

Formal Languages and Abstract Machines

Take Home Exam 2

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1 Context-Free Grammars (10 pts)

a) Give the rules of the Context-Free Grammars to recognize strings in the given languages where $\Sigma = \{a, b\}$ and S is the start symbol.

$L(G) = \{w \mid w \in \Sigma^*; |w| \geq 3;$
the first and the second from the last symbols of w are the same}

$$R = \{S \rightarrow aS \mid bS \mid aa \mid bb\}$$

$L(G) = \{w \mid w \in \Sigma^*; \text{the length of } w \text{ is odd}\}$ (2/10 pts)

$$R = \{S \rightarrow aA \mid bA, A \rightarrow aS \mid bS \mid e\}$$

$L(G) = \{w \mid w \in \Sigma^*; n(w, a) = 2 \cdot n(w, b)\}$ where $n(w, x)$ is the number of x symbols in w (3/10 pts)

$$R = \{S \rightarrow SaSaSbS \mid SaSbSbSaS \mid SbSaSaS \mid e\}$$

b) Find the set of strings recognized by the CFG rules given below: (3/10 pts)

$$\begin{aligned} S &\rightarrow X \mid Y \\ X &\rightarrow aXb \mid A \mid B \\ A &\rightarrow aA \mid a \\ B &\rightarrow Bb \mid b \end{aligned}$$

$Y \rightarrow CbaC$
 $C \rightarrow CC \mid a \mid b \mid \varepsilon$

$L(G) = \{w \mid w \in \Sigma^*; w \text{ is not } ab.\}$

2 Parse Trees and Derivations

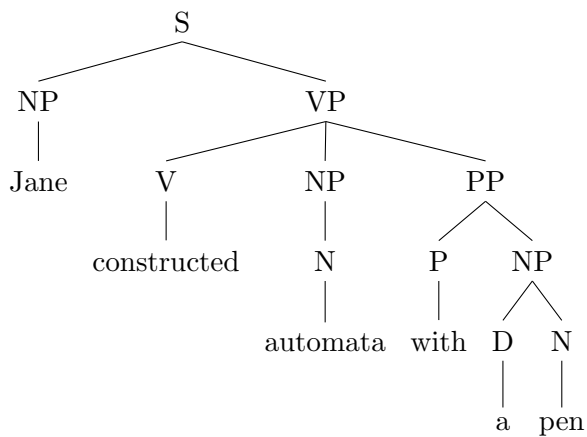
(20 pts)

Given the CFG below, provide parse trees for given sentences in **a** and **b**.

S → NP VP
VP → V NP | V NP PP
PP → P NP
NP → N | D N | NP PP
V → wrote | built | constructed
D → a | an | the | my
N → John | Mary | Jane | man | book | automata | pen | class
P → in | on | by | with

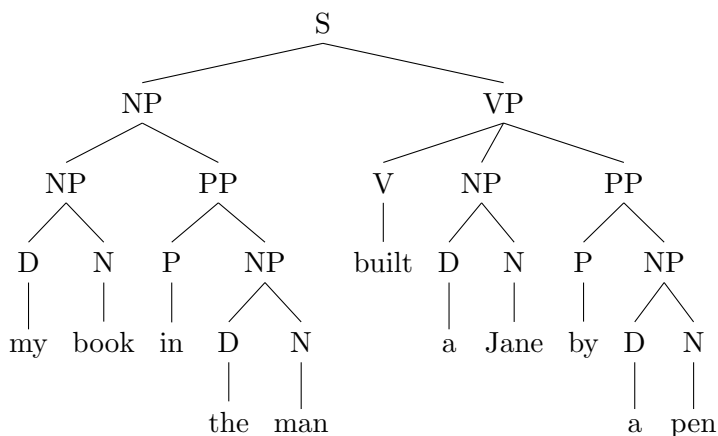
a) Jane constructed automata with a pen

(4/20 pts)



b) my book in the man built a Jane by a pen

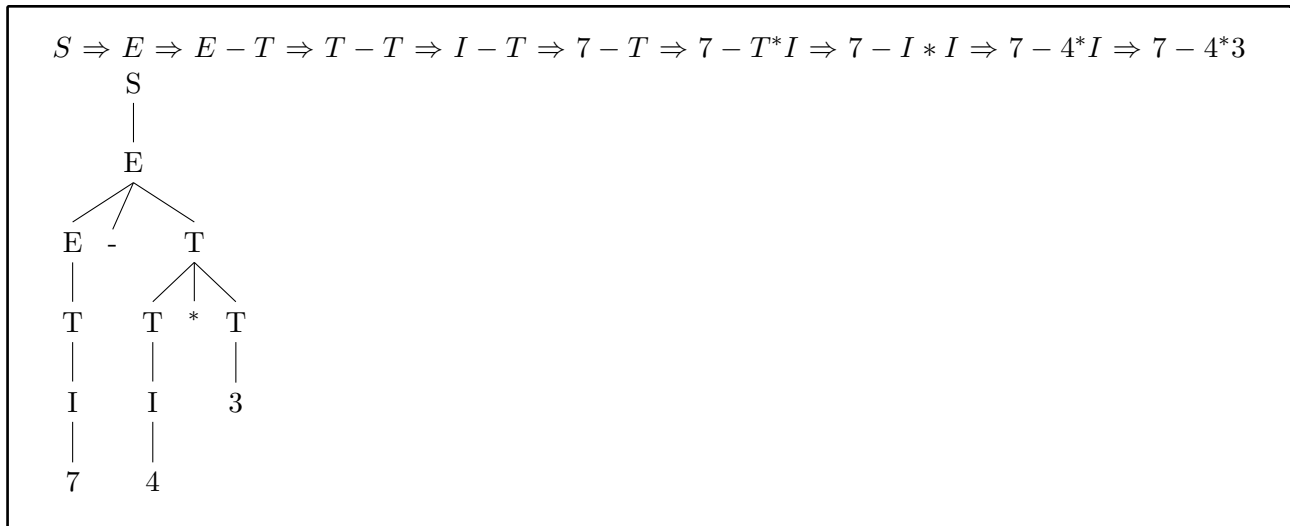
(4/20 pts)



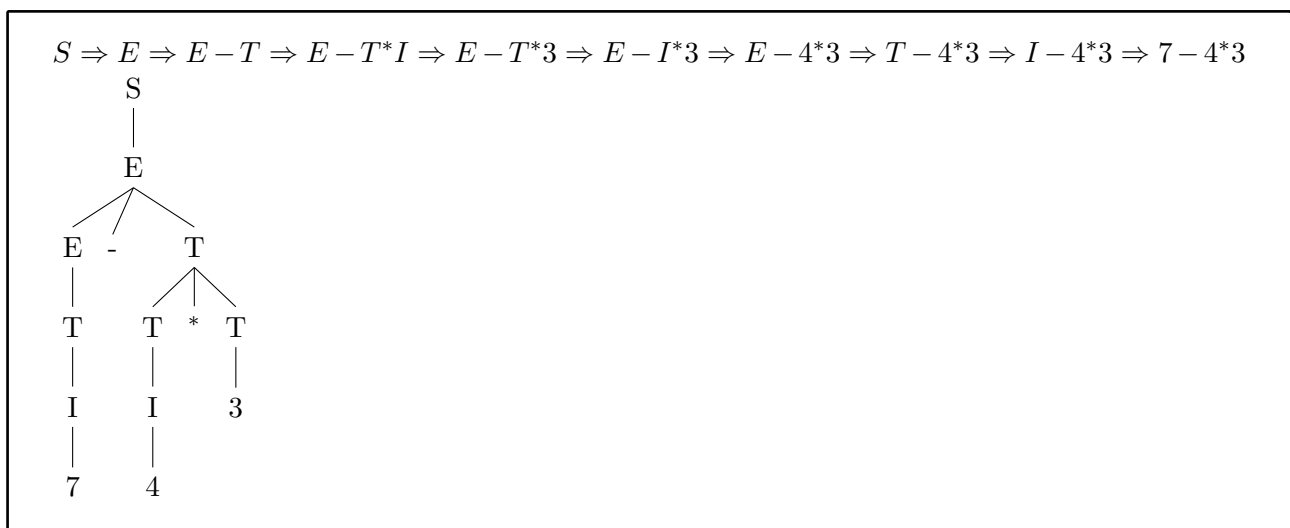
Given the CFG below, answer **c**, **d** and **e**

$S \rightarrow E$
 $E \rightarrow E + T \mid E - T \mid T$
 $T \rightarrow T * I \mid T / I \mid I$
 $I \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 6 \mid 7 \mid 8 \mid 9$

c) Provide the left-most derivation of $7 - 4 * 3$ step-by-step and plot the final parse tree matching that derivation (4/20 pts)



d) Provide the right-most derivation of $7 - 4 * 3$ step-by-step and plot the final parse tree matching that derivation (4/20 pts)



e) Are the derivations in **c** and **d** in the same similarity class? (4/20 pts)

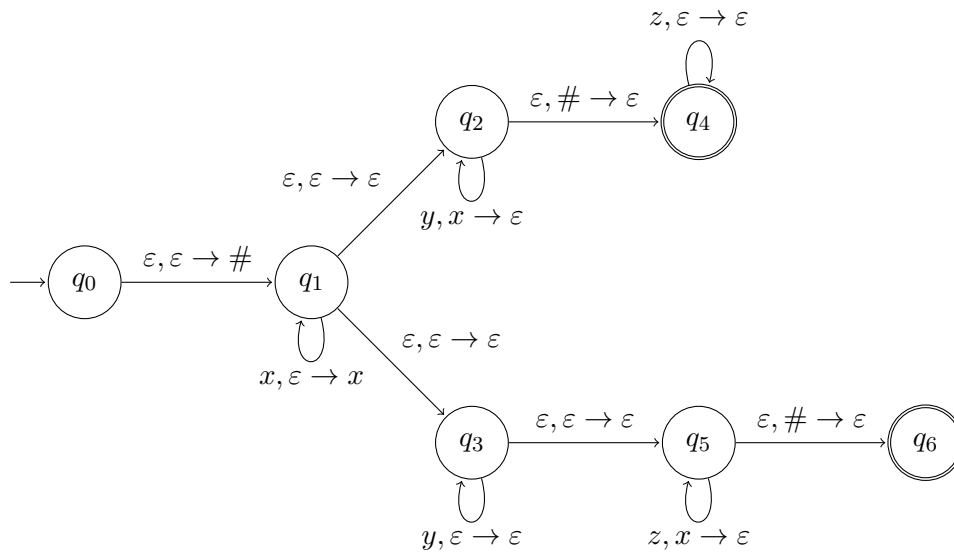
Yes, they are in the same similarity class because they represent applications of the same rules at the same positions in the strings only differing in the relative order of these applications.

3 Pushdown Automata

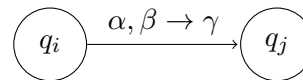
(30 pts)

a) Find the language recognized by the PDA given below

(5/30 pts)



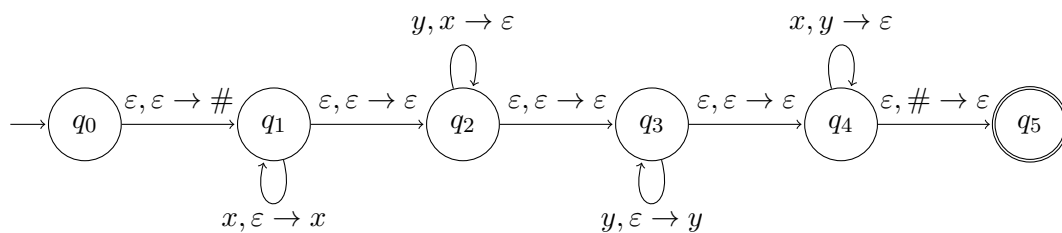
where the transition $((q_i, \alpha, \beta), (q_j, \gamma))$ is represented as:



The language $\{x^n y^n z^m \text{ or } x^n y^m z^n : m, n \geq 0\}$

b) Design a PDA to recognize language $L = \{x^n y^{m+n} x^m \mid n, m \geq 0; n, m \in \mathbb{N}\}$

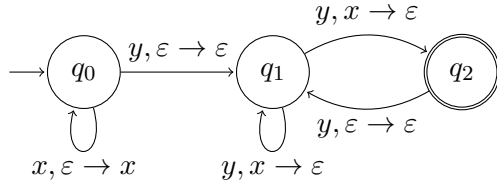
(5/30 pts)



c) Design a PDA to recognize language $L = \{x^n y^m \mid n < m \leq 2n; n, m \in \mathbb{N}^+\}$ (10/30 pts)

Do not use multi-symbol push/pop operations in your transitions.

Simulate the PDA on strings xy (with only one rejecting derivation) and $xyyyyy$ (accepting derivation) with transition tables.



$M = (K, \Sigma, \Gamma, \Delta, s, F)$, where $K = \{q_0, q_1, q_2\}$, $\Sigma = \{x, y\}$, $F = \{q_2\}$ and Δ is the following set of transitions;

- (1) $((q_0, x, e), (q_0, x))$
- (2) $((q_0, x, e), (q_1, e))$
- (3) $((q_1, y, x), (q_1, e))$
- (4) $((q_1, y, x), (q_2, e))$
- (5) $((q_2, y, e), (q_1, e))$

Table 1: transition table for xy

State	Unread Input	Stack	Transition Used
q_0	xy	ε	-
q_0	xy	x	1
q_0	y	xx	1
q_1	ε	xx	2

Table 2: transition table for $xyyyyy$

State	Unread Input	Stack	Transition Used
q_0	$xyyyyy$	ε	-
q_0	$xyyyyy$	x	1
q_0	$yyyy$	xx	1
q_1	yyy	xx	2
q_2	yy	x	4
q_1	yy	x	5
q_2	ε	ε	4

d) Given two languages L' and L as $L' = \{w \mid w \in L; |w| = 4n + 2 \text{ for } n \in \mathbb{N}\}$ (10/30 pts)

If L is a CFL, show that L' is also a CFL by constructing an automaton for L' in terms of another automaton that recognizes L .

answer here ...

4 Closure Properties

(20 pts)

Let L_1 and L_2 be context-free languages which are not regular, and let L_3 be a regular language. Determine whether the following languages are necessarily CFLs or not. If they need to be context-free, explain your reasoning. If not, give one example where the language is a CFL and a counter example where the language is not a CFL.

a) $L_4 = L_1 \cap (L_2 \setminus L_3)$

(10/20 pts)

b) $L_5 = (L_1 \cap L_3)^*$

(10/20 pts)

5 Pumping Theorem

(20 pts)

a) Show that $L = \{a^n m^n t^i \mid n \leq i \leq 2n\}$ is not a Context Free Language using Pumping Theorem for CFLs.

(10/20 pts)

answer here ...

b) Show that $L = \{a^n b^{2n} a^n \mid n \in \mathbb{N}^+\}$ is not a Context Free Language using Pumping Theorem for CFLs.

(10/20 pts)

answer here ...

6 CNF and CYK

(not graded)

a) Convert the given context-free grammar to Chomsky Normal Form.

$$S \rightarrow XSX \mid xY$$

$$X \rightarrow Y \mid S$$

$$Y \rightarrow z \mid \varepsilon$$

answer here ...

b) Use the grammar below to parse the given sentence using Cocke–Younger–Kasami algorithm. Plot the parse trees.

S → NP VP	VP → book include prefer
S → X1 VP	VP → Verb NP
X1 → Aux NP	VP → X2 PP
S → book include prefer	X2 → Verb NP
S → Verb NP	VP → Verb PP
S → X2 PP	VP → VP PP
S → Verb PP	PP → Prep NP
S → VP PP	Det → that this the a
NP → I she me Houston	Noun → book flight meal money
NP → Det Nom	Verb → book include prefer
Nom → book flight meal money	Aux → does
Nom → Nom Noun	Prep → from to on near through
Nom → Nom PP	

book the flight through Houston

Empty parse table:

<div> <div>1:5 → 1:1 2:5 1:5 → 1:2 3:5 1:5 → 1:3 4:5 1:5 → 1:4 5:5</div> </div>				
<div> <div>1:4 → 1:1 2:4 1:4 → 1:2 3:4 1:4 → 1:3 4:4</div> </div>		<div> <div>2:5 → 2:2 3:5 2:5 → 2:3 4:5 2:5 → 2:4 5:5</div> </div>		
<div> <div>1:3 → 1:1 2:3 1:3 → 1:2 3:3</div> </div>		<div> <div>2:4 → 2:2 3:4 2:4 → 2:3 4:4</div> </div>	<div> <div>3:5 → 3:3 4:5 3:5 → 3:4 5:5</div> </div>	
<div>1:2 → 1:1 2:2</div>		<div>2:3 → 2:2 3:3</div>	<div>3:4 → 3:3 4:4</div>	<div>4:5 → 4:4 5:5</div>
1:1	2:2	3:3	4:4	5:5
book	the	flight	through	Houston

rest of the answer here ...

7 Deterministic Pushdown Automata

(not graded)

Provide a DPDA to recognize the given languages, the DPDA must read its entire input and finish with an empty stack.

a) $a^*bc \cup a^nb^nc$

answer here ...

b) $(aa)^*c \cup a^nb^nc$

answer here ...