

part 1

a)

2, 4, 16, ..., n

k	2^k	i
0	1	2
1	2	4
2	4	16

$$2^k < n$$

$$\log_2(2^k) = \log n$$

$$2^k \cdot \log 2 = \log n$$

$$\log(2^k \cdot \log 2) = \log(\log n)$$

$$k \cdot \log(2 \cdot \log 2) = \log(\log n)$$

$$\boxed{\Theta k \leq \Theta \log(\log n)}$$

b) line 2 $\Theta(n)$

? line 3 will evaluate true

$\sum_{i=1}^n \sum_{j=1}^i \sum_{k=1}^j \Theta(n)$ times

$\sum_{i=1}^n \sum_{j=1}^i \Theta(n)$ will evaluate a max of i^3 times
so $\Theta(n^3)$ time

$$\Theta(n) + (\Theta(\sqrt{n})^3 + \Theta(2\sqrt{n})^3) (i = \sqrt{n}, i < n, i + \sqrt{n}) \sqrt{n} \rightarrow y$$

$$\Theta(n) + \Theta(\sqrt{n}^3 + j^3 + 2^3 + \dots + (\sqrt{n})^3) \sum_{k=1}^y x^3 \approx \frac{1}{4} y^2 (y+1)^2$$

$$1^3 + 2^3 + \dots + \sqrt{n}^3 \approx n^{7/2}$$

$$\Theta(n^{7/2}) = \frac{1}{4} y^4 + \Theta(y^3)$$

$$= \frac{1}{4} n^{4/2} + \Theta(n^{3/2}) = n^{7/2}$$

$$\boxed{\Theta(n^{7/2})}$$

$$(n) \log(n) \quad \boxed{\Theta(n^2)}$$

2)

line 5 $\Theta(n)$ (for loop)

line 7 if statement will run at most 2 times
 b/c will pass when $j \geq 10$ and then again when
 $j \geq 15$, but won't pass again b/c
 there will be an decimal value.

$$\sum_{i=0}^{n-1} (\Theta(1) + \Theta(2 \cdot \text{size}))$$

$$= \Theta(n) + \Theta(2n \cdot \text{size})$$

$$\Theta(2n \cdot \text{size}) \gg \Theta(n) \Rightarrow \boxed{\Theta = \Theta(n \cdot \text{size})}$$

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~~line 1~~ ~~line 2~~ ~~line 3~~

line 1 $\Theta(n)$ runtime (for loop)

line 2 $\Theta(n)$ runtime (for loop)

line 3 use Big O b/c we don't know how many iterations it will run

line 4 $m = m + m \Rightarrow \log \text{ run time} \Rightarrow \Theta(\log n)$

$$\sum_{i=1}^n \left(\sum_{k=1}^n (\Theta(n)) + O\left(\sum_{m=1}^n \left(\frac{1}{m}\right)\right) \right)$$

At most, will be evaluated every iteration or n times

$$= \sum_{i=1}^n (\Theta(n) + \sum (\Theta(\log n)))$$

$$= \sum_{i=1}^n (\Theta(n) + \sum (\Theta(\log Cn)))$$

$$= \sum_{i=1}^n (\Theta(n) + \sum_{i=1}^n (\Theta(\log Cn)))$$

$$= \sum_{i=1}^n (\Theta(n) + \frac{1}{2} \Theta(n \log Cn))$$

$$= 2 \cdot \Theta(n^2) + \Theta(n^2 \log Cn)$$

$$= \Theta(n^2 \log Cn) \gg \Theta(n^2)$$

$$= \Theta(n^2 \log Cn)$$