

6

Uncertainty & Utilities

Idea: Uncertain outcomes controlled by chance not an adversary!

* Expectimax Search

| | | |
|----------|----------|------------|
| Δ | ∇ | \bigcirc |
| Max | Min | Chance |

⇒ Value should now reflect average-case (expectimax) outcome, not worst-case (minimax) outcomes.

⇒ Expectimax Search: Compute the average score under optimal play.

- Max nodes as in minimax search
- Chance nodes are like min nodes but the outcome is uncertain.
- Calculate their expected utilities
(i.e. take weighted average (expectation) of children)

★ Expectimax Pseudo code

```
def value(state):
```

```
    if the state is a terminal state
        ↳ return the state's utility
```

```
    if the next agent is MAX:
```

```
        ↳ return max-value(state)
```

```
    if the next agent is EXP EXP:
```

```
        ↳ return exp-value(state)
```

```
def max-value(state)
```

```
    initialize  $V = -\infty$ 
```

```
    for each successor of state
```

```
         $V = \max(V, \text{value}(\text{successor}))$ 
```

```
    return V
```

```
def exp-value(state)
```

```
    initialize  $V = 0$ 
```

```
    for each successor of state:
```

```
         $P = \text{Probability}(\text{successor})$ 
```

```
         $V \pm P * \text{value}(\text{successor})$ 
```

```
    return V
```

★ Reminder: Probabilities

Random variable

↳ Represents an event whose outcome is unknown

Probability distribution

↳ Assignments of weights to outcomes.

⇒ The expected value of a function of a random variable is the average, weighted by the probability distribution over outcomes.

* Utility

↳ Function from outcomes (states of the world) to real number that describes an agent performance

⇒ An agent must preferences among:

- prizes: A, B etc
- Lotteries: Situation with uncertain prizes

$$L = [p, A; (1-p), B]$$

⇒ Notation:

- Preference: $A \succ B$
- Indifference: $A \sim B$

* Rational Preferences

⇒ We want some constraints on preferences before we call them rational, such as:

- Axiom of Transitivity:

$$(A \succ B) \wedge (B \succ C) \Rightarrow (A \succ C)$$

- Orderability

~~$$(A \succ B) \wedge (B \succ C) \Rightarrow (A \succ C)$$~~

$$(A \succ B) \vee (B \succ A) \vee (A \sim B)$$

- Continuity

$$A \succ B \succ C \Rightarrow \exists p [pA; (1-p), C] \sim B$$

- Substitutability

$$A \sim B \Rightarrow [p, A; (1-p), C] \sim [p, B; (1-p), C]$$

- Monotonicity

$$A \succ B \Rightarrow (p \geq q \Leftrightarrow [p, A; (1-p), B] \succ [q, A; (1-q), B])$$

Theorem [Ransey, 1931; Von Neumann & Morgenstern, 1944]

Given any preference satisfying these constraints,
there exists a real-valued function U such that

$$U(A) \geq U(B) \Leftrightarrow A \succeq B$$

$$U([p_1, s_1; \dots; p_n, s_n]) = \sum_i p_i U(s_i)$$

\Rightarrow Maximum expected utility (MEU) principle:

\hookrightarrow Choose the action that maximizes
expected utility.

