

(From Russell & Norvig)

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• Optimal decisions

• Imperfect, real-time decisions

• α-β pruning

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Games- Outline

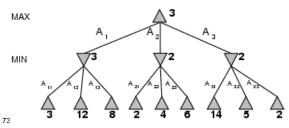
Games vs. search problems

- "Unpredictable" opponent → specifying a move for every possible opponent reply
- Time limits → unlikely to find goal, must approximate

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Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest minimax value = best achievable payoff against best play
- E.g., 2-ply game:

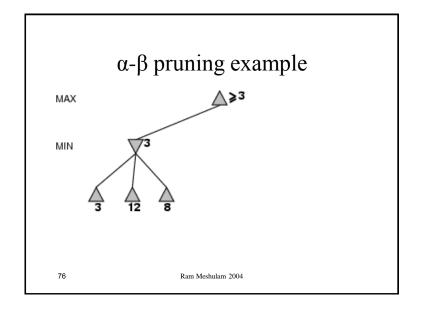


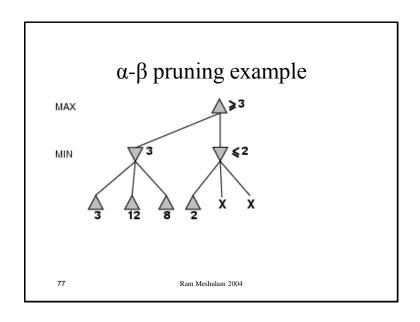
Properties of minimax

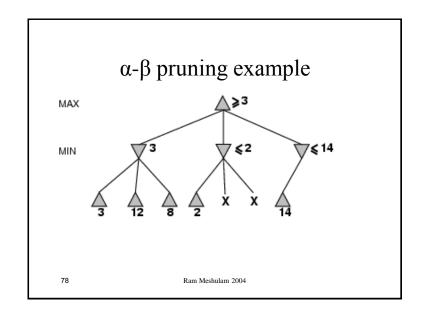
- Complete? (=will not run forever) Yes (if tree is finite)
- Optimal? (=will find the optimal response) Yes (against an optimal opponent)
- $\underline{\text{Time complexity?}} O(b^m)$
- Space complexity? O(bm) (depth-first exploration)
- For chess, b ≈ 35, m ≈100 for "reasonable" games → exact solution completely infeasible

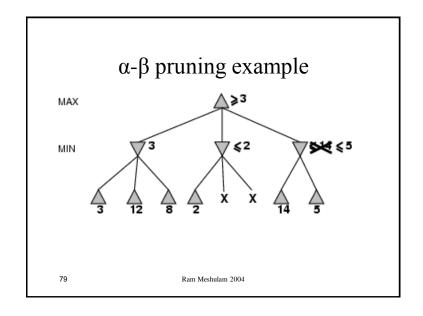
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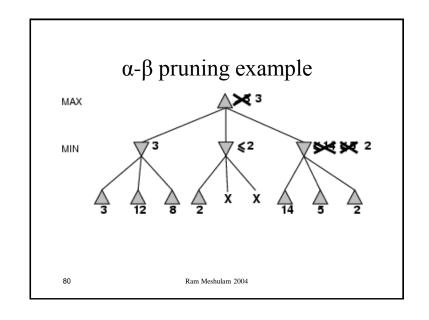
$\begin{array}{c} \textbf{Minimax algorithm} \\ \\ \textbf{function } \textbf{Minimax-Decision}(\textit{state}) \textbf{ returns } \textit{an action} \\ v \leftarrow \textbf{Max-Value}(\textit{state}) \\ \textbf{return the } \textit{action in Successors}(\textit{state}) \textbf{ with value } v \\ \hline \\ \textbf{function } \textbf{Max-Value}(\textit{state}) \textbf{ returns } \textit{a utility value} \\ \textbf{if Terminal-Test}(\textit{state}) \textbf{ then return Utility}(\textit{state}) \\ v \leftarrow -\infty \\ \textbf{for } \textit{a, s} \textbf{ in Successors}(\textit{state}) \textbf{ do} \\ v \leftarrow \textbf{Max}(v, \textbf{Min-Value}(\textit{s})) \\ \textbf{return } v \\ \hline \\ \textbf{function } \textbf{Min-Value}(\textit{state}) \textbf{ returns } \textit{a utility value} \\ \textbf{if Terminal-Test}(\textit{state}) \textbf{ then return Utility}(\textit{state}) \\ v \leftarrow \infty \\ \textbf{for } \textit{a, s} \textbf{ in Successors}(\textit{state}) \textbf{ do} \\ v \leftarrow \textbf{Min}(v, \textbf{Max-Value}(\textit{s})) \\ \textbf{return } v \\ \hline \end{array}$











Properties of α - β

- Pruning does not affect final result
- Good move ordering improves effectiveness of pruning
- With "perfect ordering" on binary tree, time complexity = O(b^{m/2})
 - → doubles depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of metareasoning)

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The α - β algorithm

```
function Alpha-Beta-Search(state) returns an action inputs: state, current state in game v \leftarrow \text{Max-Value}(state, -\infty, +\infty) return the action in Successors(state) with value v function Max-Value(state, \alpha, \beta) returns a utility value inputs: state, current state in game \alpha, \text{ the value of the best alternative for } \text{Max along the path to } state if Terminal-Test(state) then return Utility(state) v \leftarrow -\infty for a, s in Successors(state) do v \leftarrow \text{Max}(v, \text{Min-Value}(s, \alpha, \beta)) if v \geq \beta then return v \alpha \leftarrow \text{Max}(\alpha, v) return v
```

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Why is it called α - β ?

• α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for *max*

MAX

MIN

• If *v* is worse than α, *max* will avoid it

→ prune that branch

.. MAX

• Define β similarly for *min*

MIN

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V_v

The α - β algorithm

```
function Min-Value(state, \alpha, \beta) returns a utility value inputs: state, current state in game \alpha, the value of the best alternative for Max along the path to state \beta, the value of the best alternative for Min along the path to state if Terminal-Test(state) then return Utility(state) v \leftarrow +\infty for a, s in Successors(state) do v \leftarrow \min(v, \max\text{Value}(s, a, \beta)) if v \leq \alpha then return v state state for state state
```

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Resource limits

Suppose we have 100 secs, explore 10^4 nodes/sec $\rightarrow 10^6$ nodes per move

Standard approach:

- cutoff test:
 - e.g., depth limit (perhaps add quiescence search)
- evaluation function
 - = estimated desirability of position

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Cutting off search

MinimaxCutoff is identical to MinimaxValue except

- 1. Terminal? is replaced by Cutoff?
- 2. *Utility* is replaced by *Eval*

Does it work in practice?

$$b^{m} = 10^{6}, b=35 \rightarrow m=4$$

4-ply lookahead is a hopeless chess player!

- 4-ply ≈ human novice
- 8-ply \approx typical PC, human master
- 12-ply ≈ Deep Blue, Kasparov

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Evaluation functions

• For chess, typically linear weighted sum of features

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + ... + w_n f_n(s)$$

• e.g., $w_1 = 9$ with

 $f_1(s) = \text{(number of white queens)} - \text{(number of black queens)}$, etc.

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Deterministic games in practice

- Checkers: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used a precomputed endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 444 billion positions.
- Chess: Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, who are too good.
- Go: human champions refuse to compete against computers, who are too bad. In go, b > 300, so most programs use pattern knowledge bases to suggest plausible moves.

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Summary

- Games are fun to work on!
- They illustrate several important points about AI
- perfection is unattainable → must approximate
- good idea to think about what to think about

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