

Recursive evaluation of x^n

Evaluate x^n where n is a nonnegative integer

Iterative Solution

Multiply x by itself n times

$O(n)$ algorithm

```
double power(double x, int n) {
    if (n == 0)
        return 1;
    else {
        double p = power(x, n/2);
        if (n % 2 == 0)
            return p * p;
        else
            return x * p * p;
    }
}
```

Recursive Solution

$x^0 = 1$
 $x^n = (x^{n/2})^2$ $n > 0$ and even
 $x^n = x (x^{(n-1)/2})^2$ $n > 0$ and odd

$O(\log_2 n)$ algorithm

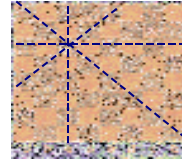
The 8 Queens Problem

Arrange 8 queens on a standard 8×8 chessboard so that none of the queens attack each other.

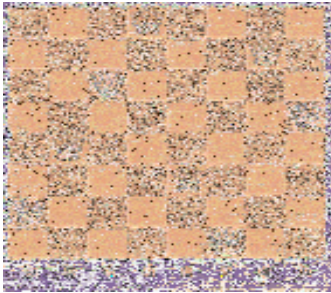
More general problem is to arrange n ($n \geq 1$) queens on an $n \times n$ chessboard so that none of the queens attack each other.

A queen placed at a particular square attacks any other queen placed

- in the **same row**
- in the **same column**
- along the **two diagonals** passing through the square

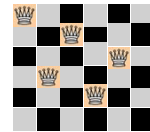
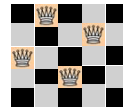


A Particular Solution to the 8 Queens Problem



This is one of 92 solutions to the problem. 12 of these solutions are distinct; the others are rotations and reflections of each other.

Solution to n Queens Problem



Solution?

Assume that queens have been correctly placed in rows $1 \dots n-1$. Then look at each column in row n to find where to place the queen in row n .

Correctness?

Base case : Yes - empty chessboard
 Smaller caller : Yes - problem reduces to base case
 General case : No - no guarantee that a solution for $n-1$ rows leads to one for n rows, or even that a solution exists at all!

Backtracking

If no solution exists for row n , "backtrack" to row $n-1$ and try to find a new column in which to place the queen, and then try row n again. If no column is suitable in row $n-1$, backtrack to row $n-2$, etc. This is now guaranteed to find a solution if one exists.

Recursive Backtracking Algorithm for n Queens Problem

Algorithm *addQueen(row)*

Assume that queens have been validly placed in rows $1 \dots n-1$. Then add a validly placed new queen to the board in row n , if this is possible. If not, backtrack to previous row(s) and try again.

```
for col = 1 to 8
    if no queen attacks square (row,col) {
        place a queen at (row,col)

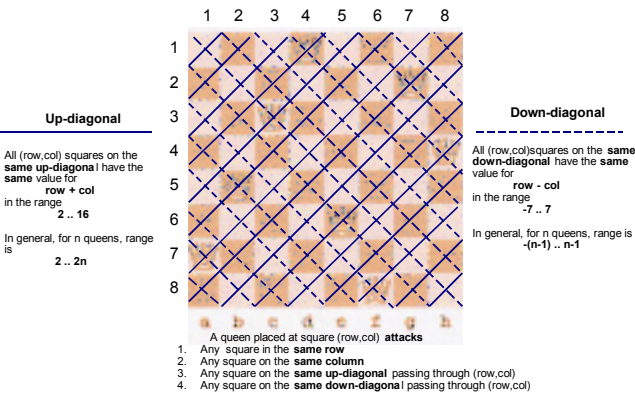
        if row = 8
            display positioned queens    // solution found!
        else
            addQueen(row+1)              // try and position next queen in next row

        remove the queen at (row,col)    // either we couldn't position in next row and we
                                         // need to backtrack, or we found a solution and
                                         // we want to backtrack anyway to find
                                         // another one
    }
}
```

Problems

- How do we represent the positions of the queens on the board?
- How do we determine whether two queens are attacking each other, i.e. are
 - in the same row or column or
 - on the same diagonal?

Squares Queen Attacks



Position Representation

Use *PArrays* since array lower limits nonzero.

PArray	Range	Purpose
queenCol	1 .. n	Store positions of queens
colStatus	1 .. n	Whether queen is located in col j
upStatus	2 .. 2n	Whether queen is located on up-diagonal j
downStatus	-(n-1) .. n-1	Whether queen is located on down-diagonal j

To position a queen at square (row,col)

Set queenCol_{row} to col
Set colStatus_{col} to ATTACKED
Set upStatus_{row+col} to ATTACKED
Set downStatus_{row-col} to ATTACKED

To check if a queen can be positioned at square (row,col)

Check if any of
colStatus_{col}
upStatus_{row+col}
downStatus_{row-col}
are set to ATTACKED.
We don't need to check if there is a queen in row row. Our algorithm never tries to place more than 1 queen in a row anyway.