

### **DTS202TC Foundation of Parallel Computing**

Lecture 4 OpenMP

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### **Group Assessment 1 realeased**



- 20%, due October 10, 23:59pm
- 60 group marks + 40 individual marks
- Work in groups, discuss with your peers, every member should make contributions.
- One group member submit the group work.
- Peer review portal will be opened later.

### Week 5 CUDA



- 6 hours face-to-face lectures + lab (live streaming provided in case you are not in Suzhou)
- Mon. 11<sup>th</sup> Oct. 9-11am
- Tue. 12<sup>th</sup> Oct. 4-6pm
- Thurs. 14<sup>th</sup> Oct. 11am-1pm
- No recording due to the copyright

# This week's topic



- OpenMP Basics
  - Our First OpenMP Program
  - Fundamental Concepts and Library Functions
- The Trapezoidal Rule
- Scope of Variables
- Task Parallelism
  - The Reduction Clause
  - The parallel for Directive

### **OpenMP**



- An API for shared-memory parallel programming.
- MP = multiprocessing
- Designed for systems in which each thread or process can potentially have access to all available memory.

### **OpenMP vs. Pthreads**



- Pthreads requires that the programmer explicitly specify the behavior of each thread. OpenMP allows the compiler and runtime system to determine some of the details of thread behavior.
- Any Pthreads program can be used with any C compiler, provided the system has a Pthreads library. OpenMP requires compiler support for some operations, and hence it's entirely possible that you may run across a C compiler that can't compile OpenMP programs into parallel programs.
- Pthreads is lower level. Cost: Specify every detail of the behavior of each thread. OpenMP can be simpler to code some parallel behaviors. Cost: Some low-level thread interactions can be more difficult to program.

### **Pragma**



- Special preprocessor instructions.
- Typically added to a system to allow behaviors that aren't part of the basic C specification.
- Compilers that don't support the pragmas ignore them.

#pragma

### "Hello, World" Using OpenMP

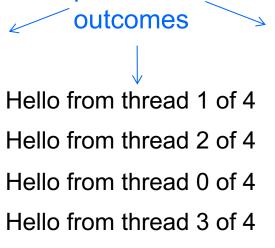


```
#include < stdio.h>
#include < stdlib.h>
#include <omp.h>
void Hello(void); /* Thread function */
int main(int argc, char* argv[]) {
   /* Get number of threads from command line */
   int thread count = strtol(argv[1], NULL, 10);
  pragma omp parallel num_threads(thread_count)
   Hello();
   return 0;
  /* main */
void Hello(void) {
   int my_rank = omp_get_thread_num();
   int thread_count = omp_get_num_threads();
   printf("Hello from thread %d of %d\n", my_rank, thread_count);
  /* Hello */
```

# **Compiling and Running**



Hello from thread 0 of 4
Hello from thread 1 of 4
Hello from thread 2 of 4
Hello from thread 3 of 4



possible

Hello from thread 3 of 4
Hello from thread 1 of 4
Hello from thread 2 of 4
Hello from thread 0 of 4

### **Outline**



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# **OpenMP Syntax**



- Most of the constructs in OpenMP are compiler directives
- -#pragma omp directive [clause list]
- A parallel directive: # pragma omp parallel
  - The number of threads that run the following structured block of code is determined by the run-time system.
- Function prototypes and types in the file
  - #include <omp.h>

### Note



- There may be system-defined limitations on the number of threads that a program can start.
- The OpenMP standard doesn't guarantee that this will actually start thread\_count threads.
- Most current systems can start hundreds or even thousands of threads.
- Unless we're trying to start a lot of threads, we will almost always get the desired number of threads.

# **Some Terminology**



 In OpenMP parlance, the collection of threads executing the parallel block — the original thread and the new threads — is called a team, the original thread is called the master, and the additional threads are called slaves.

### In Case the Compiler doesn't Support OpenMP



```
# include <omp.h>

#ifdef _OPENMP

# include <omp.h>
#endif
```



```
# ifdef OPENMP
 int my rank = omp get thread num ();
 int thread count = omp get num threads ();
# else
 int my rank = 0;
 int thread count = 1;
# endif
```

# **OpenMP Library Functions**



Control the number of threads and Processors:

```
#include <omp.h>
void omp_set_num_threads (int num_threads);
int omp_get_num_threads ();
int omp_get_max_threads ();
int omp_get_thread_num ();
int omp_get_num_procs ();
int omp in parallel ();
```

Set and monitor thread creation:

```
#include <omp.h>
void omp_set_dynamic(int dynamic_threads);
int omp_get_dynamic ();
void omp_set_nested (int nested);
int omp_get_nested ();
```

Mutex:

```
#include <omp.h>
void omp_init_lock(omp_lock_t *lock);
void omp_destroy_lock(omp_lock_t *lock);
void omp_set_lock(omp_lock_t *lock);
void omp_unset_lock(omp_lock_t *lock);
int omp_test_lock(omp_lock_t *lock);
```

 OpenMP also supports nested lock, which has similar semantics with simple lock.

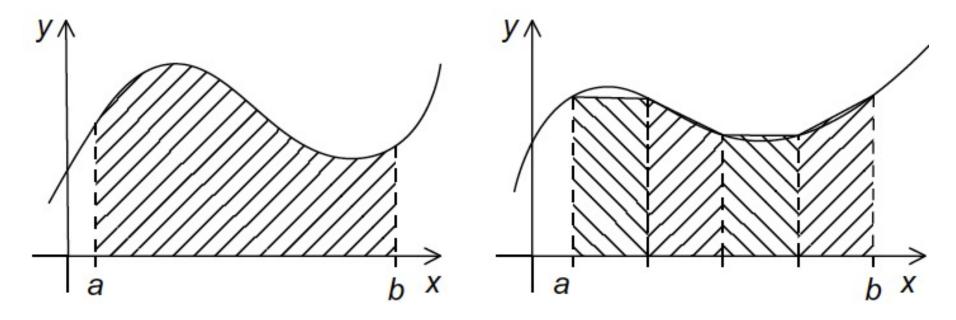
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# The Trapezoidal Rule





If each subinterval has the same length h and if we define h=(b-a)/n,  $x_i=a+ih$ , i=0,1,2,...,n, then our approximation will be

$$h[f(x_0)/2 + f(x_1) + f(x_2) + \dots + f(x_{n-1}) + f(x_n)/2]$$

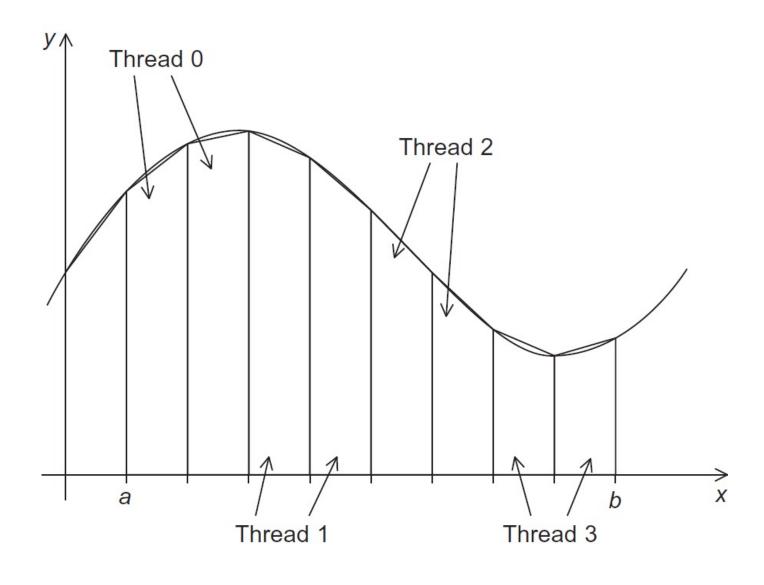
# **Serial Algorithm**



```
/* Input: a, b, n */
h = (b-a)/n;
approx = (f(a) + f(b))/2.0;
for (i = 1; i <= n-1; i++) {
    x_i = a + i*h;
    approx += f(x_i);
}
approx = h*approx;</pre>
```

### **Assignment of Trapezoids to Threads**





Time	Thread 0	Thread 1
0	global_result = 0 to register	finish my_result
1	my_result = 1 to register	global_result = 0 to register
2	add my_result to global_result	my_result = 2 to register
3	<pre>store global_result = 1</pre>	add my_result to global_result
4		<pre>store global_result = 2</pre>

Unpredictable results when two (or more)
threads attempt to simultaneously execute:
global\_result += my\_result;(critical section)

#### Recall:

Race Condition, Critical Section

### **Mutual Exclusion**



```
Critical Directive:
```

```
# pragma omp critical
global_result += my_result;
```

only one thread can execute
the following structured block at a time

### First OpenMP Trapezoidal Rule Program(1)



```
#include < stdio.h>
#include < stdlib.h>
#include <omp.h>
void Trap(double a, double b, int n, double* global_result_p);
int main(int argc, char* argv[]) {
   double global_result = 0.0; /* Store result in global_result */
  double a, b;
                /* Left and right endpoints
                                                                */
                               /* Total number of trapezoids
   int
         n:
                                                                */
   int thread count;
  thread count = strtol(argv[1], NULL, 10);
  printf("Enter a, b, and n\n");
  scanf("%lf %lf %d", &a, &b, &n);
 pragma omp parallel num_threads(thread_count)
  Trap(a, b, n, &global_result);
  printf("With n = %d trapezoids, our estimate\n", n);
  printf("of the integral from %f to %f = %.14e\n",
     a, b, global_result);
  return 0:
  /* main */
```

### First OpenMP Trapezoidal Rule Program(2)



```
void Trap(double a, double b, int n, double* global_result_p) {
   double h, x, my_result;
   double local a, local b;
   int i, local n;
   int my rank = omp get thread num();
   int thread_count = omp_get_num_threads();
   h = (b-a)/n;
   local n = n/thread count;
   local_a = a + my_rank*local_n*h;
   local b = local a + local n*h;
   my result = (f(local a) + f(local b))/2.0;
   for (i = 1; i \le local_n - 1; i++)
     x = local a + i*h;
    my result += f(x);
   my result = my result *h;
  pragma omp critical
   *qlobal result p += my result;
} /* Trap */
```

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# **Definition of Scope**



- •In serial programming, the scope of a variable consists of those parts of a program in which the variable can be used.
- •In OpenMP, the scope of a variable refers to the set of threads that can access the variable in a parallel block.

# Scope in OpenMP



- A variable that can be accessed by all the threads in the team has shared scope.
- A variable that can only be accessed by a single thread has private scope.
- The default scope for variables declared before a parallel block is shared.

### **Example: Hello, World**



```
#include < stdio.h>
#include < stdlib.h>
#include <omp.h>
void Hello(void); /* Thread function */
int main(int argc, char* argv[]) {
  /* Get number of threads from command line */
   int thread count = strtol(argv[1], NULL, 10);
  pragma omp parallel num_threads(thread_count)
  Hello();
   return 0;
                      Private Scope
  /* main */
void Hello(void)
   int my_rank = omp_get_thread_num();
   int thread_count = omp_get_num_threads();
  printf("Hello from thread %d of %d\n", my_rank, thread_count);
  /* Hello */
```

# **Example: Trapezoidal Rule(1)**



```
#include < stdio.h>
#include < stdlib.h>
#include <omp.h>
void Trap(double a, double b, int n, double* global_result_p);
int main(int argc, char* argv[]) {
   double | global_result = 0.0; /* Store result in global_result */
   double
                                 /* Left and right endpoints
          a. b:
                                 /* Total number of trapezoids
   int
          n:
                            Shared Scope
   int
           thread count;
   thread_count = strtol(argv[1], NULL, 10);
   printf("Enter a, b, and n\n");
   scanf("%lf %lf %d", &a, &b, &n);
  pragma omp parallel num_threads(thread_count)
   Trap(a, b, n, &global_result);
   printf("With n = %d trapezoids, our estimate\n", n);
   printf("of the integral from %f to %f = %.14e\n",
      a, b, global result);
   return 0;
   /* main */
```

# Example: Trapezoidal Rule(2)



#### **Private Scope**

```
void Trap(double a, double b, int n, double* global_result_p) {
  _double h, x, my_result;
   double local a, local b;
   int i, local n;
   int my rank = omp get thread num();
  int thread_count = omp_get_num_threads();
   h = (b-a)/n;
   local n = n/thread count;
   local_a = a + my_rank*local_n*h;
   local b = local a + local n*h;
   my_result = (f(local_a) + f(local_b))/2.0;
   for (i = 1; i \le local_n - 1; i++)
    x = local a + i*h;
    my result += f(x);
   my_result = my_result*h;
  pragma omp critical
   *global_result_p += my_result;
} /* Trap */
                   *Shared Scope
```

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In the OpenMP program, this more complex version is used to get global\_result by adding each thread's local calculation.

```
void Trap(double a, double b, int n, double* global_result_p);
```

Although we'd prefer this for a serial implementation.

```
double Trap(double a, double b, int n);

global_result = Trap(a, b, n);
```

### **Analysis of Trapezoidal Rule Program(2)**



For the pointer version, we need to add each thread's local calculation to get global\_result. If we use this, there's no critical section!

```
double Local_trap(double a, double b, int n);
```

If we fix it like this, we force the threads to execute sequentially.

```
global_result = 0.0;
# pragma omp parallel num_threads(thread_count)
{
# pragma omp critical
    global_result += Local_trap(double a, double b, int n);
}
```

### Analysis of Trapezoidal Rule Program(3)



We can avoid this problem by declaring a private variable inside the parallel block and moving the critical section after the function call.

```
global_result = 0.0;

pragma omp parallel num_threads(thread_count)
{
    double my_result = 0.0; /* private */
    my_result += Local_trap(double a, double b, int n);
    pragma omp critical
    global_result += my_result;
}
```

# Syntax of Reduction Clause



```
reduction(<operator>: <variable list>)
+, *, -, &, |, ^, &&, ||
```

- A reduction operator is a binary operation (such as addition or multiplication).
- A reduction is a computation that repeatedly applies the same reduction operator to a sequence of operands in order to get a single result.
- All of the intermediate results of the operation should be stored in the same variable: the reduction variable.

### Reduction



# A reduction clause can be added to a parallel directive.

```
global_result = 0.0;
# pragma omp parallel num_threads(thread_count) \
    reduction(+: global_result)
global_result += Local_trap(double a, double b, int n);
```

- When a variable is included in a reduction clause, the variable itself is shared. However, a private variable is created for each thread in the team.
- In the parallel block each time a thread executes a statement involving the variable, it uses the private variable. When the parallel block ends, the values in the private variables are combined into the shared variable.

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### **Parallel for**



- Forks a team of threads to execute the following structured block.
- However, the structured block following the parallel for directive must be a for loop.
- With the parallel for directive, the system parallelizes the for loop by dividing the iterations of the loop among the threads.
- With just a parallel directive, in general, the work must be divided among the threads by the threads themselves.

### **Example**

```
h = (b-a)/n;
approx = (f(a) + f(b))/2.0;
for (i = 1; i \le n-1; i++)
   approx += f(a + i*h);
approx = h*approx;
        h = (b-a)/n;
         approx = (f(a) + f(b))/2.0;
     # pragma omp parallel for num_threads(thread_count) \
            reduction(+: approx)
         for (i = 1; i \le n-1; i++)
            approx += f(a + i*h);
         approx = h*approx;
```

## Scope



- The default scope for all variables in a parallel directive is shared.
- •In a loop that is parallelized with a parallel for directive, the default scope of the loop variable is private.

#### **Legal Forms for Parallelizable for Statements**



- OpenMP will only parallelize for loops for which the number of iterations can be determined.
  - From the for statement itself
  - Prior to execution of the loop

# **Examples of Illegal Forms**



The "infinite loop" cannot be parallelized.

```
for (;;) {
    . . .
}
```

The following loop cannot be parallelized.

```
for (i = 0; i < n; i++) {
   if ( . . . ) break;
}</pre>
```

Since the number of iterations can't be determined from the for statement alone. This for loop is also not a structured block, since the break adds another point of exit from the loop.

## **Legal Forms**



Legal forms for parallelizable for statements

### **Caveats**



- The variable index must have integer or pointer type (e.g., it can't be a float).
- The expressions start, end, and incr must have a compatible type. For example, if index is a pointer, then incr must have integer type.
- The expressions start, end, and incr must not change during execution of the loop.
- During execution of the loop, the variable index can only be modified by the "increment expression" in the for statement.

### **Data Dependencies**



```
fibo[0] = fibo[1] = 1;
       for (i = 2; i < n; i++)
          fibo[i] = fibo[i-1] + fibo[i-2];
                                          note 2 threads
        fibo[0] = fibo[1] = 1:
     # pragma omp parallel for num_threads(2)
        for (i = 2; i < n; i++)
           fibo[i] = fibo[i-1] + fibo[i-2];
                                    but sometimes
                                    we get this
1 1 2 3 5 8 13 21 34 55
       this is correct
                          1123580000
```

### **Possible Case**



thread 0: fibo[2] fibo[3] fibo[4] fibo[5] thread 1: fibo[6] fibo[7] fibo[8] fibo[9]

- •Correct: thread 0 finishes its computations before thread 1 starts.
- •Incorrect: thread 0 has not computed fibo[4] and fibo[5], when thread 1 computes fibo[6]. It appears that the system has initialized the entries in fibo to 0, and thread 1 is using the values fibo[4] = 0 and fibo[5] = 0 to compute fibo[6] and so on.

# What happened?



- 1. OpenMP compilers don't check for dependences among iterations in a loop that's being parallelized with a parallel for directive.
- 2. A loop in which the results of one or more iterations depend on other iterations cannot, in general, be correctly parallelized by OpenMP.

# Estimating $\pi$



$$\pi = 4 \left[ 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \cdots \right] = 4 \sum_{k=0}^{\infty} \frac{(-1)^k}{2k+1}$$
**double** factor = 1.0;
**double** sum = 0.0;
**for** (k = 0; k < n; k++) {
 sum += factor/(2\*k+1);
 factor = -factor;
}
pi\_approx = 4.0\*sum;

### **OpenMP solution #1**



```
loop dependency
double factor = 1.0;
double sum = 0.0;
pragma omp parallel for num_threads(thread_count) \
   reduction (+: sum)
for (k = 0; k < n; k++) {
   sum += factor/(2*k+1);
  -factor = -factor;
pi_approx = 4.0*sum;
```

### **OpenMP solution #2**



```
Shared scope.
      double factor = 1.0;
      double sum = 0.0;
#
      pragma omp parallel for num_threads(thread_count) \
         reduction(+:sum)
      for (k = 0; k < n; k++) {
         if (k \% 2 == 0)
            factor = 1.0;
         else
            factor = -1.0:
         sum += factor/(2*k+1);
      pi_approx = 4.0*sum;
```

## **OpenMP solution #3**



```
double sum = 0.0;
pragma omp parallel for num_threads(thread_count) \
    reduction(+:sum) private(factor)

for (k = 0; k < n; k++) {
    if (k % 2 == 0)
        factor = 1.0;
    else
        factor = -1.0;
    sum += factor/(2*k+1);
}</pre>
Insures factor has
private scope.
```

## Summary



- Have a good holiday
- Looking forward to see you guys in person
- Don't forget to bring your laptop for the CUDA lectures.