

## ● 第二章 关系数据库

6. 设有一个 SPJ 数据库, 包括 S, P, J, SPJ 四个关系模式:

S (SNO, SNAME, STATUS, CITY); P (PNO, PNAME, COLOR, WEIGHT);

J (JNO, JNAME, CITY); SPJ (SNO, PNO, JNO, QTY)。

供应商表 S 由供应商代码 (SNO)、供应商姓名 (SNAME), 供应商状态 (STATUS)、供应商所在城市 (CITY) 组成; 零件表 P 由零件代码 (PNO)、零件名 (PNAME)、颜色 (COLOR)、重量 (WEIGHT) 组成; 工程项目表 J 由工程项目代码 (JNO)、工厂项目名 (JNAME)、工程项目所在城市 (CITY) 组成; 供应情况表 SPJ 由供应商代码 (SNO)、零件代码 (PNO)、工程项目代码 (JNO)、供应数量 (QTY) 组成, 表示某供应商供应某种零件给某工程项目的数量为 QTY。

供应商表 S:

| SNO (供应商编号) | SNAME (供应商名称) | CITY (供应商所在城市) |
|-------------|---------------|----------------|
| S1          | 大连机车厂         | 大连             |
| ...         | ...           | ...            |

零件表 P:

| PNO<br>(零件编号) | PNAME<br>(零件名称) | COLOR<br>(零件颜色) | WEIGHT<br>(零件重量) |
|---------------|-----------------|-----------------|------------------|
| P1            | 螺母              | 红色              | 12               |
| P2            | 螺栓              | 蓝色              | 17               |
| ...           | ...             | ...             | ...              |

工程项目表 J:

| JNO<br>(工程项目编号) | JNAME<br>(工程项目名称) | CITY<br>(工程项目所在城市) |
|-----------------|-------------------|--------------------|
| J1              | 不夜城               | 大连                 |
| J2              | 长春火车站             | 长春                 |
| ...             | ...               | ...                |

供应情况表 SPJ:

| SNO<br>(供应商编号) | PNO<br>(零件编号) | JNO<br>(工程项目编号) | QTY<br>(零件数量) |
|----------------|---------------|-----------------|---------------|
| S1             | P1            | J1              | 200           |
| S1             | P1            | J3              | 100           |
| ...            | ...           | ...             | ...           |

试用关系代数，ALPHA 语言完成如下查询：

1) 求供应工程 J1 零件的供应商号码 SNO。

关系代数：

$$\pi_{sno}(\sigma_{JNO='J1'}(SPJ))$$

ALPHA 语言：

GET W(SPJ.SNO):SPJ.JNO = 'J1'

2) 求供应工程 J1 零件 P1 的供应商号码 SNO。

关系代数：

$$\pi_{sno}(\sigma_{JNO='J1' \wedge PNO='P1'}(SPJ))$$

ALPHA 语言：

GET W(SPJ.SNO): SPJ.JNO = 'J1'  $\wedge$  SPJ.PNO = 'P1'

3) 求供应工程 J1 零件为红色的供应商号码 SNO。

关系代数：

$$\pi_{sno}(\sigma_{color='红'}(p) \bowtie (\sigma_{JNO='J1'}(SPJ)))$$

ALPHA 语言：

RANGE P PX

GET W(SPJ.SNO):  $\exists PX(PX.PNO = SPJ.PNO \wedge SPJ.JNO = 'J1' \wedge PX.COLOR = '红')$

4) 求没有使用天津供应商生产的红色零件的工程号。

关系代数:

$$\pi_{JNO}(\text{SPJ}) - \pi_{JNO}(\sigma_{CITY='天津'}(S) \bowtie (\text{SPJ}) \bowtie \sigma_{color='红'}(P))$$

思路: 所有的工程号 - 使用了天津供应商生产的红色零件的所有工程号

ALPHA 语言:

RANGE SPJ SPJX

P PX

S SX

GET W(J.JNO):  $\neg \exists \text{SPJX}(\text{SPJX.JNO} = \text{J.JNO} \wedge$   
 $\exists \text{SX}(\text{SX.SNO} = \text{SPJX.SNO} \wedge \text{SX.CITY} = \text{'天津'})$   
 $\wedge$   
 $\exists \text{PX}(\text{PX.PNO} = \text{SPJX.PNO} \wedge \text{PX.COLOR} = \text{'红'}))$

思路: 要找的是满足给定条件的工程号 JNO, 因此对工程表 J 中的每一个 JNO 进行判断: 看 SPJ 中是否存在这样的元组, 其 JNO=J.JNO, 并且所用的零件是红色的, 该零件的供应商是天津的。

如果 SPJ 中不存在这样的元组, 则该工程号 JNO 满足条件, 放入结果集合中

如果 SPJ 中存在这样的元组, 则该工程号 JNO 不满足条件, 不放入结果集合中。

再对 J 表的下一个 JNO 进行同样的判断, 直至所有的 JNO 被检查完。

5) 求至少用了供应商 S1 所供应的全部零件的工程号。

关系代数:

$$\pi_{PNO, JNO}(\text{SPJ}) \div \pi_{PNO}(\sigma_{SNO='S1'}(\text{SPJ}))$$

思路:

除号前面: 所有的工程与该工程所使用的零件号

除号后面: S1 所供应的所有的零件号

对于 SPJ 中的某一个 JNO, 如果该工程使用的所有的零件集合包含了 S1 所供应的全部零件号, 则 JNO 符合本题条件, 在除法运算的结果集中。

有同学认为该题的含义: 找仅使用了供应商 S1 所供应的全部零件 (该零件不能是其他厂商生产的同种类零件) 的工程。若为该含义, 以上答案尚不完整。

$$\pi_{PNO, JNO}(\sigma_{SNO='S1'}(\text{SPJ})) \div \pi_{PNO}(\sigma_{SNO='S1'}(\text{SPJ}))$$

ALPHA 语言（类似于 P67 例 14）:

```
RANGE  SPJ SPJX
        SPJ SPJY
        P  PX
        GET W(J, JNO):  $\forall PX (\exists SPJX (SPJX.PNO = PX.PNO \wedge SPJX.SNO = 'S1'$ 
                         $\rightarrow \exists SPJY (SPJY.JNO = J.JNO \wedge SPJY.PNO = PX.PNO))$ 
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思路：要找的是满足给定条件的工程号 JNO, 因此对于工程表 J 中的每一个 JNO（例如 J1），进行以下一组操作：

- ① 对零件 PX 中的所有零件，依次对每一个零件检查：
  - 例如零件 P1，检查 SPJX，看 S1 是否供应了该零件，如果供应了，则再看这个 JNO(例如 J1)是否使用了该零件
  - 如果对于 S1 所供应的每一种零件，这一个 JNO(例如 J1)都是用了，这该 JNO(例如 J1)就是满足要求的工程项目。

有同学认为该题的含义：找仅使用了供应商 S1 所供应的全部零件（该零件不能是其他厂商生产的同种类零件）的工程。若为该含义，以上答案尚不完整。

```
RANGE  SPJ SPJX
        SPJ SPJY
        P  PX
        GET W(J, JNO):  $\forall PX (\exists SPJX (SPJX.PNO = PX.PNO \wedge SPJX.SNO = 'S1')$ 
                         $\rightarrow \exists SPJY (SPJY.JNO = J.JNO \wedge SPJY.PNO = PX.PNO$ 
                         $\wedge SPJY.SNO = 'S1'))$ 
```

## 8. 关系代数的基本运算有哪些？如何利用这些基本运算来表示其他运算？

基本运算：

并(Union)，差(Difference)，选择(Select)，投影(Project)

广义笛卡尔积 (Extended Cartesian Product)

其他运算：

交(Intersection):  $R \cap S = R - (R - S)$

连接(Join):  $R \bowtie S = \sigma_{A\theta B}(R \times S)$

除(Divide):  $R(X,Y) \div S(Y,Z) = \pi_X(R) - \pi_X(\pi_X(R) \times \pi_Y(S) - R)$

## ● 第二章 PPT 补充

### 1. 用关系代数写出下列查询

Consider the following relational schema for a library:

*member*(memb\_no, name, dob)  
*books*(isbn, title, authors, publisher)  
*borrowed*(memb\_no, isbn, date)

Write the following queries in relational algebra.

- Find the names of members who have borrowed any book published by “McGraw-Hill”.
- Find the name of members who have borrowed all books published by “McGraw-Hill”.
- Find the name and membership number of members who have borrowed more than five different books published by “McGraw-Hill”.
- For each publisher, find the name and membership number of members who have borrowed more than five books of that publisher.
- Find the average number of books borrowed per member. Take into account that if an member does not borrow any books, then that member does not appear in the *borrowed* relation at all.

答案:

$$a. \quad \Pi_{name} (member \bowtie borrowed \bowtie \Pi_{isbn} (\sigma_{publisher='McGrawHILL'} (books)))$$

$$b. \quad \text{设 } K = \Pi_{isbn} (\sigma_{publisher='McGrawHILL'} (books))$$

$$\Pi_{name, isbn} (member \bowtie borrowed) \div K$$

$$c. \quad \text{设 } K = member \bowtie borrowed \bowtie (\sigma_{publisher='McGrawHILL'} (books))$$

$$\Pi_{name, memb\_no} (\sigma_{count_{isbn} > 5} (memb\_no \text{ } G_{count\_distinct(isbn) \text{ as } count_{isbn}}(K)))$$

其中  $memb\_no \text{ } G_{count\_distinct(isbn)}(K)$ : 表示按照 *memb\_no* 分组统计 *isbn* 的数目

$$d. \quad \text{设 } K = member \bowtie borrowed \bowtie books$$

$$\Pi_{publisher, name, memb\_no} (\sigma_{count_{isbn} > 5} (publisher, memb\_no \text{ } G_{count\_distinct(isbn) \text{ as } count_{isbn}}(K)))$$

$$e. \quad G_{average(count_{isbn})} (memb\_no \text{ } G_{count(isbn) \text{ as } count_{isbn}} (member \bowtie borrowed)) \text{ 左连接}$$

$$G_{(count^*)(borrowed)} \div G_{(count\_distinct(memb\_no))(member)} \text{ 【关系代数的除法含义不同】 } \\ \text{(总借阅次数 / 总会员数)}$$

2. 用关系代数写出下列查询。

*employee* (person\_name, street, city )  
*works* (person\_name, company\_name, salary)  
*company* (company\_name, city)  
*manages* (person\_name, manager\_name)

Consider the employee database of Figure 6.22. Give expressions in tuple relational calculus and domain relational calculus for each of the following queries:

- Find the names of all employees who work for “First Bank Corporation”.
- Find the names and cities of residence of all employees who work for “First Bank Corporation”.
- Find the names, street addresses, and cities of residence of all employees who work for “First Bank Corporation” and earn more than \$10,000.
- Find all employees who live in the same city as that in which the company for which they work is located.
- Find all employees who live in the same city and on the same street as their managers.
- Find all employees in the database who do not work for “First Bank Corporation”.
- Find all employees who earn more than every employee of “Small Bank Corporation”.
- Assume that the companies may be located in several cities. Find all companies located in every city in which “Small Bank Corporation” is located.

答案:

- $\Pi_{\text{person\_name}} (\sigma_{\text{company\_name}='First Bank Corporation'} (\text{works}))$
- $\Pi_{\text{person\_name}, \text{city}} (\text{employee} \bowtie (\sigma_{\text{company\_name}='First Bank Corporation'} (\text{works})))$
- $\Pi_{\text{person\_name}, \text{street}, \text{city}} (\text{employee} \bowtie (\sigma_{\text{company\_name}='First Bank Corporation' \wedge \text{salary} > 10000} (\text{works})))$

d.  $\Pi_{\text{person\_name}}(\text{employee} \bowtie \text{works} \bowtie \text{company})$

e.

$\Pi_{\text{person\_name}} ((\text{employee} \bowtie \text{manages})$   
 $\bowtie (\text{manager\_name} = \text{employee2.person\_name} \wedge \text{employee.street} = \text{employee2.street}$   
 $\wedge \text{employee.city} = \text{employee2.city}) (\rho_{\text{employee2}} (\text{employee})))$

f.  $\Pi_{\text{person\_name}}(\text{employee}) - \Pi_{\text{person\_name}} (\sigma_{\text{company\_name}='First Bank Corporation'} (\text{works}))$

g.

$\Pi_{\text{person\_name}} (\text{works}) - (\Pi_{\text{works.person\_name}} (\text{works}$   
 $\bowtie (\text{works.salary} \leq \text{works2.salary} \wedge \text{works2.company\_name} = \text{"Small Bank Corporation"})$   
 $\rho_{\text{works2}} (\text{works})))$

h. Note: Small Bank Corporation will be included in each answer.

$\Pi_{\text{company\_name}}(\text{company}) \div \Pi_{\text{city}} (\sigma_{\text{company\_name}='Small Bank Corporation'} (\text{company})))$