## CSS321: Theory of Computation Midterm Mock Exam

#### curated by The Peanuts

Conditions: Open Book

#### **Directions:**

- 1. This exam has 16 pages (including this page).
- 2. You may use a calculator, but it won't help you prove languages are non-regular.
- 3. Dictionaries are not allowed. Neither is asking the Pumping Lemma for help (it's not here!).
- 4. Cheating is strictly prohibited.
- 5. Good luck! May all your states be accepting.

The solution will never be released, sorry!

Consider the following statements:

- (a)  $\{a\} \in \{\{a\}, \{b\}\}\$
- (b)  $\{a,b\} \subseteq \{\{a\},\{b\},a,b\}$
- (c)  $2^{\{a,b\}} = \{\emptyset, \{a\}, \{b\}, \{a,b\}\}$
- (d) For any sets A and B, if  $A \subseteq B$  then  $2^A \subseteq 2^B$
- (e)  $A \times (B \cup C) = (A \times B) \cup (A \times C)$

Which of the above statements are true?

Let  $A = \{1, 2, 3\}$  and  $R = \{(1, 1), (1, 2), (2, 2), (2, 3), (3, 1), (3, 3)\}$  be a relation on A. Consider the following properties:

- (a) R is reflexive
- (b) R is symmetric
- (c) R is transitive
- (d) R is antisymmetric
- (e) R is a partial order

Which of the above properties hold for R?

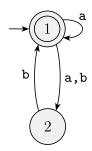
Consider the following statements about regular languages:

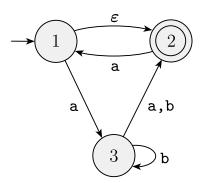
(a) 
$$L((a \cup b)^*a) = \{w \in \{a,b\}^* \mid w \text{ ends with } a\}$$

- (b)  $L(a^*b^*) \cap L(b^*a^*) = \{a^nb^n \mid n \ge 0\}$
- (c) For any regular language  $L, L^* = L^+ \cup \{\varepsilon\}$
- (d) The language  $\{a^nb^m\mid n\neq m\}$  is regular
- (e) Every finite language is regular

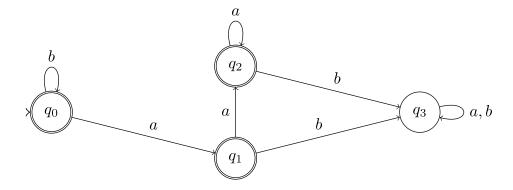
Which of the above statements are false?

Convert the following two nondeterministic finite automata to equivalent deterministic finite automata.





Consider the following finite automaton A over the alphabet  $\Sigma = \{a, b, c\}$ .



a) Is A deterministic? If not, convert A into a DFA.

b) Is A minimal? If not, convert A into a minimal DFA.

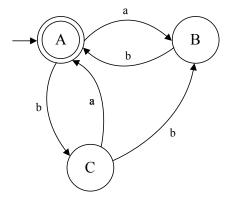
c) Convert A into a regular expression.

Prove by induction that  $n^3 + (n+1)^3 + (n+2)^3$  is divisible by 3 and n > 0

Construct a DFA equivalent to the NFA  $M=(\{a,b,c,d\},\{0,1\},\delta,\,a,\{b,d\})$  where  $\delta$  is given below and informally describe the language it accepts.

δ	0	1
a	$\{b,d\}$	$\{b\}$
b	$\{c\}$	$\{b,c\}$
c	$\{d\}$	$\{a\}$
d	Ø	$\{a\}$

Find the regular expression for the following DFA.



Show that  $(\emptyset)^* = \epsilon$  for regular expression

Let  $\Sigma = \{a, b\}$  and let  $L_1$  be the language over  $\Sigma$  given by the regular expression  $(ab \cup ba)^*$ . Design a DFA for  $L_1$ .

Let  $\Sigma = \{a, b\}$  and let  $L_2 = \{w \in \Sigma^* | \text{w does not contain bbb as a substring}\}$ . Design a DFA for  $L_2$  and write a regular expression.

Consider the following statements about cardinality and functions:

- (a) Every subset of a countably infinite set is finite or countably infinite
- (b) There exists a bijection from  $\mathbb{N}$  to  $\mathbb{N} \times \mathbb{N}$
- (c) The set  $2^{\mathbb{N}}$  is countably infinite
- (d) If  $f:A\to B$  is one-to-one and |A|=|B|, then f is onto
- (e) The diagonalization principle can prove that some infinite sets have different cardinalities

Which statements are true?
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Let  $A = \{a, b, c\}$  and  $R = \{(a, a), (a, b), (b, c), (c, a)\}$  be a relation on A. a) Find the smallest reflexive relation  $R_1$  containing R.

- b) Find the smallest reflexive and transitive relation  $R_2$  containing R.
- c) Is  $R_2$  an equivalence relation? If not, what would you need to add to make it one?
- d) For the relation R, find:
  - 1. The row set of a:  $R_a = \{x \in A \mid (a, x) \in R\} =$ \_\_\_\_\_\_
  - 2. The diagonal set:  $D = \{x \in A \mid (x, x) \notin R\} =$ \_\_\_\_\_

Construct DFAs for the following languages over  $\Sigma = \{0, 1\}$ : a)  $L_1 = \{w \mid |w| \mod 3 = 0\}$  (strings whose length is divisible by 3)

b)  $L_2 = \{ w \mid w \text{ contains an even number of 1's} \}$ 

c)  $L_3 = \{ w \mid w \text{ ends with } 01 \}$ 

Prove the following set theory identity:  $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$