

# Ceng435 Term Project Part-1 Report

Zeynep Erdoğan  
2171577

Ayşenur Bülbül  
2171403

## I. INTRODUCTION

In this project we constructed a network that consists of 5 nodes(s, r1, r2, r3 and d) that are connected to each other. We were expected to develop a UDP socket application, find the link costs of links between nodes and figure out the shortest path between s and d. We used Python as programming language.

UDP(User Datagram Protocol) is a connectionless protocol. Packets may be lost during transportation. However, UDP is faster than TCP.

Sockets are used to send and receive messages by processes. We created sockets with Python using:

```
server = socket(AF_INET, SOCK_DGRAM)
```

AF\_INET is an Internet protocol version 4 family. SOCK\_DGRAM is a datagram based socket type. Since we constructed UDP based socket application, we used these arguments to create our sockets. Python provides two basic functions to construct communication between sockets. *sendto()* is for sending messages. This function is applied to a socket and takes 2 arguments: one is the message that we are going to send, the other is a tuple consisting of receiver's IP and sender's port. This function's usage is like this:

```
socket.sendto(msg, (IP, PORT))
```

The other basic function *recvfrom()* for receiving messages. We create server sockets. Bind the server's port and receiver's IP using *bind()* function. *recvfrom()* takes buffer size as argument and returns a message and address of sender.[1] These functions' usages are like this:

```
sock.bind((IP, PORT))
```

```
msg, addr = socket.recvfrom(1024)
```

Using Python we implemented a client/server application that sends and receives multiple messages at the same time assuming 1000 discovery messages are sent. Each node runs one script. Each router finds link costs considering RTT values between each node and writes the results to a file.

We synchronized each node's clock and we configure r1 and r2 with the given configure files for the part where we find the link costs. For the experiment part we configured s, r3 and d using *tc* and *netem* commands. We reset configurations before starting each experiment.

## II. DESIGN OF THE PROJECT

We are given a network which has a topology as specified below. The main goal of the project is to find shortest path

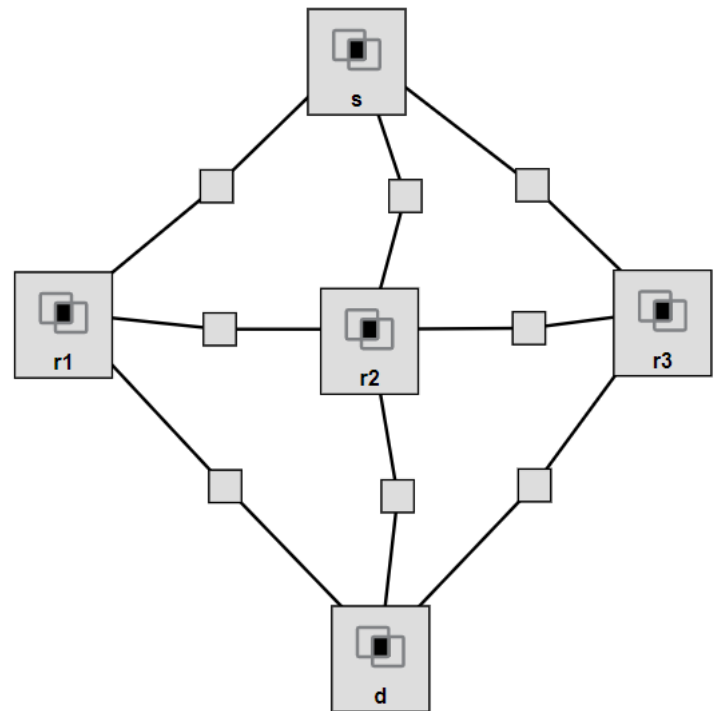


Fig. 1. Topology of the given network

from node s to d using RTT values between each node. Round-trip time (RTT) is the duration in milliseconds (ms) it takes for a network request to go from a starting point to a destination and back again to the starting point.

We designed the project such that nodes (r1, r2, r3) save the costs of the links between its neighbors. The node r1 keeps link costs of the s-r1, r1-r2 and r1-d. The node r2 keeps link costs of s-r2 and r2-d. The node r3 keeps link costs of s-r3, r3-r2 and r3-d.

To calculate RTT and keeping the values at nodes(r1, r2, r3), we have chosen node r1 as server for client nodes(s, r2, d), node r3 as server for the client nodes(s, r2, d), node r2 as server for client nodes(s, d). As understood, node s and node d are served only as clients, node r1 and r3 served only as servers and node r2 served as both client and server.

### III. IMPLEMENTATION OF THE PROJECT

Every node should be able to send and receive messages at the same time. Which means when a server is sending a message, it should be able to receive a message at the same time or when sending a message to a node, it should be able to send a message to another port at the same time. To achieve this we have used threads.

In the python script which runs in node s, there are 3 threads.

```
1 t1 = Thread(target=totalSource , args=([0]))
2 t2 = Thread(target=totalSource , args=([1]))
3 t3 = Thread(target=totalSource , args=([2]))
```

The totalSource function implements node s as the client for the server  $r_{i+1}$  where i is the given argument.

In the python script which runs in node r1 there are 3 threads.

```
1 t1 = Thread(target=r1Server , args=([0]))
2 t2 = Thread(target=r1Server , args=([1]))
3 t3 = Thread(target=r1Server , args=([2]))
```

The r1Server function implements node r1 as the server for the clients s, r2 and d(which we understood from the given argument i).

In the python script which runs in node r2 there are 4 threads.

```
1 t1 = Thread(target=r2Client , args=([1]))
2 t2 = Thread(target=r2Client , args=([2]))
3 t3 = Thread(target=r2Server , args=([0]))
4 t4 = Thread(target=r2Server , args=([3]))
```

The function r2Client implements node r2 as the clients for the servers r1 and r3. The function r2Server implements node r2 as the servers for the clients s and d.

In the python script which runs in node r3 there are 3 threads.

```
1 t1 = Thread(target=r3Server , args=([0]))
2 t2 = Thread(target=r3Server , args=([1]))
3 t3 = Thread(target=r3Server , args=([2]))
```

The r3Server function implements node r3 as the server for the clients s, r2 and d.

In the python script which runs in node d, there are 3 threads.

```
1 t1 = Thread(target=totalDest , args=([0]))
2 t2 = Thread(target=totalDest , args=([1]))
3 t3 = Thread(target=totalDest , args=([2]))
```

The totalDest function implements node d as the client for the server  $r_{i+1}$  where i is the given argument.

When implementing a server our main algorithm is

```
1 create a UDP socket.
2 while message left
3     start timer
4     send the message to client's IP through
   server's port.
5     wait for the feedback from the client
6     end timer
7     add passed time to totalTime
8 close socket
9 avg rtt = (totalTime/#messages)x1000 (in msec)
```

(We did not divide the total time with number of messages and then multiply it again with 1000 since we sent exactly 1000 messages.)

Since we are calculating RTT, starting the timer just before sending message and ending it right before feedback comes would gives us the total time.

When implementing a client our main algorithm is

```
1 create a UDP socket.
2 bind it to server's port
3 while message comes
4     receive message
5     send feedback
6 close socket
```

After running the scripts in each node, for 1000 messages, we have got the following average RTT values for each link;

TABLE I  
AVG RTT FOR LINKS(S-R1, R1-R2 AND R1-D)

|    | s       | r2       | d       |
|----|---------|----------|---------|
| r1 | 60.7022 | 140.7246 | 60.6613 |

TABLE II  
AVG RTT FOR LINKS(S-R2 AND R2-D)

|    | s       | d       |
|----|---------|---------|
| r2 | 80.6631 | 80.6530 |

TABLE III  
AVG RTT FOR LINKS(S-R3, R3-R2 AND R3-D)

|    | s      | r2     | d      |
|----|--------|--------|--------|
| r3 | 0.4819 | 80.632 | 0.4541 |

### IV. FINDING SHORTEST PATH USING DIJKSTRA ALGORITHM

Using Dijkstra Algorithm shortest path between certain nodes in graphs can be found.

The approach in Dijkstra Algorithm is, set distance from source to source, to 0. Then set other distances to  $\infty$ . Starting from the source traverse each neighbour and set their distance from  $\infty$  to distance from source. Then traverse from each neighbour of s to their neighbours. Change the distance if the path is shorter from previously found path. Apply this process until you traverse from every neighbour. We used this approach and get the following result:

|                 | s | r1       | r2       | r3       | d        |
|-----------------|---|----------|----------|----------|----------|
| min dist from s | 0 | $\infty$ | $\infty$ | $\infty$ | $\infty$ |
| min dist from s | 0 | $\infty$ | $\infty$ | 0.4819   | $\infty$ |
| min dist from s | 0 | $\infty$ | $\infty$ | 0.4819   | 0.4541   |
| min dist from s | 0 | 60.7022  | $\infty$ | 0.4819   | 0.4541   |
| min dist from s | 0 | 60.7022  | 80.6631  | 0.4819   | 0.4541   |

Then we concluded that the shortest path between s and d is s-r3-d.

## V. EXPERIMENT

We find the shortest path between *s* and *d* as *s-r3-d* using Dijkstra Algorithm. Thus, for experiment part, we modified our discovery files for *s*, *r3* and *d*. We changed *s* as server. It will send messages to *r3*. *r3* is client for *s* and server for *d*. *r3* waits for the message from *s* and then sends this message to *d*. *d* waits for the message from *r3*. All nodes runs one thread. This way we send all the messages to node *d*.

*s* sends the time of the start time of the sending process as a message to the *r3*. Then waits for feedback from *r3* to confirm message is taken from *r3*. *r3* sends the same message(starting time of the message sending) to *d*. Then waits for the feedback that confirms message is taken from *d*. In node *d*, we hold a variable *e2e* which stores total end-to-end delay. When *d* takes the message it takes the time the message is received and subtract it from the message that is sent from *r3*(which is start time of sending the message. We add this values to *e2e* variable 1000 times since we were expected to send 1000 messages.

To configure *s*, *r3* and *d* we find the interface *eth* between them using the commands that are explained in README.txt file. The *eth* interface between *s* and *r3* is *eth2* and between *d* and *r3* is *eth3*. We configured *r3*'s both *eth2* and *eth3*, *s*'s *eth2* and *d*'s *eth3* interface. We run the experiment scripts in the following order: *experimentD*, *experimentR3*, *experimentS*. The reasoning behind this is we run the clients first, they wait for the messages to be sent and then we run servers so that they have clients that are waiting for messages. The results of the experiments are 41.1603ms for 20ms emulation delay, 80.8184ms for 40ms emulation delay and 101.1108ms for 50ms emulation delay.

To find confidence interval, standard deviation is calculated. For %95 confidence interval the Z value is 1.960.  $error = 1.960 * (deviation/math.sqrt(1000))$ , since we send 1000 messages. Then we put this values in our MATLAB code to create graph. In our MATLAB code (*graph.m*), we create x axis values, y axis values and error values. Then plot the following graph.

## VI. GRAPH

The graph of emulation delay vs end-to-end delay is shown below.

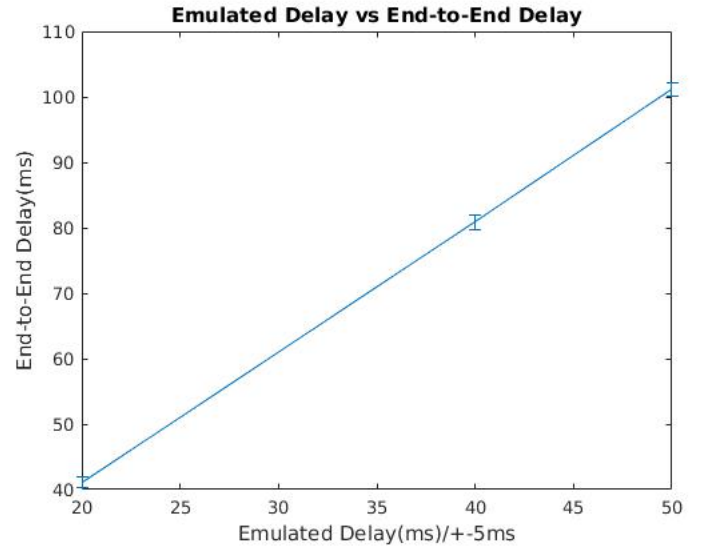


Fig. 2. Emulated Delay vs End-to-End Delay

When we examine the graph, we can see that emulation delay is doubled. Since we waited for feedback from each client that confirms the message is received this result makes sense. We sent the message from the link and waited for the feedback from the same link which calculates the time two times. This is why we get  $\sim 40ms$  for emulation delay 20ms,  $\sim 80ms$  for emulation delay 40ms and  $\sim 100ms$  for emulation delay 50ms. Delay is doubled up in end-to-end result.

## VII. CONCLUSION

In this part of the term project we implemented constructed a network that consists of 5 nodes that are connected to each other. We developed a UDP socket application, find the link costs of links between nodes using the files in *discoveryScripts* folder and figure out the shortest path between *s* and *d* using Dijkstra Algorithm. Then we modified our discovery files for experiment part. Configured each node (*s*, *d* and *r3*) and plotted the results. We see that end-to-end delay is 2 times of emulation delay. We learned UDP: How to send and receive messages, how to create sockets and bind them to IP and ports.

## REFERENCES

- [1] Retrieved from <https://wiki.python.org/moin/UdpCommunication>.
- [2] <https://www.cloudflare.com/learning/cdn/glossary/round-trip-time-rtt/>