

Assembly Language and Computer Architecture Lab

Week 3

Assignment 4

- Program Counter (PC) is the register that is affected by jump/branch instructions.
- Below are explanations of calculating target addresses:

Branch instruction (bne, beq)

Text Segment				
Bkpt	Address	Code	Basic	Source
<input type="checkbox"/>	4194304	0x24120001	addiu \$18,\$0,1	3: li \$s2, 1
<input type="checkbox"/>	4194308	0x24110002	addiu \$17,\$0,2	4: li \$s1, 2
<input type="checkbox"/>	4194312	0x0251402a	slt \$8,\$18,\$17	6: slt \$t0,\$s2,\$s1 # j < i ?
<input type="checkbox"/>	4194316	0x15000003	bne \$8,\$0,3	7: bne \$t0,\$zero,else
<input type="checkbox"/>	4194320	0x21290001	addi \$9,\$9,1	8: addi \$t1,\$t1,1
<input type="checkbox"/>	4194324	0x200b0001	addi \$11,\$0,1	9: addi \$t3,\$zero,1
<input type="checkbox"/>	4194328	0x08100009	j 4194340	10: j endif
<input type="checkbox"/>	4194332	0x214affff	addi \$10,\$10,-1	11: else: addi \$t2,\$t2,-1
<input type="checkbox"/>	4194336	0x016b5820	add \$11,\$11,\$11	12: add \$t3,\$t3,\$t3

The **bne** instruction has machine code **0x15000003** and address **4194316**.

Let analyse its machine code:

- ➔ Convert to binary: **0001 0101 0000 0000 [0000 0000 0000 0011]** (the last 16 bits are *jump addr*).
- ➔ Get *jump addr* and left shift 2 bits: **0000 0000 0000 0011 00**
- ➔ Sign-extended to 32 bits: **00 0000 0000 0000 0000 0000 0011 00**
- ➔ Add $4_{(10)}$ for the next instruction, we then get:
0000 0000 0000 0000 0000 0000 0001 0000₍₂₎ = $16_{(10)}$
- ➔ Add **16** to the address in base 10, we have the next instruction: **4194316 + 16 = 4194332**

Jump instruction (j)

Text Segment					Labels	
Bkpt	Address	Code	Basic	Source	Label	Address
	4194304	0x24120003	addiu \$t8,\$0,3	3: li \$s2, 3	demo.asm	
	4194308	0x24110002	addiu \$t7,\$0,2	4: li \$s1, 2	start	4194304
	4194312	0x0251402a	slt \$t8,\$t8,\$t7	6: slt \$t0,\$s2,\$s1 # j < i ?	else	4194332
	4194316	0x15000003	bne \$t8,\$0,3	7: bne \$t0,\$zero,else	endif	4194340
	4194320	0x21290001	addi \$t9,\$t9,1	8: addi \$t1,\$t1,1		
	4194324	0x200b0001	addi \$t11,\$t0,1	9: addi \$t3,\$zero,1		
	4194328	0x08100009	j 4194340	10: j endif		
	4194332	0x214affff	addi \$t10,\$t0,-1	11: else: addi \$t2,\$t2,-1		
	4194336	0x016b5820	add \$t11,\$t11,\$t11	12: add \$t3,\$t3,\$t3		

The **j** instruction has machine code **0x08100009** and address **4194328**.

Let analyse its machine code:

- ➔ Convert to binary: **0000 10[00 0001 0000 0000 0000 0000 1001]** (the last 26 bits are *jump addr*).
- ➔ Get *jump addr* and left shift 2 bits: **00 0001 0000 0000 0000 0000 1001 00 = 12₍₁₀₎**
- ➔ Current address: **4194328₍₁₀₎ = 0000 0000 0100 0000 0000 0000 0001 1000₍₂₎**
- ➔ Get first **4 bits** of current address to get 32 bits address of destination:

$$\mathbf{0000\ 00\ 0001\ 0000\ 0000\ 0000\ 0000\ 1001\ 00 = 4194340_{(10)}}$$