WODULE - 03.

(NIO-SOL - schema pleniblity, dynamice schemas, replication, simple relationship, replicalation Integn shorizontal scalability, distributed tuples, somi Fring structured data.

Iscues in Mosol

elyba Voice showays

- Standauirestion de not three
- Procusing and difficult in executing complex queves.

Characteristics

1) MOSO! - non relation database storage system that finible data model.

> Key-Value Pair consunder (Eare book)

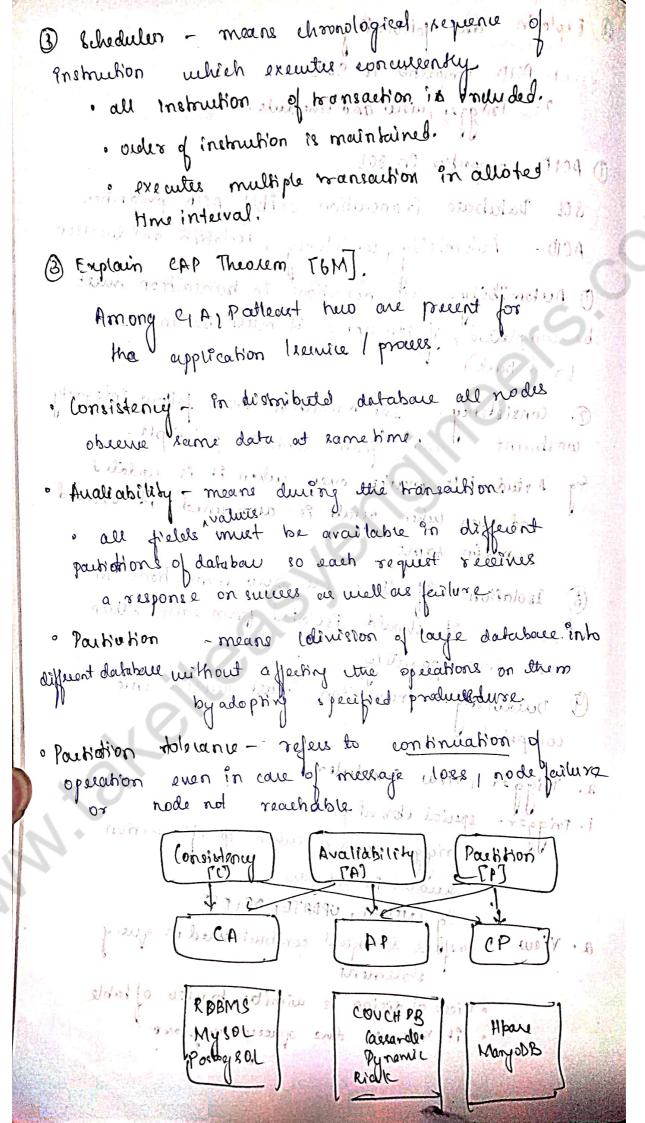
HBOSE HOZT

Rey-Value Pair.

- no fixed rehema such astable or yoins. @ 11080L data repléation le possible control data store is fault torelent.

109 full

@ Emplain the following 1. ACID properties in SOL 2. Tot gger, view and schedule. Amidiation of Newhork by subject . O ACID propadies in sol 301 Database Transcution oxibit ACID properties. ACID - Automicity, consistency, isolation and duration 1) Automikity- all operations in transaction must be completed, if itempted it must be undone (roll beuli) 2 Consistency - The transaction must follow integrity unstraint and follow contraint principle. Eg Astudent 12 meules even when it is updated and even when result is augmovened. The should remain same. (3) Isolation- af there are turo transactions in databare it should isolated from each other and done separately. (9) Durability - transaction must perusist once completed. 2. Trigger, view scheduler. 1. Trigger - special (bro ed produce. docident trigger executes, enhen specific action aciones in détabace! Eg, INSERT, UPDATE, DELETE. a. View -1. reflere to logical construct med in query stalement · View of abelon is winibe to view of table it redecesse dhe queey complexe



@ Emplain BASE properties. [BM]. BASE is flerible model for no sor databave BA - means Basic Availibility g - means eat - means eventual constituency - ensures by distribution of thrads across may data nodes with high degree of replication. O soft - ensures processes y over in the presente inconsistancies but achieving consistency 3) Eventual consistency means consistency requirement in NOSO L database meeting at some point 10 februse. (3) Emploien tey Malue Datastore its advantages and limitation's HMJ · The simplest way to represent a scheme-less data close je trey value pairs. I so sond o mando e Rata reteival is fast in key-vair pair data store · key maps to laye data iting Cey value store allers ule a areceing the values. Chalue ! Judue 2) [roalue H-1]

Value.

" (ategory: student; class: B. Tech i Somestor : VIII; Branch : Engineui of; MOBILE: 9999876532;

200 grang 3 80 2 150 cm

Advantages of Cey Value Store.

1. Data store can store any data in field value.

2. A query just request the values and reliens. the value as a single item.

3. Loy-value store is consistent

4. tey-value datastore may be Westerchy so may be ordered key-value store

5. Rehuned values on queues can be conveiled ento list, table - comme y data fram fields ad columns.

6. Have o) scalability is rewabialty iii) portability

(v) law operational wet

7. The bey can be synthetic or auto gonerated.

dimitations of key value architecture.

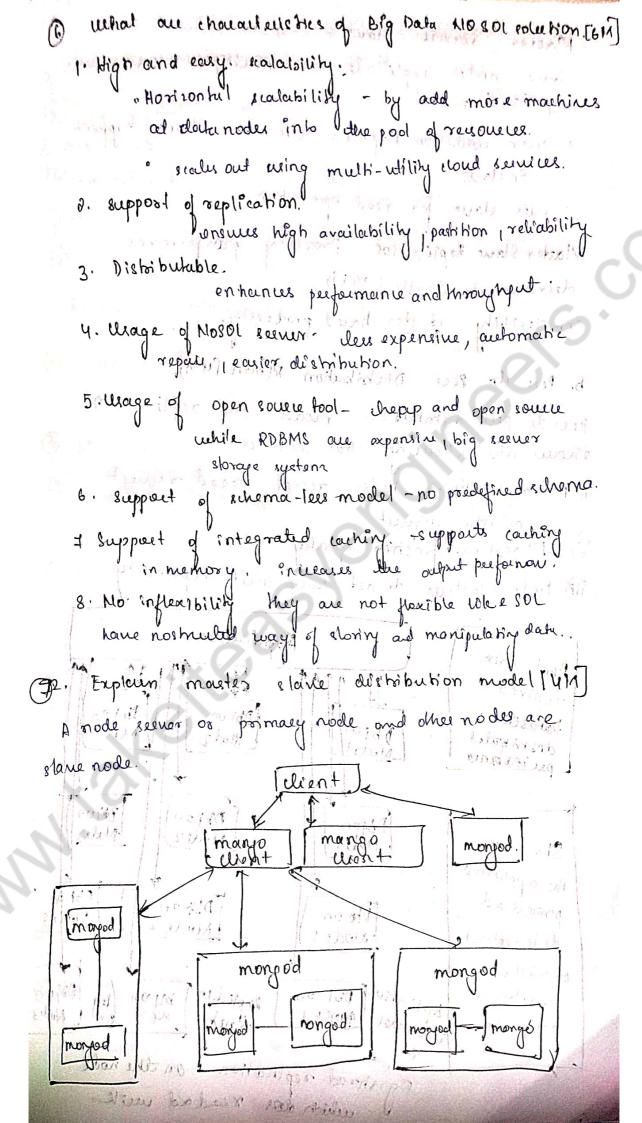
1. No Endexes are maintained on values.

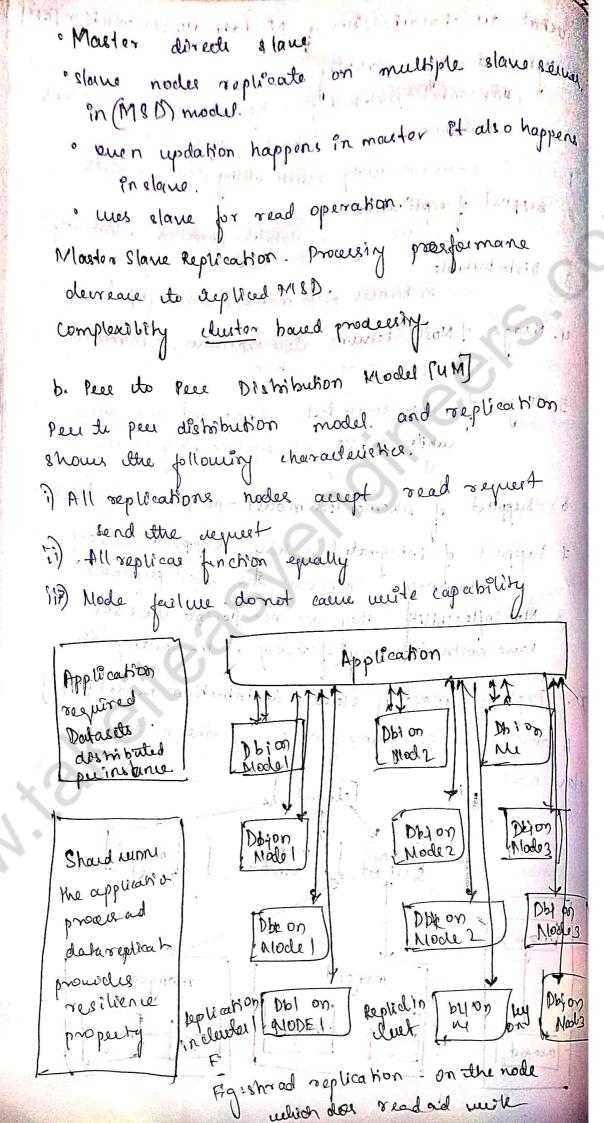
2. 91 does not provide traditional dutabale.

capabilitées such as AICD transaction.

3. Maintaining unique values as key may become difeult when of data incrases.

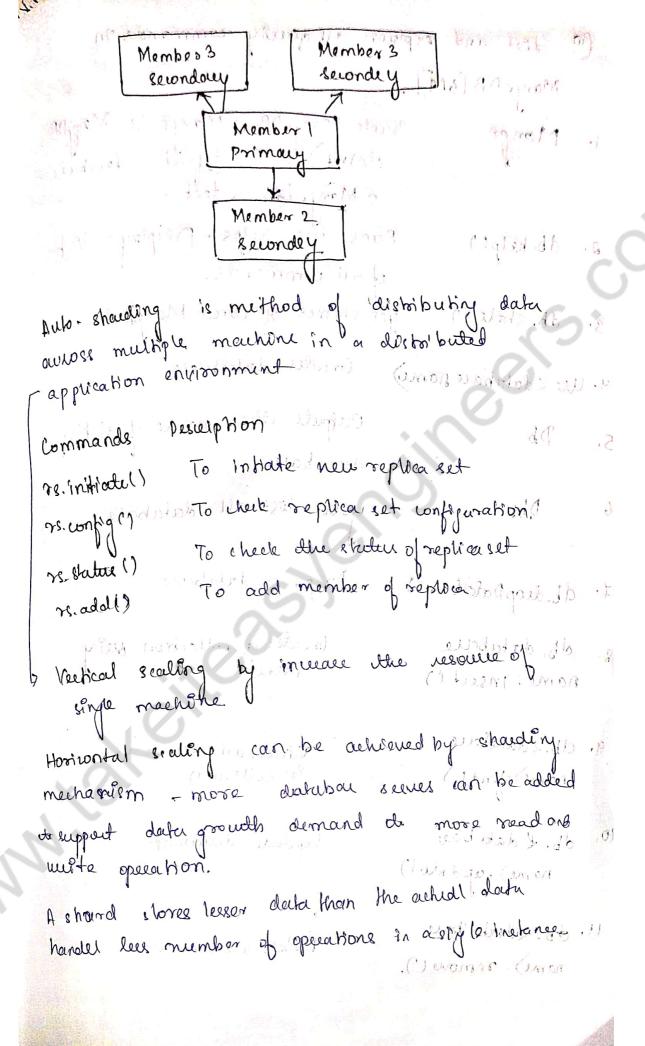
4. Quevies can not be pérférence on moi individuel values





B Ust and explain any 10 features of Mangolis (312). Mango DB is open source DBMS. Monyo program croales and manage destabale. · Margo DB madrayes collection and document close. (has acteels tie) NO SOL ii) distribuled iii) open-source. in) non-relational v) Scalable villonible vili) Indexed. in) multi-moister xi) fault-tole rant. Feature of Margo DB. Distributed DB - makes et highly available and provides nomontal ecalability. wal M art of Levishin moderal fri (2) No Complex Joins. Quelty, indexing and ocal time agrayation allows accessing and analysing the data efficiently (4) Deep query abôlity supports dynamic quevées on documents ining a document bould query lenguge that is nearly as powerful as soll. 5 Mango DB Ps a doument dake vore en which collection holds difficient doughentes Data store is en formand IsoNI document. Osbreng of data is perible: datastore vonsiste of IsoM document, The fields in one document can vouy from whel dataitoute can be changed one How. of downerst in JEDN format 1 Doument model le well defined, Stouherse of document is clear. Document is the unit of sonin date en Manyol B databale.

- O Collection closes a number of Many a DB documents
 Collection uxists with style DB to achieve
 a style purpose. Collection may shore document
 that do not have same fields.
- 1 Indenes on any field
- 1 Atomic operation to the third
- 1 Bast en place of updote
- (3) Conversion (mapping)
- 9 How does reproation achieved in Many BB [6M].
- Replécation ensure ligh availlability in Big Data
- Replicat set la group of mongod (Many PB seerus)
 processes that store same dataset
 - · A replica set usually has minimum three nodes.
 - one of whem breaked primary, receives all muite operation.
 - " all other nodes are secondary
 - a primary node can choose among the secondary node at the time of automatic failouer or meintenance



Lost and explain all operary commands In Maryo DB (6M) . 1 10000 Shute ManyorB. Manyo 13 ManyorB 1. Maryo. client). The default database In Manyo DB 98 test Runs the helps. Displaye 1891 db. help() of all commande. yet shines of about Mary B seener. db. state () Creates databall 4. Use < database name) Outpute the name of earlstory Db 5. to sign databall di Dhe may to yet west all databale in 7. db.deophalabari) Drops a databare. breates a collection being g. db. databolle name. Insect () 9. ab. Ldarehouse Vilceus all documents mounts in collection name) find () Update document 10. db. K dake bowe name), up duto!) will month of the world server broads a 11. db. 2d vota boue Delete a doument name). remove ().

D List and explain feedure and components dan lasandea DB ToM. Feature of laesandes. 1. Manimires the number of unites - units are not very worky 2. Maxinures data dupléation. 3. Does not support join, group by OR clause and eggregation. 4. the clares consisting of ordered keys somi-structured datustorage system 6.9s fact and easily realable with with operation spreed accross the cluster. 6. 98 dostributed DBMR designed for handiby a ligh volume of instructured data a cross multiple cloud servers 7. Ites per do per doctribution in system ougross pri India. componente of caerandles, a magne mon & Mode sur place where data closes for proceeding Data unter Collection of many related nodes Collection of many data center Used for wash reversely each unite Momory resistant data conchere Mem table wilton in committees blata wei An "in membeble temproniy

SST Touble

uchen mem table realis a wetain threehood I date fruet ing. 8ST Table.

Bloom Filter

Fact and memory officient 1
probabilistic data structure de
find whether an olement 10
present in a set 1 Bloom fitel
are accused after every query

(1a) Explain

a. Consistency Command. [4M]

Consistency command shows the weent consistency level.

1. ALC: loigh consistent. A unite must be unition to committing and membable and all replica nodes in actuster.

J. EACH _OUDRAM: A willte muet be wilten in committee and mentable on all s replica nodes in datacenter

3. local -OUDRAM! A with noust be weitten in commit log and membelle on replice node in same duston.

4. ONE: A milte muit be mitten to committer and numbable, af a least one replies node in duster.

5. LOCAL-OME: same as one but least hum or Inner replier nodes, respectivel

remained placement of an area

Memberlae a la la la manda de la compresa la la manda de la manda

6. ANY: A write must be writen to ordearst more l'one mode minit! 1.10cAL-ONE: A milte muit be muiton for al loast one replica, node in local data untor. s serial: Linearable consietency to prevent unconditional update. q. vocal-serial! same as sevial but restricted de local data center