
  //...
}
```

- Aborting with no information sucks
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- The preprocessor can give us more help
- Delegating to a pluggable function helps that
- Leave all behavior up to the violation handler

```
int main() {
  bb::Assert::setViolationHandler(&xkcd::violationHandler);
  //..
}
```

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```
std::abort()
```

```
• while (true) {std::this_thread::sleep_for(std::chrono::years(1));}
```

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• std::abort()
```

```
• while (true) {std::this_thread::sleep_for(std::chrono::years(1));}
```

• throw std:::exception("Oops?");

- The violation handler can notify in different ways
 - Custom logging frameworks
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 - std::abort()
 - while (true) {std::this_thread::sleep_for(std::chrono::years(1));}
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 - Custom logging frameworks
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- ... do different things
 - std::abort()
 - while (true) {std::this_thread::sleep_for(std::chrono::years(1));}
 - throw std:::exception("Oops?");
- ... or try to recover?
- main gets to decide

• Checks use state already in cache, are often very fast

```
T* binsearch(T*begin, T*end, const T& val) {
   ASSERT(begin);
   ASSERT(end);
   ASSERT(begin < end)7
   //..
}</pre>
```

- Checks use state already in cache, are often very fast
- Algorithmic complexity can still ruin that

```
T* binsearch(T*begin, T*end, const T& val) {
   ASSERT(is_sorted_range(begin,end));
   //..
}
```

- Checks use state already in cache, are often very fast
- Algorithmic complexity can still ruin that
- 3 levels of complexity

```
#define ASSERT_OPT(X) ...
#define ASSERT(X) ...
#define ASSERT_SAFE(X) ...
```

- Checks use state already in cache, are often very fast
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- 3 levels of complexity
- ... 2 levels probably sufficient

```
[[ assert default : X ]];
[[ assert audit : X ]];
```

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- 3 levels of complexity
- ... 2 levels probably sufficient
- Linear scale of enablement.

```
#if defined(ASSERT_LEVEL_NONE) ? 1 : 0 \
    + defined(ASSERT_LEVEL_OPT) ? 1 : 0 \
    + defined(ASSERT_LEVEL_ASSERT) ? 1 : 0 \
    + defined(ASSERT_LEVEL_SAFE) ? 1 : 0 \
    > 1
#error Multiple ASSERT_LEVEL macros defined
#endif
```

- Checks use state already in cache, are often very fast
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- ... 2 levels probably sufficient
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```
#if !defined(ASSERT_LEVEL_NONE) \
&& !defined(ASSERT_LEVEL_OPT) \
&& !defined(ASSERT_LEVEL_ASSERT) \
&& !defined(ASSERT_LEVEL_SAFE)
#define ASSERT_LEVEL_ASSERT
#endif
```

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```
#if defined(ASSERT_LEVEL_NONE)
#define ASSERT_OPT(X) ASSERT_DISABLED_IMP(X)
#define ASSERT(X) ASSERT_DISABLED_IMP(X)
#define ASSERT_SAFE(X) ASSERT_DISABLED_IMP(X)
//..
```

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```
//..
#elif defined(ASSERT_LEVEL_OPT)
#define ASSERT_OPT(X) ASSERT_IMP(X)
#define ASSERT(X) ASSERT_DISABLED_IMP(X)
#define ASSERT_SAFE(X) ASSERT_DISABLED_IMP(X)
//..
```

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//..
#elif defined(ASSERT_LEVEL_ASSERT)
#define ASSERT_OPT(X) ASSERT_IMP(X)
#define ASSERT(X) ASSERT_IMP(X)
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//..
```

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Checking is slow!

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```
#if defined(ASSERT_LEVEL_OPT) \
  defined(ASSERT_LEVEL_ASSERT) \
  defined(ASSERT_LEVEL_SAFE)

#define ASSERT_OPT(X) ASSERT_IMP(X)

#else
  // defined(ASSERT_LEVEL_NONE)

#define ASSERT_OPT(X) ASSERT_DISABLED_IMP(X)

#endif
```

Checking is slow!

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```
#if defined(ASSERT_LEVEL_ASSERT) \
  defined(ASSERT_LEVEL_SAFE)

#define ASSERT(X) ASSERT_IMP(X)

#else
  // defined(ASSERT_LEVEL_OPT)
  // defined(ASSERT_LEVEL_NONE)

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#endif
```

Checking is slow!

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- Bloomberg 2005-2018

```
#if defined(ASSERT_LEVEL_SAFE)
#define ASSERT_SAFE(X) ASSERT_IMP(X)
#else
// defined(ASSERT_LEVEL_OPT)
// defined(ASSERT_LEVEL_ASSERT)
// defined(ASSERT_LEVEL_NONE)
#define ASSERT_SAFE(X) ASSERT_DISABLED_IMP(X)
#endif
```

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- Changing is hard

What we did

- What we did
 - Developement ASSERT_LEVEL_ASSERT

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 - Beta testing ASSERT_LEVEL_ASSERT_OPT
 - Production ASSERT_LEVEL_ASSERT_OPT
- What we wanted

- What we did
 - Developement ASSERT_LEVEL_ASSERT
 - Unit tests ASSERT_LEVEL_ASSERT hopefully
 - Beta testing ASSERT_LEVEL_ASSERT_OPT
 - Production ASSERT_LEVEL_ASSERT_OPT
- What we wanted
 - Developement ASSERT_LEVEL_ASSERT or ASSERT_LEVEL_SAFE

- What we did
 - Developement ASSERT_LEVEL_ASSERT
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- Adding more assertions
 - ~bsl::string() { ASSERT(m data[m size] == 0); }
 - Time ABI change
- Changing levels of assertions
 - SAFE to ASSERT
 - ASSERT to OPT
- Changing deployed assertion levels
- Everyone will need to do this in 202x!
 - Using language contracts when they come will be a case of adding new assertions to existing code.

Mis-Step #1

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 - Blanket continuation unsafe

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 - Tracking failure counts
 - Alternate logging
- Still unsuccessful
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 - No way to indicate in code that a check is "new"
 - Rarely used, minimal progress

• BSLS_REVIEW (2018)

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Parallel structure to BSLS_ASSERT

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- Lifecycle BSLS_ASSERT_SAFE to BSLS_REVIEW to BSLS_ASSERT
- Alternately, <nothing> to BSLS_REVIEW_? to BSLS_ASSERT_?

Initially a copy of ASSERT

```
#define REVIEW_IMP(X) if (!(X)) {
    bb::assert_violation violation(__FILE__, __LINE__, #X);
    bb::Review::invoke_violation_handler(violation);
}
```

- Initially a copy of ASSERT
- Number of failures is important

```
#define REVIEW_IMP(X) if (!(X)) {
    static std::atomic<int> count;
    bb::review_violation violation(_FILE__, __LINE__, ++count, #X);\
    bb::Review::invoke_violation_handler(violation);
}
```

- Initially a copy of ASSERT
- Number of failures is important
- Default violation handler logs only

- Initially a copy of ASSERT
- Number of failures is important
- Default violation handler logs only
- With expeonential backoff

```
void Review::default violation handler(
                              const bb::review violation &violation)
  int count = violation.count();
  if (0 == (count & (count-1))) {
    // Log a message, with contents of violation
    // Log a stack trace
  // Return
```

Mutually exlusive

```
#if defined(REVIEW_LEVEL_NONE) ? 1 : 0 \
    + defined(REVIEW_LEVEL_OPT) ? 1 : 0 \
    + defined(REVIEW_LEVEL_REVIEW) ? 1 : 0 \
    + defined(REVIEW_LEVEL_SAFE) ? 1 : 0 \
    > 1
#error Multiple REVIEW_LEVEL macros defined
#endif
```

- Mutually exlusive
- Default to assert level

```
#if defined(ASSERT_LEVEL_NONE)
#define REVIEW_LEVEL_NONE
#elif defined(ASSERT LEVEL OPT)
#define REVIEW LEVEL OPT
#elif defined(ASSERT LEVEL ASSERT)
#define REVIEW LEVEL REVIEW
#elif defined(ASSERT LEVEL SAFE)
#define REVIEW LEVEL SAFE
#else
#define REVIEW LEVEL REVIEW
#endif
```

- Mutually exlusive
- Default to assert level (In reality copies assert logic)

```
#if defined(ASSERT_LEVEL_NONE)
#define REVIEW_LEVEL_NONE
#elif defined(ASSERT LEVEL OPT)
#define REVIEW LEVEL OPT
#elif defined(ASSERT LEVEL ASSERT)
#define REVIEW LEVEL REVIEW
#elif defined(ASSERT LEVEL SAFE)
#define REVIEW LEVEL SAFE
#else
#define REVIEW LEVEL REVIEW
#endif
```

- Mutually exlusive
- Default to assert level
- Controls just like ASSERT

```
#if defined(REVIEW_LEVEL_NONE)
#define REVIEW_OPT(X) REVIEW_DISABLED_IMP(X)
#define REVIEW(X) REVIEW_DISABLED_IMP(X)
#define REVIEW_SAFE(X) REVIEW_DISABLED_IMP(X)
//..
```

- Mutually exlusive
- Default to assert level
- Controls just like ASSERT

```
//..
#elif defined(REVIEW_LEVEL_OPT)
#define REVIEW_OPT(X) REVIEW_IMP(X)
#define REVIEW(X) REVIEW_DISABLED_IMP(X)
#define REVIEW_SAFE(X) REVIEW_DISABLED_IMP(X)
//..
```

- Mutually exlusive
- Default to assert level
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```
//..
#elif defined(REVIEW_LEVEL_REVIEW)
#define REVIEW_OPT(X) REVIEW_IMP(X)
#define REVIEW(X) REVIEW_IMP(X)
#define REVIEW_SAFE(X) REVIEW_DISABLED_IMP(X)
//..
```

- Mutually exlusive
- Default to assert level
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```
//..
#elif defined(REVIEW_LEVEL_SAFE)
#define REVIEW_OPT(X) REVIEW_IMP(X)
#define REVIEW(X) REVIEW_IMP(X)
#define REVIEW_SAFE(X) REVIEW_IMP(X)
#endif
```

- Mutually exlusive
- Default to assert level
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```
#if defined(REVIEW_LEVEL_OPT)
  defined(REVIEW_LEVEL_REVIEW) \
  defined(REVIEW_LEVEL_SAFE)

#define REVIEW_OPT(X) REVIEW_IMP(X)

#else
  // defined(REVIEW_LEVEL_NONE)

#define REVIEW_OPT(X) REVIEW_DISABLED_IMP(X)

#endif
```

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- Default to assert level
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```
#if defined(REVIEW_LEVEL_REVIEW) \
  defined(REVIEW_LEVEL_SAFE)

#define REVIEW(X) REVIEW_IMP(X)

#else
  // defined(REVIEW_LEVEL_OPT)
  // defined(REVIEW_LEVEL_NONE)

#define REVIEW(X) REVIEW_DISABLED_IMP(X)

#endif
```

- Mutually exlusive
- Default to assert level
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```
#if defined(REVIEW_LEVEL_SAFE)
#define REVIEW_SAFE(X) REVIEW_IMP(X)
#else
// defined(REVIEW_LEVEL_OPT)
// defined(REVIEW_LEVEL_REVIEW)
// defined(REVIEW_LEVEL_NONE)
#define REVIEW_SAFE(X) REVIEW_DISABLED_IMP(X)
#endif
```

ASSERT and REVIEW interaction

• Changing build levels requires reviewing all asserts at the target level

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- BSLS_ASSERT again

```
#if defined(BSLS_ASSERT_LEVEL_ASSERT) \
  defined(BSLS_ASSERT_LEVEL_SAFE)
#define BSLS_ASSERT(X) ASSERT_IMP(X)
#else
#define BSLS_ASSERT(X)
#endif
```

ASSERT and REVIEW interaction

- Changing build levels requires reviewing all asserts at the target level
- BSLS_ASSERT again

```
#if defined(BSLS_ASSERT_LEVEL_ASSERT) \
    defined(BSLS_ASSERT_LEVEL_SAFE)
#define BSLS_ASSERT(X) ASSERT_IMP(X)
#elif defined(BSLS_REVIEW_LEVEL_REVIEW) \
    defined(BSLS_REVIEW_LEVEL_SAFE)
#define BSLS_ASSERT(X) REVIEW_IMP(X)
#else
#define BSLS_ASSERT(X) ASSERT_DISABLED_IMP(X)
#endif
```

ASSERT and REVIEW interaction

- Changing build levels requires reviewing all asserts at the target level
- BSLS_ASSERT again
- Same for BSLS_ASSERT_OPT and BSLS_ASSERT_SAFE

```
#if defined(BSLS_ASSERT_LEVEL OPT)
  defined(BSLS_ASSERT_LEVEL_ASSERT) \
  defined(BSLS ASSERT LEVEL SAFE)
#define BSLS ASSERT OPT(X) ASSERT IMP(X)
#elif defined(BSLS REVIEW LEVEL OPT)
  defined(BSLS REVIEW LEVEL REVIEW) \
  defined(BSLS REVIEW LEVEL SAFE)
#define BSLS ASSERT OPT(X) REVIEW IMP(X)
#else
#define BSLS ASSERT OPT(X) ASSERT DISABLED IMP(X)
#endif
```

ASSERT and REVIEW interaction

- Changing build levels requires reviewing all asserts at the target level
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```
#if defined(BSLS_ASSERT_LEVEL_SAFE)
#define BSLS_ASSERT_SAFE(X) ASSERT_IMP(X)
#elif defined(BSLS_REVIEW_LEVEL_SAFE)
#define BSLS_ASSERT_SAFE(X) REVIEW_IMP(X)
#else
#define BSLS_ASSERT_SAFE(X) ASSERT_DISABLED_IMP(X)
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 - Can conrol from code
 - Can control at build time

What do you do if you can't prove a contract is being followed?

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Believe

- Introduction
- 2 Basic Contracts
- Oping Stuff With Contracts
 - A Dream
 - Less Bugs
 - Deploying it
 - Faster Code
- 4 SG21
- Conclusion

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- If you believe the contract is being followed, __builtin_assume can give you the benefit without the costt
- The risk is the strength of your belief

• Let's add another choice for mapping the BSLS_ASSERT macros

```
#define BSLS_ASSERT_ASSUME(X) if (!(X)) { std::unreachable(); }
```

- Let's add another choice for mapping the BSLS_ASSERT macros
- Lots of ways to impelement, different tradeoffs and portability

```
#define BSLS\_ASSERT\_ASSUME(X) if (!(X)) { std::unreachable(); } #define BSLS\_ASSERT\_ASSUME(X) __builtin_assume(X) #define BSLS\_ASSERT\_ASSUME(X) if (!(X)) { int *p = nullptr; *p = 17; }
```

- Let's add another choice for mapping the BSLS_ASSERT macros
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- This almost made it to the standard

```
#define BSLS_ASSERT_ASSUME(X) [[ assert assume : X ]]
```

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- This almost made it to the standard
- Coming to BDE with an extended BSLS_ASSERT_LEVEL scale

```
//..
#elif defined(ASSERT_LEVEL_ASSUME_OPT)
#define ASSERT_OPT(X) ASSERT_ASSUME(X)
#define ASSERT(X) ASSERT_DISABLED_IMP(X)
#define ASSERT_SAFE(X) ASSERT_DISABLED_IMP(X)
//..
```

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 - P1607 proposed two options remove all but a nicer c assert, or remove all but add in literal semantics
- P1607's two options were the only consensus reached that day by EWG

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- On Wednesday, July 17th in Cologne P1823 was proposed and accepted by a massive margin
- On Saturday, July 20th, P1823 was ratified and SG21 was announced to pursue contracts again

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 - As a C++ API Developer, in order to Maintain a class hierarchy, I want to Ensure overriding methods have same or wider preconditions

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 - As a C++ API Developer, in order to Maintain a class hierarchy, I want to Ensure overriding methods have same or wider preconditions
- Some require vastly more than was possible before:

- Use cases were gathered from all SG21 participants and edited for clarity
- All of the form "As an X in order to Y I need to Z".
- This hopefully be the first wg21 paper published "by" SG21.
- 29 different classes of users with 196 different use cases
- Use cases range from very general to very specific:
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- Some require vastly more than was possible before:
 - As a C++ API Developer In Order to Enforce contracts in async code I want to Express contracts on callbacks such as std::function, function pointers, or references to functions, lambdas, or function objects

Prioritization

- SG21 members have been asked to rate each use case on behalf of whatever users they feel they best represent:
 - Not important
 - Nice to have
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- Expect these results to be ready to analyze and discuss by Belfast.

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Conclusion

- Bloomberg's BSLS_ASSERT and BSLS_REVIEW provide a rich set of contract enforcement utility. Grab the open source BDE to play with it today
- The needs of that facility will hopefully be met by language level contracts in the future, SG21 is working hard to see that happen