Concepts as she is spoke

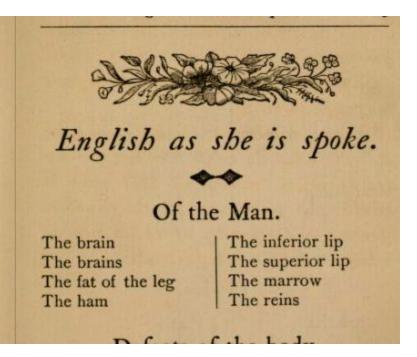
Concepts in the C++2a Working Draft

What's with the title?

4 English as she is spoke.

portuguese; indispensable object whom wish to speak the english and portuguese languages correctly.

We expect then, who the little book (for the care what we wrote him, and for her typographical correction) that may be worth the acceptation of the studious persons, and especially of the Youth, at which we dedicate him particularly.



https://publicdomainreview.org/collections/english-as-she-is-spoke-1884/

What's with the title?

```
Appréndi ô francêz.
                                       I have learned the french language.
 Dôe-me à cabéca.
                                       My head is sick.
 Não tênho têmpo.
                                       I have no time.
 Não pósso demorár-me.
                                       I cannot to stayme.
 Conheço-o de vista.
                                       I know him by sight.
 Tênho dê ôu dêvo sahir.
                                       I have to go out.
Tenho de ou devo receber dinheiro no I have any money to receive at last
  fim độ mêz.
                                         month.
 Dár-vôs-hèi ôu dár-hè-hèi ó sêu ende- I will give you on's adress.
  rêco.
 Refiro-me ao que dizeis ou diz.
                                       I report me at this you tell me.
 Não sôu tôlo.
                                       I am not to silly.
 Não entêndo ôu intendo isso.
                                       I not understand that
```

https://books.google.com/books?id=lUd5AAAAIAAJ&pg=PA75

Outline

- Why requires-clauses? [9–17]

 - What's wrong with good old-fashioned enable_if?
- The *requires-expression* as she is spoke [18–24]
 - Subsumption and normalization [25–56]
- The requires-clause as she is spoke [57–73]
 - Pitfalls for the unwary [74–82]
- Terse syntax(es) [83–88]
- Miscellaneous tidbits [89–97]

Hey look! Slide numbers!

What's the state of Concepts?

There's the "Concepts Technical Specification (Concepts TS)," and then there's the part of it that has been merged into the C++2a working draft.

GCC supports (almost all of) Concepts TS.

```
g++ -std=c++2a -fconcepts -Dconcept="concept bool" test.cc
```

Saar Raz's fork of Clang supports (almost all of) C++2a Concepts.

```
clang++* -std=c++2a test.cc
```

MSVC 2018 can kinda parse Concepts TS syntax but basically treats it as whitespace; it doesn't "support" Concepts in any meaningful sense.

```
cl -std:c++latest -experimental:concepts test.cc
```

https://godbolt.org

What's the state of Concepts?

void f(SomeConcept x); void f(auto x);

Concepts TS concepts far Saar Raz with Concepts Concepts NSVC 2018 with NSVC 2018 with Concepts Υ Υ Ν Ν Ν template<class T> concept Foo = ...; Υ Υ IGN Ν template<class T> concept bool Foo = ...; Ν Υ ICE IGN Ν Ν template<class T> concept bool Foo() { return ...; } Υ Υ Υ IGN Υ template<int I> concept ...; Υ Υ Υ ICE bool b = SomeConcept<int>; Υ Υ Υ Υ IGN template<class T> requires SomeConcept<T> void f(); Ν template<SomeConcept T> void f();

Υ

Υ

Ν

Ν

Ν

What's the state of Concepts?

This is a very silly diagram.		ncepts	is included in the second of t	epts to	in Gaar	Razinental 2018 with a 2018 wi
<pre>template<class t=""> concept Foo =;</class></pre>	N	N S	Y	A Ch	N N	sx, cour
<pre>template<class t=""> concept bool Foo =;</class></pre>	Y	Υ	N	N	IGN	
<pre>template<class t=""> concept bool Foo() { return; }</class></pre>	Υ	ICE	N	N	IGN	
template <int i=""> concept;</int>	Υ	Υ	Υ	Y	IGN	
<pre>bool b = SomeConcept<int>;</int></pre>	Υ	Υ	Υ	Υ	ICE	
<pre>template<class t=""> requires SomeConcept<t> void f();</t></class></pre>	Υ	Υ	Υ	Υ	IGN	
<pre>template<someconcept t=""> void f();</someconcept></pre>	Υ	Υ	Υ	Υ	N	
<pre>void f(SomeConcept x); void f(auto x);</pre>	Υ	Υ	N	N	N	,

What is a concept, anyway?

```
template < class T>
std::string stringify(const T& t) {
    std::ostringstream oss;
    oss << t;
    return std::move(oss).str();
}</pre>
```

What *kinds* of arguments can we pass to this function?

- "Anything with an << operator"
- Just that one requirement? Probably doesn't need Concepts.

What is a concept, anyway?

```
template<class Body>
struct http request {
    int version;
    string method, target;
    map<string, string> fields;
    Body::value type body;
    void read(istream& is) {
        read header(is, *this);
        Body::read(is, this->body);
    void write(ostream& os) const {
        write header(os, *this);
        Body::write(os, this->body);
```

This example comes from
Vinnie Falco's CppCon 2017 session
"Make Classes Great Again"
https://youtu.be/WsUnnYEKPnI

What is a concept, anyway?

```
template<class Body>
struct http request {
    int version;
    string method, target;
    map<string, string> fields;
    Body::value type body;
    void read(istream& is) {
        read header(is, *this);
        Body::read(is, this->body);
    void write(ostream& os) const {
        write header(os, *this);
        Body::write(os, this->body);
```

Many requirements:

- Body::value_type
- Body::read
- Body::write

Let's also say we want to prevent callers from using http_request with a "half-implemented" Body type. Implementing Body::write() but not Body::read(), for example. Let's forbid that.

This example comes from
Vinnie Falco's CppCon 2017 session
"Make Classes Great Again"
https://youtu.be/WsUnnYEKPnI

Defining a concept (old-school)

```
template<class B, class = void> struct is body : false type;
template<class B>
struct is body<B, void t<</pre>
  typename B::value type,
  decltype(
    B::read(declval<istream&>(), declval<typename B::value type&>())
  decltype(
    B::write(declval<ostream&>(), declval<typename B::value type const&>())
>> : true type {};
                                                              This example again comes from
static assert(is body<file body>::value);
                                                            Vinnie Falco's CppCon 2017 session
static assert(is_body<string_body>::value);
                                                               "Make Classes Great Again"
                                                             https://youtu.be/WsUnnYEKPnI
static assert(not is body<int>::value);
```

Defining a concept (C++2a)

```
template<class B>
concept Body = requires(
   istream& is,
    ostream& os,
    typename B::value type& b,
   typename B::value type const& cb)
    typename B::value type; // redundant
   B::read(is, b);
   B::write(is, cb);
};
static assert(Body<file body>);
static assert(Body<string body>);
static assert(not Body<int>);
```

Asserting and SFINAEing (old-school)

```
template<class Body>
struct http request {
    static assert(is body<Body>::value, "Body requirements are not satisfied");
};
template<class Body, class = enable if t<is body<Body>::value>>
void do something() { puts("yes"); }
template<class NonBody, class = enable if t<not is body<NonBody>::value>>
void do something() { puts("no"); }
do something<file body>(); // prints "yes"
                                                    This enable if prevents
do something<int>();  // prints "no"
                                                   ambiguity during the
                                                   overload resolution of
                                                   do something<file body>().
```

Asserting and SFINAEing (C++2a)

```
template<class B>
struct http request {
    static assert(Body<B>, "Body requirements are not satisfied");
};
                                                            This template is
template<class B> requires Body<B>
                                                            "constrained."
void do something() { puts("yes"); }
template<class B>
void do_something() { puts("no"); }
                                                            This template is
                                                            "unconstrained,"
                                                            and therefore a
                                                            "less good" match
do something<file body>(); // prints "yes"
                                                            during overload
do something<int>();  // prints "no"
                                                            resolution.
```

Something else about the old way...

```
template<class Body, class = enable_if_t<is_body<Body>::value>>
void do_something() { puts("yes"); }
```

We're basically abusing template parameters here — we have a template type *parameter* that is not being used to *parameterize!*A caller who knows we're doing this can "hijack" our implementation.

```
do_something<int>();  // correctly does not compile
do_something<int, char**>(); // prints "yes" (yikes!)
```

requires clauses aren't template parameters

```
template<class Body> requires is_body<Body>::value
void do_something() { puts("yes"); }
```

The constraint is properly relegated to a requires clause, which is not a template parameter and thus cannot be hijacked.

```
do_something<int>();  // correctly does not compile
do_something<int, char**>(); // also does not compile
```

Q: How do requires-clauses interact with name-mangling?

A: They don't. (Yet.)

requires doesn't affect mangling*

```
template<class T, class = enable if t<is integral v<T>>>
T foo(T t) { return t + 1; }
                                                          old-school
                        // _Z3fooIi<mark>v</mark>ET_S0_
                                                          enable_if
template int foo<int>(int);
template float foo<float>(float); // does not compile
template<class T> requires is integral v<T>
T bar(T t) { return t + 1; }
                        // _Z3barIiET_S0_
template int bar<int>(int);
template short bar<short>(short);  // Z3barIsET S0
template float bar<float>(float); // does not compile
```

requires doesn't affect mangling*

```
template<class T, enable if t<is integral v<T>, int> = 0>
T foo(T t) { return t + 1; }
template<class T, enable if t<not is integral v<T>, int> = 0>
                                                                    old-school
T foo(T t) \{ return t + 2; \}
                                                                    enable if
                             // _Z3fooIi<mark>Li0</mark>EET S0
template int foo<int>(int);
template short foo<short>(short);  // _Z3fooIsLi0EET_S0_
template float foo<float>(float); // Z3fooIfLi0EET S0
template<class T> requires is integral v<T>
T bar(T t) { return t + 1; }
template<class T> requires not is integral v<T>
                                                                    new-school
                                                                    requires
T bar(T t) \{ return t + 2; \}
template int bar<int>(int);
                                     // Z3barIiET S0
template short bar<short>(short);  // Z3barIsET S0
template float bar<float>(float);
                                     // Z3barIfET S0
```

Duplicate function templates are OK

We need a uniquely mangled name for each *template function*, but not necessarily for each *function template*.

Different story with class templates

A constraint on a class, variable, or alias template is forever. But you can place different constraints on different partial specializations.

```
template<class T>
class Widget { ... };

template<class U> requires is_pointer_v<U>
class Widget { ... };  // Error: redefinition of class template "Widget"

template<class V> requires is_pointer_v<V>
class Widget<V> { ... };  // OK: constrained partial specialization
```

GCC doesn't fully support constrained alias templates.

Remember this old-school trap?

Recall a trap I discussed in "Template Normal Programming, Part 1."

At this point, if we call baz(&i), overload resolution will consider two candidates:

- the template function baz<int*>(int*) that could be instantiated from template "A"
- the template function baz<int>(int*) that could be instantiated from template "B"

"B" is a better match than "A".

So the compiler uses template "B" to instantiate baz<int>(int*) a.k.a. _Z3bazIiEPT_S1_ ... and our explicitly specialized baz<int*>(int*) a.k.a. _Z3bazIPiET_S1_ is never called!

With requires, the stakes are higher*

At this point, if we call baz(&i), overload resolution will consider two candidates:

- the template function baz<int*>(int*) that could be instantiated from template "A"
- the template function baz<int*>(int*) that could be instantiated from template "B"

"B" is a better match than "A", because "B" is constrained and "A" is not.

So the compiler uses template "B" to instantiate baz<int*>(int*) a.k.a. _Z3bazIPiET_S1_ ... which has exactly the same mangled name as our explicitly specialized baz<int*>(int*) a.k.a. Z3bazIPiET_S1 ! This is a violation of the One Definition Rule, and causes undefined behavior.

Two distinct "better match" criteria

When we call foo(&i),

- "C" is a better match than "A" because its function parameter list is more highly specialized.
- "C" is a better match than "B" because its function parameter list is more highly specialized.
- "B" is a better match than "A" because it is more constrained.

Two distinct "better match" criteria

The old-school rule has highest priority! Look first at the parameter list; *then* (and only then) consider the requires clause.

When we call foo(&i),

- "C" is a better match than "A" because "unction parameter list is more highly specialized.
- "C" is a better match than "B" because its function parameter list is more highly specialized.
- "B" is a better match than "A" because it is more constrained.

```
template < class T>
T foo(T t) { return t + 1; }  // "A"

template < class T> requires is_integral_v < T>
T foo(T t) { return t + 2; }  // "B"
```

• "B" is a better match than "A" because it is more constrained.

• "B" is a better match than "A" because it is more constrained.

```
template < class T > requires (sizeof(T) == 4)
T foo(T t) { return t + 3; }  // "C"
```

Intuitively, neither "B" nor "C" is "more constrained" than the other.

• "B" is a better match than "A" because it is more constrained.

```
template < class T > requires (sizeof(T) == 4)
T foo(T t) { return t + 3; }  // "C"
```

• Intuitively, neither "B" nor "C" is "more constrained" than the other.

```
template < class T> requires (sizeof(T) != 0)
T foo(T t) { return t + 4; } // "D"
```



• "B" is a better match than "A" because it is more constrained.

```
template < class T > requires (sizeof(T) == 4)
T foo(T t) { return t + 3; }  // "C"
```

• Intuitively, neither "B" nor "C" is "more constrained" than the other.

```
template < class T > requires (sizeof(T) != 0)
T foo(T t) { return t + 4; }  // "D"
```

 Logically "D" is less constrained than "C", but we don't seriously expect the compiler to realize that. So intuitively it makes sense that all of "B", "C", and "D" are equally much constrained.

So do we only have two levels — "constrained" and "unconstrained"? Let's try some things we think the compiler might be able to figure out.

What do we expect to happen here? Will the compiler pick "G", or say there's an ambiguity?

So do we only have two levels — "constrained" and "unconstrained"? Let's try some things we think the compiler might be able to figure out.

Trick question! The Concepts TS (GCC) is actually happy with this. But the C++2a version of Concepts (Clang) says "no":

```
template < class T >
concept Integral = is_integral_v < T >;

template < class T > requires Integral < T >
T foo(T t) { return t + 10; }

// "J"

template < class T > requires Integral < T > && is_signed_v < T >
T foo(T t) { return t + 11; }

int x = foo(0);
```

What do we expect to happen here? Will the compiler pick "K", or say there's an ambiguity?

```
template < class T >
concept Integral = is_integral_v < T >;

template < class T > requires Integral < T >
T foo(T t) { return t + 10; }

// "J"

template < class T > requires Integral < T > && is_signed_v < T >
T foo(T t) { return t + 11; }

int x = foo(0);
```

The compiler picks "K"!

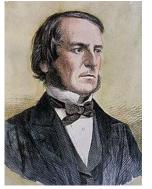
Subsumption

Constraints are logical formulas!



requires-clause:
requires constraint-logical-or-expression

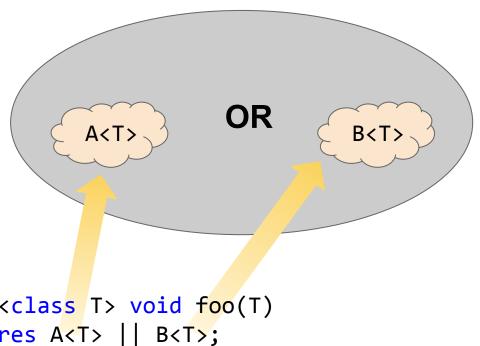
constraint-logical-or-expression :
 constraint-logical-and-expression
 constraint-logical-or-expression | | constraint-logical-and-expression



constraint-logical-and-expression :
 primary-expression
 constraint-logical-and-expression && primary-expression

Both Clang and GCC accept inclusive-or-expression where the wording currently says "primary-expression." I expect the wording to catch up soon.

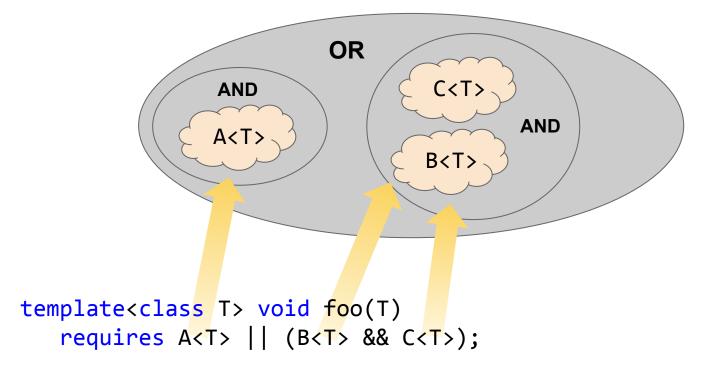
Normalization and normal form



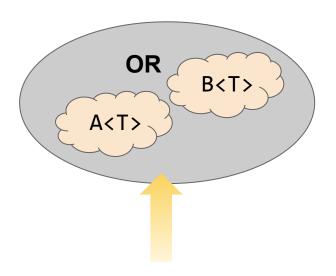
This cloudy Platonic diagram is called the constraint's "normal form." It is a mathematical construct, not a C++ source code construct.

template<class T> void foo(T) requires A<T> || B<T>;

Normalization and normal form



Logical operations are commutative



```
template<class T> void foo(T)
    requires A<T> || B<T>;
template<class T> void foo(T)
    requires B<T> || A<T>;
```

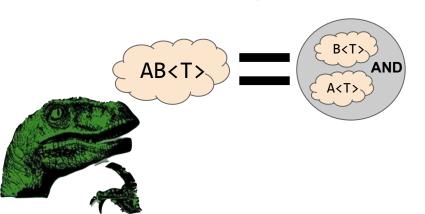
These two requires-clauses have exactly the same *normal* form.

Non-atomic concepts

```
template<class T> concept AB =
    A<T> && B<T>;
```

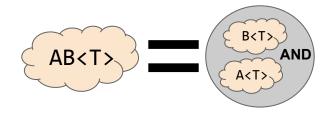
Non-atomic concepts

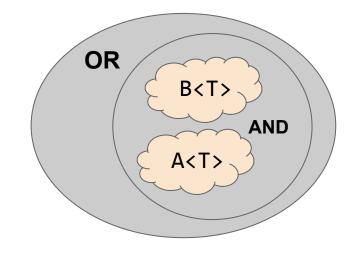
```
template<class T> concept AB =
    A<T> && B<T>;
```



Non-atomic concepts

```
template<class T> concept AB =
    A<T> && B<T>;
```

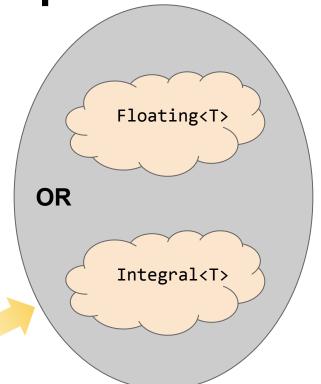




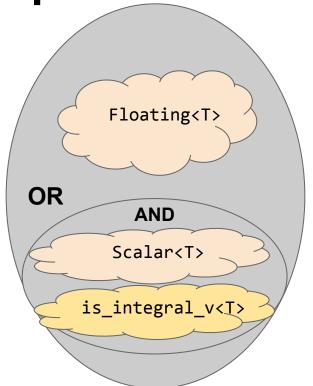
template<class T> void foo(T)
 requires AB<T>;
template<class T> void foo(T)
 requires B<T> && A<T>;

These two requires-clauses have exactly the same *normal* form.

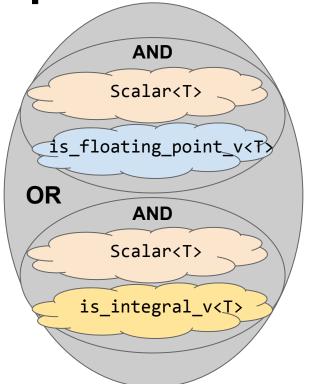
```
template<class T> concept Scalar =
   is scalar v<T>;
template<class T> concept Integral =
   Scalar<T> && is integral v<T>;
template<class T> concept Floating =
   Scalar<T> && is_floating point v<T>;
template<class T> void foo(T)
   requires Integral<T> || Floating<T>;
```



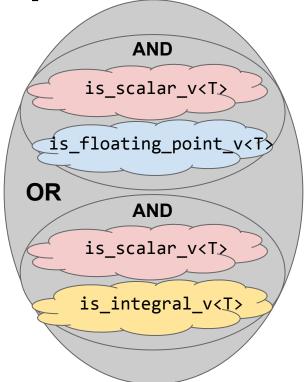
```
template<class T> concept Scalar =
   is scalar v<T>;
template<class T> concept Integral =
   Scalar<T> && is integral v<T>;
template<class T> concept Floating =
   Scalar<T> && is floating point v<T>;
template<class T> void foo(T)
   requires Integral<T> || Floating<T>;
```



```
template<class T> concept Scalar =
   is scalar v<T>;
template<class T> concept Integral =
   Scalar<T> && is integral v<T>;
template<class T> concept Floating =
   Scalar<T> && is floating point v<T>;
template<class T> void foo(T)
   requires Integral<T> || Floating<T>;
```

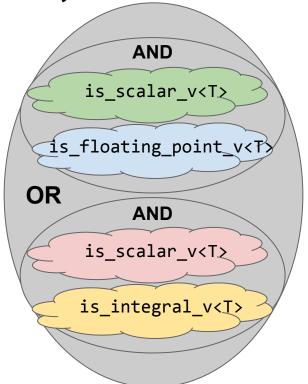


```
template<class T> concept Scalar =
   is scalar v<T>;
template<class T> concept Integral =
   Scalar<T> && is integral v<T>;
template<class T> concept Floating =
   Scalar<T> && is floating point v<T>;
template<class T> void foo(T)
   requires Integral<T> || Floating<T>;
```

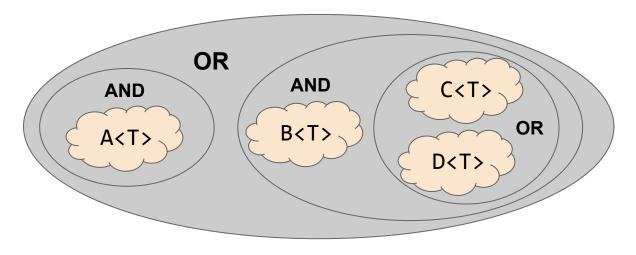


...indexed by source position, not text*

```
template<class T> concept Integral =
   is scalar v<T> &&
                 is integral v<T>;
template<class T> concept Floating =
   is scalar v<T> &&
                 is floating point v<T>;
template<class T> void foo(T)
   requires Integral<T> || Floating<T>;
```

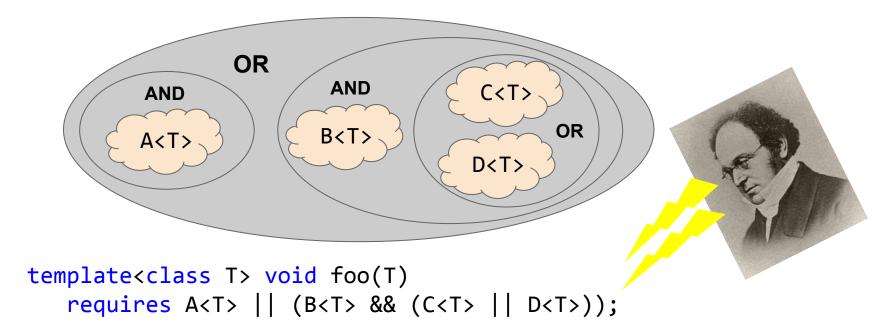


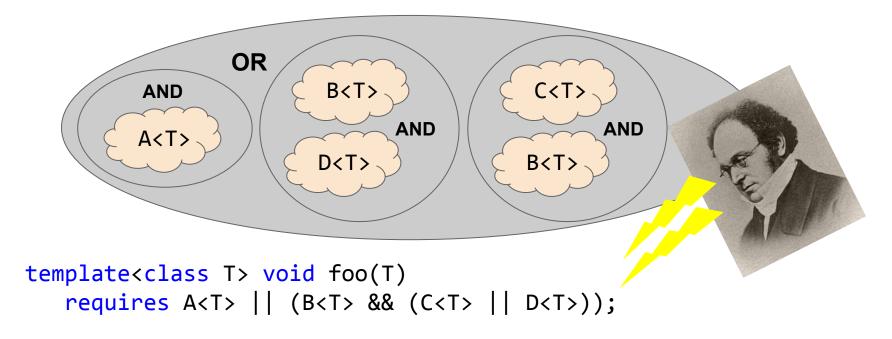
This is slide 31's trick question. The Concepts TS actually did try to deduplicate expressions based on textual spelling, not source position.



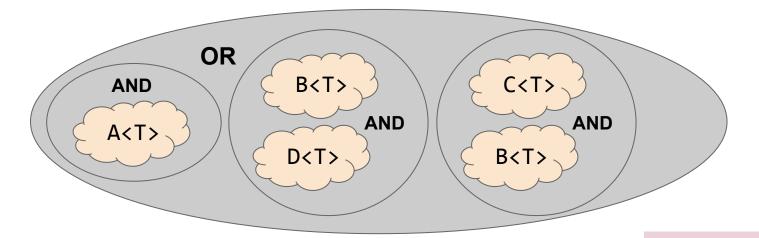
```
template<class T> void foo(T)
  requires A<T> || (B<T> && (C<T> || D<T>));
```

??





```
template<class T> void foo(T)
  requires A<T> || (B<T> && C<T>) || (B<T> && D<T>);
```

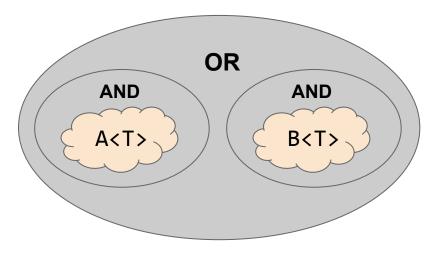


```
template<class T> void foo(T)
  requires A<T> || (B<T> && (C<T> || D<T>));
```

These two requires-clauses have exactly the same *normal* form.

```
template<class T> void foo(T)
  requires A<T> || (B<T> && C<T>) || (B<T> && D<T>);
```

What about negation?

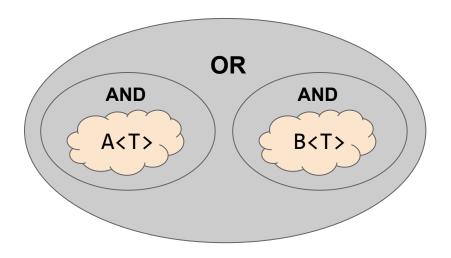


```
template < class T> void foo(T)
    requires not (not A<T> || not B<T>);

template < class T> void foo(T)
    requires (A<T> && B<T>);
```

??

What about negation?

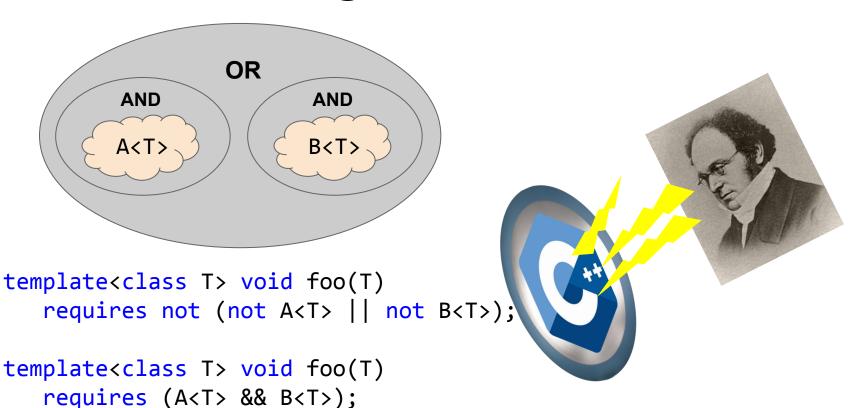


```
template<class T> void foo(T)
   requires not (not A<T> || not B<T>);
```

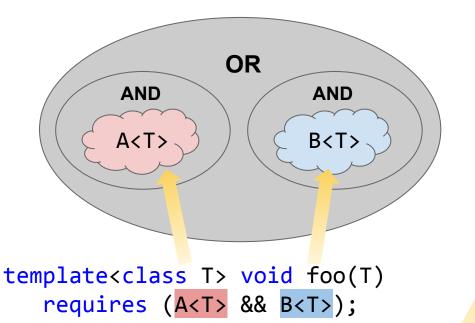
```
template<class T> void foo(T)
  requires (A<T> && B<T>);
```



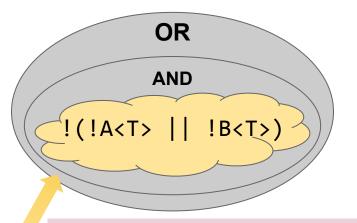
What about negation? NO!



What about negation? NO!



```
template<class T> void foo(T)
   requires not (not A<T> | not B<T>);
```



The grammar has special cases for top-level "| | " and "&&", but not for "!". These two requires-clauses are functionally equivalent, but neither one **subsumes** the other.

What about negation? NO!



The situation is similar to "covariance and contravariance" with virtual methods. Just as in that case, C++2a supports covariance but not contravariance.

```
template<class T> void foo(T) requires Scalar<T>;
   // ...is subsumed by the more constrained template...
template<class T> void foo(T) requires Integral<T>;

template<class T> void foo(T) requires !Integral<T>;
   // ...has no subsumption relationship at all with...
template<class T> void foo(T) requires !Scalar<T>;
```

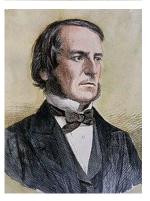
Defining a Concept

Defining a concept

```
template <template-parameter-list>
concept concept-name = constraint-expression;
```

Kinds of constraint-expressions





- Disjunction ab
- Conjunction a && b
- Primary expression of type bool
 - Type trait is_integral_v<T>
 - A requires-expression

Kinds of constraint-expressions





- Disjunction ab
- Conjunction a && b
- Primary expression of type bool
 - Type trait is_integral_v<T>
 - A requires-expression

Noexcept-expression vs. clause

Noexcept-clause is part of a declaration: "This is noexcept when..."

```
template<class T>
void xyzzy() noexcept ( boolean-expression );
```

Noexcept-expression is a constant expression of type bool: "Is this expression noexcept?"

```
constexpr bool plugh =
   noexcept ( unevaluated-expression );
```

Requires-expression vs. clause

Requires-clause is part of a declaration: "This participates in overload resolution when..."

```
template<class T>
void xyzzy() requires boolean-expression ;
```

Requires-expression is a constant expression of type bool: "Is this (set of) requirement(s) satisfied?"

```
constexpr bool plugh =
   requires ( parameter-list ) { requirement-seq };
```

Noexcept-expression vs. clause

Noexcept-clause sometimes contains a noexcept-expression:

```
template<class T>
void xyzzy() noexcept(noexcept(unevaluated-expression ));
```

Or we separate the two:

```
template < class T >
bool is_nothrow_fooable =
    noexcept ( unevaluated-expression );

template < class T >
void xyzzy() noexcept( is_nothrow_fooable_v<T> );
```

Requires-expression vs. clause

Requires-clause sometimes contains a requires-expression:

```
template<class T>
void xyzzy()
    requires requires ( parameter-list ) { requirement-seq };
```

Or we separate the two:

```
template<class T>
concept Fooable =
    requires ( parameter-list ) { requirement-seq };

template<class T>
void xyzzy() requires Fooable<T>;
```

Requires-expression vs. clause

Requires-clause sometimes contains a requires-expression:

```
template<class T>
void xyzzy()
    requires requires ( parameter-list ) { requirement-seq };
```

Or we separate the two:

```
template < class T >
concept Fooable =
    requires ( parameter-list ) { requirement-seq -};

template < class T >
void xyzzy() requires Fooable < T >;
```

```
constexpr bool Fooable = requires ( parameter-list ) {
    typename type-expression;
    { value-expression };
    { value-expression } noexcept ;
    { value-expression } -> type-or-concept ;
    requires constraint-expression;
```

```
template<class T>
concept Fooable = requires ( parameter-list ) {
    typename T::value type;
    { value-expression };
    { value-expression } noexcept ;
    { value-expression } -> type-or-concept ;
    requires constraint-expression;
```

```
template<class T>
concept Fooable = requires (const T ct, int i) {
    typename T::value type;
    { ct + i };
    { ct += i } noexcept;
    \{ ct += i \} -> T\&;
   { ct - ct } -> Integral;
```

```
template<class T>
concept Fooable = requires (const T ct, int i) {
    typename T::value type;
                                      Exactly equivalent to
    { ct + i };
                                        requires noexcept(ct += i);
    { ct += i } noexcept;
    \{ ct += i \} -> T\&;
    { ct - ct } -> Integral;
```

```
template<class T>
concept Fooable = requires (const T ct, int i) {
    typename T::value type;
                                       Exactly equivalent to
    { ct + i };
                                         requires is convertible v<
                                          decltype(ct += i), T&
    { ct += i } noexcept;
                                         >;
    \{ ct += i \} -> T\&;
    { ct - ct } -> Integral;
```

```
template<class T>
concept Fooable = requires (const T ct, int i) {
    typename T::value type;
    { ct + i };
    { ct += i } noexcept;
                                      Exactly equivalent to*
    \{ ct += i \} -> T\&;
                                       requires Integral<decltype(ct - ct)>;
    { ct - ct } -> Integral;
```

Sometimes, requirements can be chained together on a single line of code.

GCC (not C++2a) supports auto

```
template<class T>
concept Fooable = requires (const T ct) {
    { *ct } -> auto&&;
                                          Equivalent to
    { +ct } -> auto*;
                                           requires !is void v<decltype(*ct)>;
                                          Equivalent to
                                           requires is_pointer_v<</pre>
                                            remove reference t<decltype(+ct)>>;
```

Fun traps for the unwary

```
template<class T>
concept Negatable = requires (T t) {
    -t -> T;
};
static_assert(Negatable<char>); // FAILS
```



??

Fun traps for the unwary

```
template<class T>
concept Negatable = requires (T t) {
    -t -> T;
};
struct S {
    int T;
static_assert(Negatable<S*>); // OK
```



??

Fun traps for the unwary

```
template<class T>
concept Negatable = requires (T t) {
    -t -> T;
};
                                        Forgot the curly
                                           braces!
struct S {
    int T;
static_assert(Negatable<S*>); // OK
```





```
template < class T >
concept IntSized = requires {
    sizeof(T) == 4;
};
static_assert(IntSized < char > ); // OK
```





```
template < class T >
concept IntSized = requires {
    sizeof(T) == 4;
};
static_assert(IntSized < char > ); // OK
```

Forgot the "requires" keyword!



```
template<class T>
concept IntSized = requires {
    requires sizeof(T) == 4;
};
// Or even better...
template<class T>
concept IntSized = (sizeof(T) == 4);
```

```
template < class T >
concept NothrowIncrementable = requires (T t) {
    noexcept ( ++t );
};
```



Forgot the "requires" keyword!



```
template<class T>
concept NothrowIncrementable = requires (T t) {
    requires noexcept( ++t );
    // Or even better...
    { ++t } noexcept;
```

Terse syntax(es)

C++2a supports a terser syntax, but I currently recommend not to use it.

```
template<Integral T>
void foo(T t) {
    // ...
} "Constrained parameter"
```

means the same thing as

```
template<class T> requires Integral<T>
void foo(T t) {
    // ...
}
```

Terse syntax(es)

C++2a supports a terser syntax, but I currently recommend not to use it.

```
template < class T > concept Reference = is_reference_v < T >;
template < class T > concept Function = is_function_v < T >;

template < Reference R > void foo(R r) { ... }

template < Function F > void foo(F f) { ... }
```



Perfect forwarding?

```
template<class T> concept Integral = is_integral_v<T>;
template<class T> requires Integral<T>
void oops(T&& value) { ... }
int ie; oops(ie); // does not compile
```

Perfect forwarding?

```
template < class T > concept Range = requires(T&& t) {
    ranges::begin(std::forward<T>(t)) -> Iterator;
    ranges::end(std::forward<T>(t)) -> Iterator;
};

template < class T > requires Range < T >
void okay(T&& range) { ... }
Oops! I fell into the exact trap identified on slide 74! I did not
```

mean to do that.

Terse syntax(es)

Concepts TS supports an *even terser* syntax, commonly known as "the terse syntax." This syntax was **not** adopted into C++2a.

```
void foo(Integral t) { // Look ma, no "template"!
    // ...
}
```

would mean the same thing as

```
template<class T> requires Integral<T>
void foo(T t) {
    // ...
}
```

Terse syntax(es)

The Concepts TS also extends constrained return types and variable type deduction.

```
template<class T>
Integral foo(T i) { // if the deduced type is not Integral, static-assert
    Integral j = i; // if the deduced type is not Integral, static-assert
    return j;
Exactly analogous to C++11's constrained return types:
template<class T>
auto* foo(T i) {  // if the deduced type is not a pointer, static-assert
    auto* j = i;  // if the deduced type is not a pointer, static-assert
    return j;
```

Non-type concepts

C++2a rejects concepts that are *not* templates:

```
concept True = true; // error
But C++2a permits concepts that are weird templates:
template<template<class...> class Ctr>
concept BigContainer = sizeof(Ctr<int>) >= 24; // OK!
static assert(BigContainer<vector>);
template<int I>
concept EvenValue = (I % 2 == 0); // OK!
static assert(EvenValue<42>);
```

Metaprogramming with concepts

Concepts can be nested inside namespaces.
Concepts *cannot* be nested inside classes or structs.
Concepts *cannot* be passed as template parameters.

```
namespace N { template<class T> concept NestedC = ...; } // OK
struct S { template<class T> concept MemberC = ...; } // ERROR
template <template<class T> concept ParamC> // ERROR
struct ST { ... };
```

C++2a permits concepts with multiple template parameters.

```
template<class A, class B, class C>
concept BiggerThan =
    sizeof(A) >= sizeof(B) && sizeof(B) >= sizeof(C);
static_assert(BiggerThan<int[3], int[2], int[1]>);
```

C++2a permits concepts with multiple template parameters.

```
template<class A, class B, class C>
concept BiggerThan =
    sizeof(A) >= sizeof(B) && sizeof(B) >= sizeof(C);
```

Combined with C++2a's constrained-parameter syntax, this gets weird.

```
template<BiggerThan<int[2], int[1]> D>
void foo();
```

```
Equivalent to

template<class D> requires BiggerThan<D, int[2], int[1]>
void foo();
```

Even though I say not to use constrained-parameter syntax, its existence suggests some style guidelines for naming of concepts:

- Your first template parameter is magic. Make sure your name reflects that.
 - ConvertibleTo<Src, Dst>
 - GeneratorOf<F, ReturnType>
- But don't strain yourself. Maybe this concept just doesn't need to play well with constrained-parameter syntax.
 - requires Mergeable<InIter1, InIter2, OutIter>
 - Does it even need to be a concept at all?
 - requires is_mergeable<I1, I2, 0>

Concepts enable subsumption.
We'll never use Mergeable in a context requiring subsumption. So it needn't be a concept.

Even though I say not to use constrained-parameter syntax, its existence suggests some style guidelines for naming of concepts:

- Constrained-parameter syntax eliminates the visual distinction between type and non-type parameters. Make sure your concept's name is unambiguous.
- I strongly recommend the suffix -Value for non-type concepts!

```
o template<class T> concept Arithmetic = ...;
```

o template<auto V> concept ArithmeticValue = ...;

```
template<Arithmetic A> void foo();
template<ArithmeticValue B> void bar();
... foo<int>(); ... bar<42>(); ...
```

We already saw this convention with concept EvenValue.

Constrained non-template functions?

On the surface, constrained non-template functions seem to provide an alternative to #ifdef / if constexpr.

But neither Clang nor GCC supports this idiom. Don't try it.

```
unsigned byteswap(unsigned x) requires (sizeof(int) == 2) {
   return (x << 8) | (x >> 8);
unsigned byteswap(unsigned x)
                requires (sizeof(int) == 4) {
    return (x << 24) | (x >> 8 << 24 >> 8)
         | (x >> 16 << 24 >> 16) | (x >> 24);
```

Because name-mangling. requires means "does not participate in overload resolution"; it doesn't mean "does not get codegenned."

More information on Concepts

- "Constraints and concepts (since C++20)"
 https://en.cppreference.com/w/cpp/language/constraints
- Cpplang Slack channel #concepts https://cpplang-inviter.cppalliance.org
- For examples of usage in practice:
 P0896 "The One Ranges Proposal" warning! 233-page PDF alert!
 http://wg21.link/p0896
- Herb Sutter's P0745 "Concepts in-place syntax" http://wg21.link/p0745
- Corentin Jabot's "The tightly-constrained design space of convenient syntaxes for generic programming" https://cor3ntin.github.io/posts/concepts_syntax/



Sidebar: Linkers hate return-type SFINAE

```
template<class T>
enable if t<is integral v<T>, T> foo(T t) { return t + 1; }
                                                                        enable if on
template<class T>
                                                                        the function's
enable_if_t<not is_integral_v<T>, T> foo(T t) { return t + 2; }
                                                                        return type
template int foo<int>(int);
                              // Z3fooIiENSt9enable ifIX13is_integral_vIT_EES1_E4typeES1_
template short foo<short>(short); // _Z3fooIsENSt9enable_ifIX13is_integral_vIT_EES1_E4typeES1_
template float foo<float>(float);
                                   // Z3fooIfENSt9enable_ifIXnt13is_integral_vIT_EES1_E4typeES1_
template<class T> requires is integral v<T>
T bar(T t) { return t + 1; }
template<class T> requires not is_integral_v<T>
                                                                       requires
T bar(T t) { return t + 2; }
template int bar<int>(int);
                                       // Z3barIiET S0
template short bar<short>(short);
                                       // Z3barIsET S0
template float bar<float>(float);
                                       // Z3barIfET_S0_
```