

# **PAPER 4 — MATHEMATICS & CRYPTOGRAPHY**

## **The Formal Cryptographic Model of NOVAK: Deterministic Execution, HVET, RGAC, and Provable Action Boundaries**

### **Abstract**

This paper presents the mathematical foundations of the NOVAK Protocol, focusing on deterministic proof-before-action execution, hash verification primitives, and the Recursive Global Audit Chain. We formally define the system's correctness properties, failure boundaries, and cryptographic assumptions.

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# 1. Introduction

The NOVAK Protocol enforces deterministic computational integrity through:

- **Execution Identity Receipts (EIR)**
- **Hash-Verified Execution Traces (HVET)**
- **Recursive Global Audit Chains (RGAC)**
- **Deterministic Safety Gates**

This aligns NOVAK with research traditions in:

- formal verification
- cryptographic commitments
- model checking
- attested execution
- verifiable computing

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## 2. Mathematical Preliminaries

Let:

- $R$  = rule function
- $D$  = input data
- $O = R(D)$  = correct output
- $HR = H(R)$
- $HD = H(D)$
- $HO = H(O)$

NOVAK constructs:

**$HVET = H( HR \parallel HD \parallel HO \parallel \text{timestamp} )$**

Where  $H$  is a cryptographic hash satisfying:

- preimage resistance
- second-preimage resistance
- collision resistance

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### 3. Deterministic Safety Gate

The Safety Gate enforces:

If  $\text{HNET\_proposed} \neq \text{HNET\_expected} \rightarrow \text{Reject Execution}$

No probabilistic interpretations.

No heuristics.

No learning-based inference.

It is a **hard, algebraic execution rule**.

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## 4. Recursive Global Audit Chain (RGAC)

Each action appends:

$$\text{RGAC}(n+1) = H( \text{RGAC}(n) \parallel \text{HVET}(n) )$$

This produces:

- immutability
- global consistency
- independent verifiability
- temporal ordering guarantees

Aligned with blockchain **only in audit**,  
NOT in execution, consensus, or decentralization.

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## 5. Correctness Theorems

### Theorem 1: Undetectable Tampering is Impossible

If  $H$  is secure and HVET includes  $HR$ ,  $HD$ ,  $HO$ , then:

Any attempt to modify:

- the rule
- the input
- the output
- intermediate values

produces a mismatch.

### Theorem 2: Unauthorized Execution is Undecidable Under NOVAK

Given deterministic  $R$  and validated HVET:

$R'(D')$  cannot execute if  $H(R')$ ,  $H(D')$ ,  $H(R'(D'))$  differ

### Theorem 3: Forward Integrity

No function of future state can retroactively alter past RGAC entries without breaking hash hardness assumptions.

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## 6. Computational Hardness Model

NOVAK relies solely on:

- SHA-256 class hash functions
- Concatenation commitments
- No need for asymmetric cryptography
- No need for zero-knowledge proofs
- No need for consensus protocols

This increases:

- speed
- scalability
- implementability
- verifiability

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## 7. Conclusion

NOVAK is a deterministic integrity enforcement system whose mathematical properties ensure unforgeable execution. It represents a new category of cryptographically enforced computation.