DES227: Algorithms Design Final Mock Exam

curated by The Peanuts

Name Nonpravid I. ID. 66227724422 Section Seat No. ...

Conditions: Semi-Closed Book (A4 Both Sides)

Directions:

- 1. This exam has 14 pages (including this page).
- 2. Calculators (Casio 991 Series) are allowed.
- 3. Any attempt to bring a quantum computer will be considered... impressive, but still cheating.
- 4. Write your name clearly at the top of each page.
- 5. Stay calm. Breathe in. Breathe through. Breathe deep. Breathe out.
- 6. No brute-force cheating. We will prune you with a branch-and-bound penalty.
- 7. May your answers be correct, your logic flawless, and your pencil always sharp.



Multiply $(100)_2$ with $(01)_2$ using the Recursive-Multiply algorithm. Assume the base case is when both x and y are 2-digit binary numbers. Show your complete calculation steps.



Consider the array: [5, 2, 8, 1, 9, 3, 7, 4]

Apply the divide-and-conquer approach (Sort-and-Count algorithm) to count the number of inversions in this array. Show all steps of the recursive calls and how the "across-half" inversions are counted during the merge steps.

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 \begin{bmatrix} \mathsf{S}_{9}2_{9}\mathsf{S}_{9}1_{9}q_{9}3_{9}q_{9}4 \end{bmatrix} \\ \mathsf{left} \ \mathsf{half} \ \left[ \mathsf{S}_{9}2_{9}\mathsf{S}_{9}q_{1} \right] \xrightarrow{\mathsf{right}} \left[ \mathsf{S}_{9}1 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}5 \right] \\ \mathsf{right} \ \mathsf{half} \ \left[ \mathsf{S}_{9}2_{9}\mathsf{S}_{9}q_{1} \right] \xrightarrow{\mathsf{right}} \left[ \mathsf{S}_{4}1 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{right} \ \mathsf{half} \ \left[ \mathsf{S}_{9}2_{9}\mathsf{Sorti:} \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{right} \ \mathsf{half} \ \left[ \mathsf{S}_{9}2_{9}\mathsf{Sorti:} \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{Sorti:} \left[ \mathsf{S}_{4}9 \right] \\ \mathsf{sorti:} \left[ \mathsf{S}_{4}9 \right] \xrightarrow{\mathsf{Sorti:}} \left[ \mathsf{Sorti:} \left[
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Use dynamic programming to solve the following coin-changing problem:

Given coins of denominations {1, 3, 4, 5} and an amount of 7, find the minimum number of coins needed to make this amount. Show your work by filling in the complete DP table and explaining your approach.

	\$0	\$1	\$2	\$3	\$4	\$ 5	\$ 6	\$7
d1 = 1	0	1	2	3	4	5	6	7
d ₂ = 3	0	1	2	1	2	3	2	3
d3 = 4	0	1	2	1	1	2	V 2	2
d4 = 5	0	1	2	1	1	1	2	2

Starting here

tebes only two coins of deads}

Solve the following 0/1 knapsack problem:

item	weight	value
1	3	\$25
2	2	\$20
3	4	\$40
4	5	\$50

$$W = 7$$

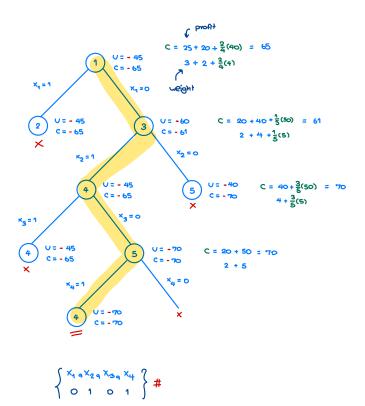
(a) Solve this problem using the dynamic programming approach. Show the complete DP table.

			0	1	2	3	4	5	6	7
Р	ω	0	(·)	0	0	0	0	0	0	0
25	3	1	P	0	0	25	٤ŷ	25	25	25
20	2	2	0	0	8	25	25	4	4 5	45
40	4	3	0	0	20	25	40	45	60	65
50	5	4	0	0	(20)	25	40	8	8	70

(b) Solve the same problem using the branch-and-bound algorithm. Draw the state space tree and show the calculation of upper bounds at each node.

item	weight	value
1	3	\$25
2	2	\$20
3	4	\$40
4	5	\$50

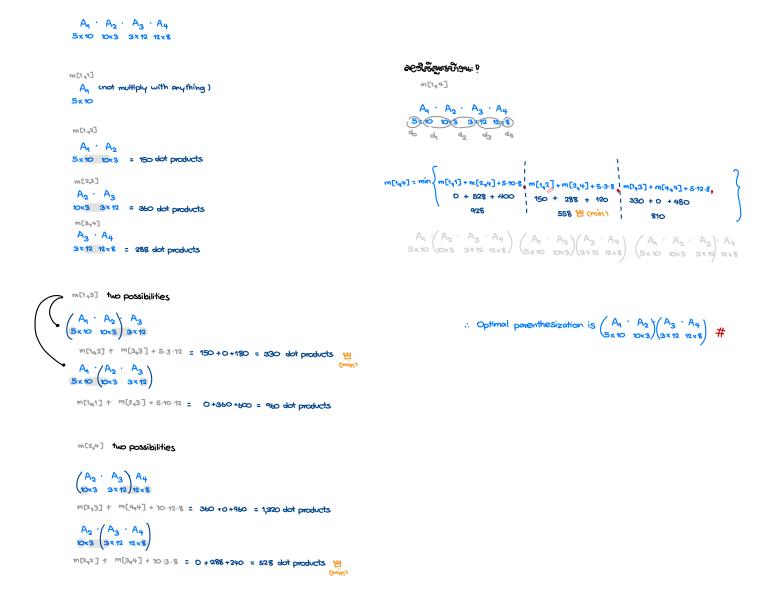
$$W = 7$$



Complete the given table m[i,j] to solve the matrix chain multiplication problem by dynamic programming. There are 4 matrices with dimensions as follows: $M_1: 5\times 10,\ M_2: 10\times 3,\ M_3: 3\times 12,\ {\rm and}\ M_4: 12\times 8.$

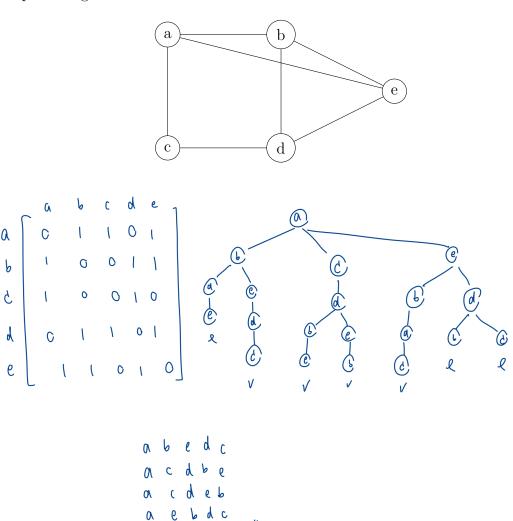
m[i,j]	1	2	3	4	
1	0	150	330	<i>5</i> 58	K=2
2		0	360	528	
3			0	288	
4				0	

After completing the table, determine the optimal parenthesization for multiplying these matrices.





Apply backtracking to find a Hamiltonian circuit in the following graph. Start from vertex a and show your search tree, including promising and non-promising nodes.

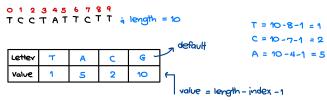


Use Horspool's algorithm to search for the pattern "TCCTATTCTT" in the DNA sequence:

$\hbox{``TTATAGATCTCGTATTCTTTATAGATCTCCTATTCTT''}$

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TCCTATTCTT

TCCTATTCTT
```

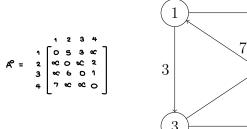


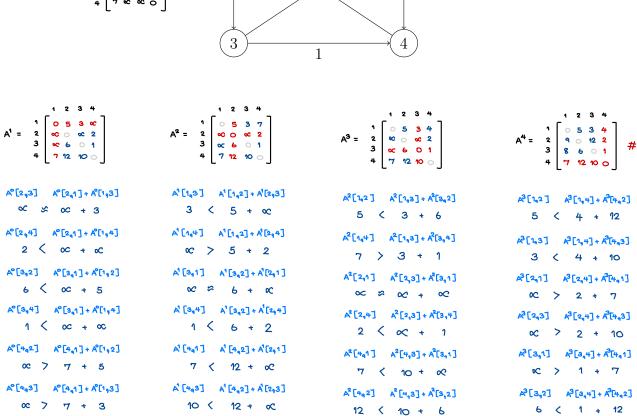
Apply Floyd-Warshall algorithm to find the shortest path distances between all pairs of vertices in the following weighted directed graph. Show the intermediat1e matrices after each iteration.

5

6

2

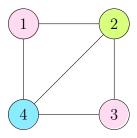




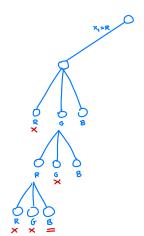
How many comparisons does the brute force string matching algorithm make when searching for the pattern "124" in the text "6125212901"?

```
6 1 2 5 2 1 2 9 0 1
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
1 2 4
```

Consider the following graph with 4 vertices:

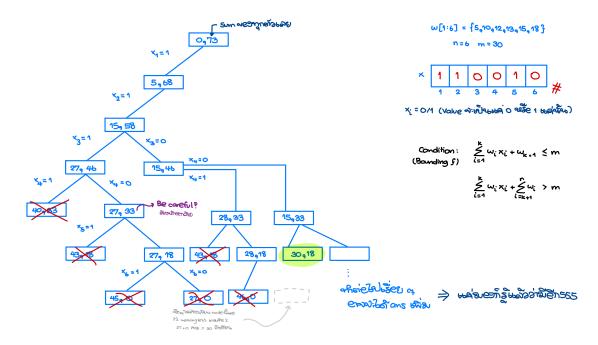


m=3 {RqGqB} Apply a graph coloring algorithm to find at least one valid coloring for this graph using the 3 number of colors. Explain your approach and show your work.



.: One of the solutions is RGRB #

Apply backtracking to solve the following instance of the subset sum problem: A = 5, 10, 12, 13, 15, 18 and target sum d = 30. Find at least one subset that sums to exactly 30. Show your state space tree.



For each of the following statements, indicate whether it is True or False. No need to justify your answer..

(a) Every problem in NP is also in P. (assumudae)

(a) True (b) False.

(b) Taylor Swift was born in 1989.

(a) True (b) False.

(c) Branch-and-bound algorithms always produce optimal solutions faster than brute-force approaches.

(a) True (b) False.

(d) If a problem is NP-complete, then there exists a polynomial-time algorithm to verify a solution, but no known polynomial-time algorithm to find a solution.

(a) True (b) False.

(e) The Floyd-Warshall algorithm can handle negative edge weights but not negative cycles.

(a) True (b) False.