

# Subqueries

# Where can a subquery go?

- Relational algebra syntax is so elegant that it's easy to see where subqueries can go.
- In SQL, a bit more thought is required ...

# Subqueries in a FROM clause

- In place of a relation name in the FROM clause, we can use a subquery.
- The subquery must be parenthesized.
- Must name the result, so you can refer to it in the outer query.

## Worksheet, Q1:

```
SELECT sid, dept || cnum as course, grade
FROM Took,
    (SELECT *
     FROM Offering
     WHERE instructor='Horton') Hofferings
WHERE Took.oid = Hofferings.oid;
```

- This FROM is analogous to:  
 $\text{Took} \times \rho_{\text{Hofferings}} (\llbracket \text{subquery} \rrbracket)$
- Can you suggest another version?

# Subquery as a value in a WHERE

- If a subquery is guaranteed to produce exactly one tuple, then the subquery can be used as a value.
- Simplest situation: that one tuple has only one component.

## Worksheet, Q2:

```
SELECT sid, surname
FROM Student
WHERE cgpa >
      (SELECT cgpa
       FROM Student
       WHERE sid = 99999);
```

- We can't do the analogous thing in RA:

$$\pi_{\text{sid, surname}} \sigma_{\text{cgpa} > (\text{«subquery»}) \text{Student}}$$

# Special cases

- What if the subquery returns **NULL**?
- What if the subquery could return more than one value?

# Quantifying over multiple results

- When a subquery can return multiple values, we can make comparisons using a quantifier.

- Example:

```
SELECT sid, surname
FROM Student
WHERE cgpa >
      (SELECT cgpa
       FROM Student
       WHERE campus = 'StG');
```

- We can require that
  - $cgpa >$  all of them, or
  - $cgpa >$  at least one of them.



# The Operator ANY

- Syntax:

$x \text{ «comparison» ANY («subquery»)$

or equivalently

$x \text{ «comparison» SOME («subquery»)$

- Semantics:

Its value is true iff the comparison holds for at least one tuple in the subquery result, i.e.,

$\exists y \in \text{«subquery results»} \mid x \text{ «comparison» } y$

- $x$  can be a *list* of attributes,  
but this feature is not supported by psql.

# The Operator ALL

- Syntax:

$x \text{ «comparison» ALL («subquery»)}$

- Semantics:

Its value is true iff the comparison holds for every tuple in the subquery result, i.e.,

$\forall y \in \text{«subquery results»} \mid x \text{ «comparison» } y$

- $x$  can be a list of attributes, but this feature is not supported by psql.

- Example: any-all

# The Operator IN

- Syntax:  
     $x \text{ IN } (\langle\langle \textit{subquery} \rangle\rangle)$
- Semantics:  
    Its value is true iff  $x$  is in the set of rows generated by the subquery.
- $x$  can be a list of attributes, and `psql` does support this feature.

## Worksheet, Q3:

```
SELECT sid, dept || cnum AS course, grade
FROM Took NATURAL JOIN Offering
WHERE
    grade >= 80 AND
    (cnum, dept) IN (
        SELECT cnum, dept
        FROM Took NATURAL JOIN Offering
            NATURAL JOIN Student
        WHERE surname = 'Lakemeyer' );
```

## Worksheet, Q4:

Suppose we have tables  $R(a, b)$  and  $S(b, c)$ .

1. What does this query do?

```
SELECT a  
FROM R  
WHERE b IN (SELECT b FROM S);
```

2. Can we express this query without using IN?

# The Operator EXISTS

- Syntax:  
**EXISTS** («*subquery*»)
- Semantics:  
Its value is true iff the subquery has at least one tuple.
- Read it as “exists a row in the subquery result”

# Example: EXISTS

```
SELECT surname, cgpa
FROM Student
WHERE EXISTS (
    SELECT *
    FROM Took
    WHERE Student.sid = Took.sid and
           grade > 85 );
```

## Worksheet, Q5:

```
SELECT instructor
FROM Offering Off1
WHERE NOT EXISTS (
    SELECT *
    FROM Offering
    WHERE
        oid <> Off1.oid AND
        instructor = Off1.instructor );
```



## Worksheet, Q6:

```
SELECT DISTINCT oid
FROM Took
WHERE EXISTS (
    SELECT *
    FROM Took t, Offering o
    WHERE
        t.oid = o.oid AND
        t.oid <> Took.oid AND
        o.dept = 'CSC' AND
        took.sid = t.sid );
```

# Scope

- Queries are evaluated from the inside out.
- If a name might refer to more than one thing, use the most closely nested one.
- If a subquery refers only to names defined inside it, it can be evaluated **once** and used repeatedly in the outer query.
- If it refers to any name defined outside of itself, it must be evaluated **once for each tuple in the outer query**.

These are called **correlated subqueries**.

# Renaming can make scope explicit

```
SELECT instructor
FROM Offering Off1
WHERE NOT EXISTS (
    SELECT *
    FROM Offering Off2
    WHERE
        Off2.oid <> Off1.oid AND
        Off2.instructor = Off1.instructor );
```

# Summary: where subqueries can go

- As a relation in a FROM clause.
- As a value in a WHERE clause.
- With ANY, ALL, IN or EXISTS in a WHERE clause.
- As operands to UNION, INTERSECT or EXCEPT.
- Reference: textbook, section 6.3.

# Modifying a Database

# Database Modifications

- Queries return a relation.
- A modification command does not; it changes the database in some way.
- Three kinds of modifications:
  - Insert a tuple or tuples.
  - Delete a tuple or tuples.
  - Update the value(s) of an existing tuple or tuples.

# Two ways to insert

- We've already seen two ways to insert tuples into an empty table:

```
INSERT INTO «relation» VALUES «list of tuples» ;
```

```
INSERT INTO «relation» («subquery») ;
```

- These can also be used to add tuples to a non-empty table.

# Naming attributes in INSERT

- Sometimes we want to insert tuples, but we don't have values for all attributes.
- If we name the attributes we *are* providing values for, the system will use `NULL` or a default for the rest.
- Convenient!



# Example

```
CREATE TABLE Invite (  
    name TEXT,  
    campus TEXT DEFAULT 'StG',  
    email TEXT,  
    age INT);
```

```
INSERT INTO Invite(name, email)  
(  
    SELECT firstname, email  
    FROM Student  
    WHERE cgpa > 3.4 );
```

Here, name and email get values from the query, campus gets the default value, and age gets **NULL**.

# Deletion

- Delete tuples satisfying a condition:

```
DELETE FROM «relation»  
WHERE «condition»;
```

- Delete all tuples:

```
DELETE FROM «relation»;
```

# Example 1: Delete Some Tuples

```
DELETE FROM Course
WHERE NOT EXISTS (
    SELECT *
    FROM Took JOIN Offering
              ON Took.oid = Offering.oid
    WHERE
        grade > 50 AND
        Offering.dept = Course.dept AND
        Offering.cnum = Course.cnum
);
```

# Updates

- To change the value of certain attributes in certain tuples to given values:

UPDATE *«relation»*

SET *«list of attribute assignments»*

WHERE *«condition on tuples»*;

# Example: update one tuple

- Updating one tuple:

```
UPDATE Student  
SET campus = 'UTM'  
WHERE sid = 999999;
```

- Updating several tuples:

```
UPDATE Took  
SET grade = 50  
WHERE grade >= 47 and grade < 50;
```